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EXPERT SYSTEM
IN
SOFTWARE ENGINEERING
USING
STRUCTURED ANALYSIS AND
DESIGN TECHNIQUE(SADT)

THESIS Intaek Kim Captain, ROKAF

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AFIT/GCS/ENG/90J-02

# EXPERT SYSTEM IN SOFTWARE ENGINEERING USING STRUCTURED ANALYSIS AND DESIGN TECHNIQUE(SADT)

# **THESIS**

Presented to the Faculty of the School of Engineering of the Air Force Institute of Technology

Air University

In Partial Fulfillment of the
Requirements for the Degree of
Master of Science (Computer Systems)

Intaek Kim, B.S. Captain, ROKAF

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# Acknowledgments

This thesis effort consists of designing and implementing an Expert System which checks the syntax of the Structured Analysis and Design Technique (SADT) language from a drawing model generated during the requirements analysis phase of the software life cycle.

The development of the Expert System is separated into two parts. The first part is to translate a Structured Analysis diagram into a set of predicate data forms for SADT syntax analysis. The second part is to check the syntax of a Structured Analysis language through an inference engine and a rule-base knowledge.

I wish to express my gratitude to Dr. Gary B. Lamont, my thesis advisor, for his guidance and inspiration throughout this effort. Also, I wish to thank my committee members, Dr. Thomas C. Hartrum and Dr. Frank M. Brown, for their contribution to this thesis. I also would like to thank the R.O.K. and U.S. Governments for allowing me to have this experience.

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To my father and mother.

Intack Kim

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#### Abstract

This thesis effort focuses on designing and implementing the Knowledge-Based Software Assistant System (KBSAS) for the Structured Analysis Design Technique (SADT) method developed by Softech, Inc.

A Graphics Editor 2 is used to create specific Structured Analysis (SA) diagrams and a graphical symbol syntax is derived from these diagrams. The development of the KBSAS is divided into two parts: the design and implementation of a graphics translator and an application of a knowledge-based system for syntax checking.

First, the objective of the translator is to map a subset of the graphical symbol syntax from a SA diagram into the first order predicate calculus. The SA diagram information is represented in a set of predicate data forms.

Secondly, the objective of a knowledge-based system is to evaluate adherence to proper SADT syntax. This is accomplished by generating SA rules associated with either an activity box or boundary arrows. The requirements analyst and the designer are provided with a means of recovering from a graphical symbol syntax error(s) through a display window.

Specific emphasis focuses on a comprehensive mapping of the graphical symbol syntax to predicate logic as well as development of an application of a rule-based system using this capability.

<sup>&</sup>lt;sup>1</sup>SADT is a trademark of Softech, Inc.

<sup>&</sup>lt;sup>2</sup>developed by Steven E. Johnson at the Air Force Institute of Technology.

# EXPERT SYSTEM IN SOFTWARE ENGINEERING USING STRUCTURED ANALYSIS AND DESIGN TECHNIQUE(SADT)

### I. Introduction

Background

There have been several thesis efforts in the field of Computer-Aided Software Engineering (CASE) tools to support the analysis and design stages of the software process at the Air Force Institute of Technology (AFIT). One of these efforts was a Graphics Editor for structured analysis with data dictionary support <sup>1</sup> (9).

The SAtool is one of several requirements analysis and design CASE tools based on the IDEF<sub>0</sub> syntax<sup>2</sup>. This tool provides a means of developing IDEF<sub>0</sub> diagrams and data dictionary support. This tool also saves information derived from the diagrams; however, SAtool does not have the capability of checking IDEF<sub>0</sub> syntax. To solve this problem, a validation scheme for checking the consistency of IDEF<sub>0</sub> methodology and providing error messages with error recovery is required. Using the predicate data form, a specific IDEF<sub>0</sub> diagram from the SAtool is evaluated for adherence to proper SADT syntax through the use of a knowledge-based expert system (KBES). The earlier version of the validation tool was hosted on the SUN workstations<sup>4</sup> with the expert system written in Prolog-1 and run on a Zenith Z-248 computer (10).

<sup>&</sup>lt;sup>1</sup>The Graphics Editor is called SAtool.

<sup>&</sup>lt;sup>2</sup>IDEF<sub>0</sub> is a version of SofTech's SADT.

<sup>&</sup>lt;sup>3</sup>Sometimes IDEF<sub>0</sub> diagrams are called SADT diagrams.

<sup>&</sup>lt;sup>4</sup>SUN is a trademark of SUN Microsystems, Inc.

The intent of this thesis effort is to continue the earlier investigation by expanding the rule set, thereby developing a more integrated environment and analyzing its performance.

# Statement of the Problem

The specific objectives of this thesis investigation are to extend the earlier predicate calculus definitions of SADT syntax to the more complete set of SADT constructs, to extend the expert system rule set based on the new definitions, to integrate the graphical translation process in C with the expert system on a SUN workstation, and also if time permits, to use the structure of the knowledge-based SADT syntax system to incorporate the design knowledge of a specific software application.

# Assumptions

For the purpose of this investigation, several assumptions were made.

- 1. The primary users of this tool are AFIT graduate students and faculty.
- 2. The users of this tool are familiar with the Structured Analysis methodology.
- 3. The users of this tool are familiar with the SAtool.
- 4. Although not necessary, the users of this tool are also familiar with Prolog.

#### Research Approach

The thesis objective is accomplished through the development of two major components: an IDEF<sub>0</sub> Diagram Translator and an IDEF<sub>0</sub> Syntax Expert System. The IDEF<sub>0</sub> syntax for SAtool was studied, followed by a review of the design and implementation of SAtool. The previous syntax rules of the syntax validation tool were also reviewed. These were updated and changed as necessary to reflect the development of the new system. The reusable components of the syntax validation

tool were extracted, new code was written, and new syntactical rules were generated to implement the changes considered by the analysis.

The IDEF<sub>0</sub> Diagram Translator has three parts. The function of the first part is to create a set of predicate data forms from a SAtool activity box in the IDEF<sub>0</sub> Diagram in order to check the activity IDEF<sub>0</sub> syntax. The function of the second part is also to generate a set of the predicate data forms from the current IDEF<sub>0</sub> Diagram and its parent IDEF<sub>0</sub> diagram in order to check the boundary IDEF<sub>0</sub> syntax. In the third part, the objective is to create a file for the current IDEF<sub>0</sub> diagram in the form of a set of the predicate data to speedically check the boundary IDEF<sub>0</sub> syntax.

The IDEF<sub>0</sub> Syntax Expert System consists of two major parts: an inference engine and a rule base. The inference engine was selected as backward chaining strategy (search) called BC3<sup>5</sup> which is a directed problem-solving (pattern matching) process written in prolog. The rule base consists of activity rules and boundary rules. The activity rules check the 'activity' IDEF<sub>0</sub> syntax and the boundary rules are to check the 'boundary' IDEF<sub>0</sub> syntax.

The system checks IDEF<sub>0</sub> syntax, displays error messages, and provides editing suggestions interactively.

All software conformed to the software engineering standards in AFIT's System Development Documentation Guidelines and Standards draft #4 (8).

# Materials and Equipment

The materials and equipment for this effort were provided by the AFIT Department of Electrical and Computer Engineering. The following items were used:

- 1. SUN workstations.
- 2. Berkeley Unix <sup>6</sup> version 4.3.

<sup>&</sup>lt;sup>5</sup>developed by Dr. Frank M. Brown at AFIT.

<sup>&</sup>lt;sup>6</sup>Unix is a trademark of AT&T.

- 3. Suncore graphics and Suntools environment.
- 4. The software developed in this thesis effort.
- 5. Prolog environment.

# Overview of Thesis

This thesis is divided into six chapters. Chapter I explains the history of AFIT's CASE tools based on SADT syntax and defines some of the terms to be used. Chapter II presents a literature review of the current issues that affect this thesis effort. The requirements for the translator and the expert system for this thesis effort are presented in chapter III. Chapter IV and V describe the high level and detailed design and implementation phases respectively. In chapter VI, the conclusions and the recommendations are addressed for this thesis effort.

### II. Literature Review

#### Introduction

The focus of this thesis investigation is to design and implement an application of expert system formulation for checking the syntax of IDEF<sub>0</sub> diagrams as derived from SAtool. The current SAtool is one of requirements analysis CASE tools based on the IDEF<sub>0</sub> syntax which is a subset of SADT syntax. The purpose of this literature review is to discuss the issues of the requirements analysis phase, AFIT Structured Analysis Syntax as a subset of IDEF<sub>0</sub> syntax, and a rule-based expert system architecture.

# Requirements Analysis Phase

The software life cycle represents the functionally distinct portions of development and use of a software product from birth to death. The classic life cycle model as shown in Figure 2.1 may be divided into six major phases: system engineering and analysis phase, software requirements analysis phase, design phase, implementation, testing, and finally maintenance phase (3). The requirements analysis process focuses specifically on software by definition. To understand the nature of the software to be built, the software analysts must understand the information domain for the software, as well as the required function, performance expectations, and interfacing (12). Analysts should develop the software specification using a documentation tool so that they may later compare the requirements against the solution because a complete specification of software requirements is imperative to the success of a software development effort. No matter how well-designed or well-coded, poorly specified software will disappoint the user and bring grief to the developer (12). A number of software analysis and specification methods have been developed and each method has its own notation and point of view; however, there is a set of general principles for requirements analysis:

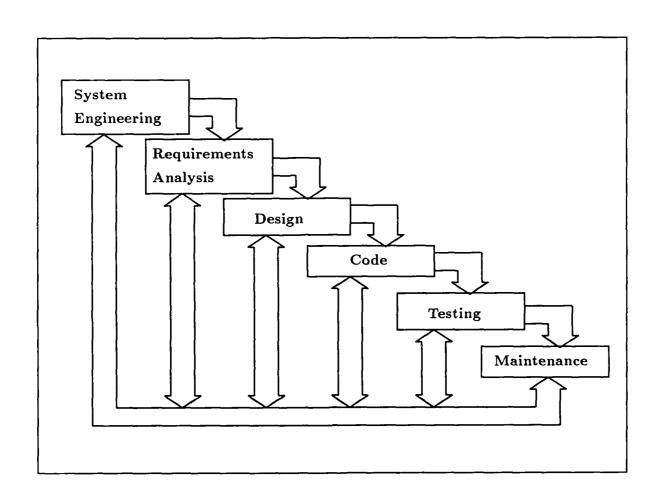


Figure 2.1. Classic Software Life Cycle Model(12:20)

- 1. The information domain and the functional domain of a problem must be represented by syntax and understood by humans.
- 2. The problem must be partitioned in a manner that uncovers detail in a hierarchical fashion.
- 3. Logical and physical representations of the system should be developed (12).

Many requirements analysis methods and tools have been developed during the past decade. The methods and tools may be divided into three broad analysis categories: data flow-oriented analysis, data structure-oriented analysis and language-based formal specification (12). The software requirements analysis methods were originally developed to be applied manually; however, each of these methods is available in a computer-aided format (12). Several of computer-aided analysis tools have been developed to automate the generation and maintenance of what was originally a manual method. These tools make use of a graphical notation for analysis. This class of tools produces diagrams, aids in problem partitioning and maintains a hierarchy of information about the system (12). These CASE tools enable the analyst to update information and compare the connections between new and existing representations of the system. For example, the SAtool enables the analyst to produce a structured analysis diagram and a data dictionary and maintain these in a data base that can be analyzed for correctness, consistency, and completeness. The computer-aided requirements analysis approach provides benefits as followings:

- improved documentation quality through standardization and reporting
- better coordination among analysts in that the data base is available to all
- gaps, omissions, and inconsistency are more easily uncovered through crossreference maps and reports
- the impact of modifications can be more easily traced
- maintenance costs for the specification are reduced (12:200).

# AFIT Structured Analysis Syntax

The SADT syntax is based on a tabulation of some 40 notations in a paper by Douglas T. Ross of Softech, Inc. (14). The notations give the definitions and the semantics of the SADT graphic language. The SADT methodology provides a means of handling large complex system problems. The SADT notations consist of two major constructs: rectangular boxes and arrows. Rectangular boxes, identified as verbs (activities), provide for the decomposition of the system parts. Arrows, labeled with nouns (data structures), represent the data flow relationship among the rectangular boxes. The rectangular boxes, arrows and English text build a diagram which represents the whole system.

The U.S. Air Force Program for Integrated Computer-Aided Manufacturing (ICAM) has developed the IDEF<sub>0</sub> (ICAM Definition Method Zero) <sup>1</sup> language. IDEF<sub>0</sub> syntax is a derivative of the SADT syntax and is used for software requirements analysis. The AFIT Structured Analysis syntax implemented by SAtool is represented in Table 2.1 (9).

Column 1 in the table shows the line numbers of the SADT graphical notations (14:20). Column 2 shows the name by which each notation was referenced in SAtool. Column 3 indicates the page in the SAtool User's Manual (9).

The SAtool provides a means of drawing the IDEF<sub>0</sub> diagrams and storing information derived from the diagrams. Each diagram is drawn and stored individually. An example of an IDEF<sub>0</sub> diagram drawn by SAtool is shown in Figure 2.2. Box1, Box2, and Box3 are for ACTIVITY NAMEs and the numbers in the rectangular boxes represent NODE NUMBERs. The numbers in small rectangular boxes show FOOTNOTEs. The label in a small circle is for TO/FROM ALL. In1, Out2, Con22, Mech1, and etc. represent line LABELs for INPUT, OUTPUT, MECHANISM, and

<sup>&</sup>lt;sup>1</sup>See the reference 17. ICAM Definition method IDEF<sub>0</sub> Sep.1979.

Ross Article Line Number	Term	User's Manual Reference
1	вох	2-2,3
2	ARROW	2-2,3
3	INPUT	3-26 (FIG)
3	OUTPUT	3-26 (FIG)
4	CONTROL	3-26 (FIG)
5	MECHANISM	3-11
6	ACTIVITY NAME	2-3,4
7	LABEL	2-3,4
12	BRANCH	3-9
13	JOIN	3-9
14	BUNDLE	6-14
15	SPREAD	6-14
18	BOUNDARY ARROW	3-17
20	DETAILED REFERENCE NUMBER	2-3
22	2-WAY ARROW	3-26 (FIG)
24	TUNNEL ARROW	2-3
25	TO/FROM ALL	6-21
27	FOOTNOTE	3-26 (FIG)
28	META-NOTE	2-3
29	SQUIGGLE	3-26 (FIG)
30	C-NUMBER	2-3
31	NODE NUMBER	2-3
32	MODEL NAME	3-26 (FIG)
33	ICOM CODE	4-8
37	FACING PAGE TEXT	4-1
38	FEO (FOR EXPOSITION ONLY)	6-5
39	GLOSSARY	2-3
40	NODE INDEX	2-3

Table 2.1. AFIT SADT syntax used by SAtool (9:A-3)

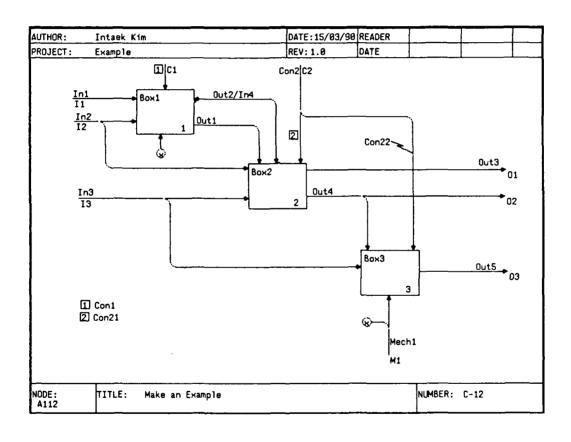


Figure 2.2. An example of IDEF<sub>0</sub> diagram

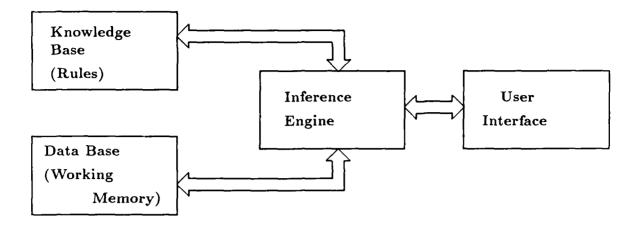


Figure 2.3. A typical Expert System Structure

CONTROL. 11, C1, O1, M1, and etc. also represent ICOM CODEs for boundary arrows.

#### Expert System

Knowledge-based expert systems are likely to be applied to requirements analysis tasks. However, the definition of the knowledge base (facts, rules, and necessary inferences to perform analysis) will remain a significant challenge in the foreseeable future. (12:201)

Figure 2.3 presents the components of a general expert system: knowledge base (rules), data base (facts; working memory), inference engine, and user interface.

The knowledge base can be represented by many different methods, such as predicate calculus, lists, frames, semantic nets, production rules, etc. In this thesis effort, the language of if-then rules was selected to represent the knowledge base, since it provides several features which are modularity, incrementability, modifiability, and transparency of the system. The if-then rule consists of two parts: condition and conclusion. The logical condition part may contain one or more premises linked by the conjunctions AND, OR, or NOT. If the conditions are true (met), the conclusion part is also true (fired).

The data base is a portion of working memory where the current status of problem is stored. Initially, the lists of object, attribute, and value (OAV) triple derived from a IDEF<sub>0</sub> diagram are stored. Then the new lists of OAV triple from the inference process are added. The data base also stores a list of rules that have been examined, and fired in some order. The rule order can be given later if the user or developer requires an explanation of the reasoning process.

The inference engine is called a rule interpreter because its operation is like a software interpreter for a computer language. The rule interpreter examines the rules in a specific order searching for matches to the initial and current conditions given in the working memory. As the rules continue to fire, they will reference one another and form an inference chain. The firing of a rule may add new facts to the working memory, which gives the rule interpreter additional information on which to proceed. This process continues until the solution is found.

Given a desired goal, there are two basic approaches in searching for a solution: forward chaining and backward chaining. In forward chaining, the rule interpreter tries to match a fact in the working memory to the situation stated in the condition part. If a fact in the working memory has been matched, then that rule is fired. The conclusion part could generate a new fact that is stored in the knowledge base. This new fact may be used in the search for the next proper rule. This process continues until the solution of the given goal is satisfied. In backward chaining, the rule interpreter starts examining the conclusion part to look for a match. If a match is found, then the working memory is updated recording the conditions that the rule stated as necessary for supporting the matched conclusion (13). The backward chaining process continues with the system repeatedly attempting to match the conclusion part against the current system's status. The corresponding condition parts matched are used to produce new intermediate goal states which are recorded in the working memory. This process continues until the given goal is proved.

Finally, the user interface provides a means of communication between the

user and the system. The user interface asks questions or presents menu choices for entering initial facts in the data base. Any intermediate communications during the problem-solving process are taken care of by the user interface.

Expert systems are far more useful if they have the following additional features:

# • Modularity:

Each rule defines a small, relatively independent piece of knowledge.

## • Incrementability:

New rules can be added to the knowledge base relatively independently of other rules.

- Modifiability (as a consequence of modularity):
   Old rules can be changed relatively independently of other rules.
- Support system's transparency (4:316-317).

#### Summary

This thesis effort concentrates on translating an IDEF<sub>0</sub> diagram drawn by the SAtool into a set of predicate data forms during requirements analysis. It also focuses upon developing an application of a rule-based expert system for evaluating IDEF<sub>0</sub> syntax. The literature review in this chapter provides understanding concerning the main subjects of this thesis investigation: requirements analysis phase of software development, IDEF<sub>0</sub> syntax, SAtool, and rule-based expert system components.

# III. System Requirements Analysis

#### Introduction

This thesis investigation is classified into two major categories. First, the  $IDEF_0$  Diagram Translator is to be redesigned and reimplemented to translate any  $IDEF_0$  diagrams drawn by SAtool into a set of predicate data forms. It should create a necessary file to be used for checking  $IDEF_0$  syntax. The second category is to design and implement the  $IDEF_0$  Syntax Expert System which is an application of a knowledge-based expert system.

This chapter presents the considerations related to the development of the IDEF<sub>0</sub> Diagram Translator requirements, IDEF<sub>0</sub> Syntax Expert System requirements, formalization criteria, the IDEF<sub>0</sub> diagrams of this tool, and validation test requirements.

#### Considerations of the Previous Studies

The SAtool provides an interactive graphics editor for drawing IDEF<sub>0</sub> diagram and a means of generating data dictionary information (9:3-2). The SAtool also provides the capability to check IDEF<sub>0</sub> syntax for an activity box in any IDEF<sub>0</sub> diagrams(10:2-1). The SAtool does not provide a means of checking IDEF<sub>0</sub> syntax for any boundary arrows with the parent IDEF<sub>0</sub> diagram, however, the SAtool provides a means of checking IDEF<sub>0</sub> syntax for only one activity box in any IDEF<sub>0</sub> diagram. The current SAtool checks IDEF<sub>0</sub> syntax through Zenith Z-248 workstations using the expert system written in Prolog-1. The improved tool in this thesis effort should interface with the SAtool, produce a set of predicate data forms for the activity boxes and the boundary arrows information in any IDEF<sub>0</sub> diagrams, and provide the more completed capability of checking IDEF<sub>0</sub> syntax. Also, the revised tool should satisfy all requirements of the previous SAtool (9).

# IDEF<sub>0</sub> Diagram Translator Requirements

The function of the IDEF<sub>0</sub> Diagram Translator is to generate the predicate data forms from any IDEF<sub>0</sub> diagrams. The IDEF<sub>0</sub> Diagram Translator should act as a bridge between the current SAtool and the IDEF<sub>0</sub> Syntax Expert System. The current SAtool was written in the C language and used graphics software packages called SunView and SunWindow environment. Thus, the IDEF<sub>0</sub> Diagram Translator should operate with the current SAtool. Since the current SAtool provides a means of checking IDEF<sub>0</sub> syntax for only one activity box, to check syntax for boundary arrows, the IDEF<sub>0</sub> Diagram Translator should be redesigned and reimplemented in the C language. The predicate data forms generated by the IDEF<sub>0</sub> Diagram Translator should be the initial data base for the IDEF<sub>0</sub> Syntax Expert System. It is necessary to formalize the IDEF<sub>0</sub> syntax to produce the predicate data forms and to check the IDEF<sub>0</sub> syntax in the IDEF<sub>0</sub> diagrams. IDEF<sub>0</sub> syntax is formalized as follows:

- The formal definition of IDEF<sub>0</sub> syntax must contain the syntax information in any IDEF<sub>0</sub> diagram and be described syntactically,
- 2. provide the means to determine syntax errors in any IDEF<sub>0</sub> diagram,
- 3. provide a domain where the definition of consistency can be given,
- 4. serve as the final arbiter in cases where there is disagreement concerning the exact meaning of the representation, and
- 5. should be able to be implemented in a computer system (10:2-3).

The next section, Requirements analysis Diagrams, describes the functional decompositions for the IDEF<sub>0</sub> Diagram Translator.

# IDEF<sub>0</sub> Syntax Expert System Requirements

The IDEF<sub>0</sub> Syntax Expert System should allow the user to check the activity IDEF<sub>0</sub> syntax and the boundary IDEF<sub>0</sub> syntax in any IDEF<sub>0</sub> diagrams using a created predicate data file.

The backward chaining search is useful when there are many more rules than desired goals. A backward chaining inference engine was selected called BC31 which directed problem-solving processes and acts as a rule interpreter (6) because the backward chaining strategy is useful when there are many more rules than desired goals. The rule base should be able to support the knowledge of IDEF<sub>0</sub> syntax with the inference engine in accordance with the activity boxes and the boundary arrows. To simplify the rule base, the rule base should be consisted of two separate parts because the boundary arrows information is not necessary in checking the activity IDEF<sub>0</sub> syntax and the activity boxes information is not needed for examining the boundary IDEF<sub>0</sub> syntax. The first part, called the Activity IDEF<sub>0</sub> Syntax rule base, should allow the checking of the Activity IDEF<sub>0</sub> Syntax using a created predicate file which includes all information pertaining to an activity box in any IDEF<sub>0</sub> diagram. The second part, called the Boundary IDEF<sub>0</sub> syntax rule base, should allow for the checking of the Boundary IDEF<sub>0</sub> Syntax using a created predicate file which includes all information pertaining to the boundary arrows in any IDEF<sub>0</sub> diagram and its parent IDEF<sub>0</sub> diagram.

## Validation Test Requirements

Some parameters can be used for evaluating the conformity of the requirements with the tool. As mentioned in Chapter I, the important parameters are the accuracy with which the tool checks the IDEF<sub>0</sub> syntax and the capability of which the tool interactively displays error messages and editing suggestions. Other parameters to be considered are user friendliness, maintainability, compatibility, and consistency.

<sup>&</sup>lt;sup>1</sup>developed by Dr. Frank M. Brown at AFIT.

# Requirements Analysis Diagrams

This section presents the functional model which defines and describes the functional decompositions for the IDEF<sub>0</sub> Diagram Translator and the IDEF<sub>0</sub> Syntax Expert System discussed in the previous sections. The IDEF<sub>0</sub> diagrams in this section are based on the analysis of processes or activities and illustrate one level of the functional decomposition with the facing page text. The facing page text provides additional information that is not easily inferred from the diagram. The Data Dictionary entries provide even more detailed information in accordance with each activity and data item (22:4-5).

Figure 3.1 shows the top most level IDEF<sub>0</sub> diagram for the overall system of the SAtool. The purpose of the *Provide SAtool* function is to create and edit an IDEF<sub>0</sub> diagram, its data dictionary information, and the facing page text interactively (9). The function also involves the process which produces the predicate data forms to be used for the knowledge base of the IDEF<sub>0</sub> Syntax Expert System.

Figure 3.2 displays the first decomposition of the top most level of *Provide SAtool* activity. This decomposition shows the two primary functions: *Provide SA Editor* and *Translate Diagram*. The *Provide SA Editor* function is to draw the activity IDEF<sub>0</sub> diagram and to generate its data dictionary information and its pacing page text for activity and data entry. The *Translate Diagram* activity provides a means of translating IDEF<sub>0</sub> diagrams into predicate data forms.

Figure 3.3 shows the decomposition of the Translate Diagram activity into three functional components: Translate Activity, Translate Boundary, and Save Diagram. The Translate Activity provides a means of translating an activity box in any IDEF<sub>0</sub> diagrams into a set of predicate data forms through the translation rules of the activity box and saving it into a temporary file. The Translate Boundary operation is the process which translates the boundary arrows in any IDEF<sub>0</sub> diagrams into a set of predicate data forms through the translation rules of the boundary arrows and saves it into a temporary file. Finally, Save Diagram provides a means of translating

#### A-0 Provide SAtool

Abstract: Provide SAtool provides a means of mechanism by which the user is able to draw Activity IDEFO Diagrams. From these diagrams, Facing Page Text and Data Dictionaries for Activities and Data and Predicate Data forms are generated.

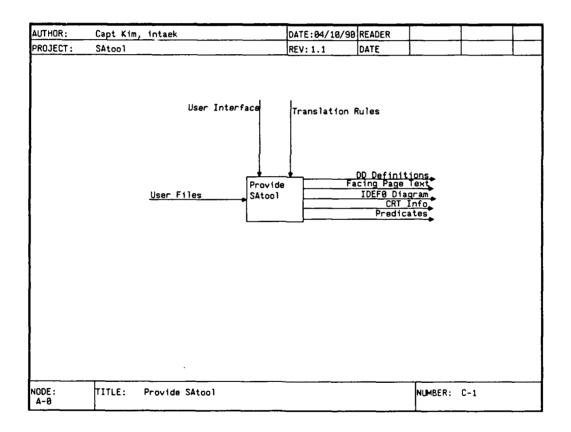


Figure 3.1. Provide SAtool

#### A0 Provide SAtool

Abstract: Provide SAtool provides the user a mechanism by which the user is able to draw Activity IDEFO Diagrams. From these diagrams, Facing Page Text, Data Dictionaries for Activities and Data, and Predicate Data forms are generated.

Al Provide SA Editor provides a means of drawing Activity IDEFO Diagrams. From these diagrams, Facing Page Text and Data Dictionaries for Activities and Data are generated.

A2 Translate Diagram provides a means of translating IDEFO Diagrams into Predicate Data Forms.

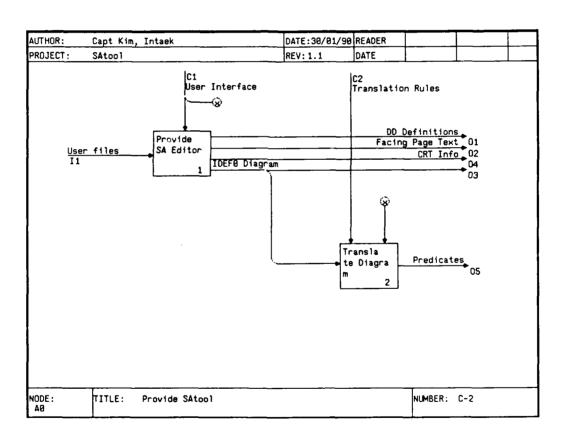


Figure 3.2. Provide SA Editor

the IDEF<sub>0</sub> diagram into a set of predicate data forms and saving it into a file which is specified by the user. Further functional decompositions are presented in Appendix A.

# Summary

This chapter presented the requirements analysis for the development of IDEF<sub>c</sub> Diagram Translator and IDEF<sub>o</sub> Syntax Expert System. Since this investigation extends the earlier version of the SAtool, this tool should satisfy all requirements of the earlier version. This tool should also provide a means for displaying error messages and editing suggestions. The functional decompositions for this tool was presented in *Requirements Analysis Diagrams* section.

The next two chapters use the requirements developed in this chapter as the fundamental for designing, implementing, and testing of IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System.

#### A2 Translate Diagram

Abstract: Translate Diagram provides the user a means of translating IDEFO Diagrams into Predicate Data Forms.

- A21 Translate Activity provides a means of translating an activity box in any IDEFO Diagrams into a set of predicate data forms through the translation rules of the activity box.
- A22 Translate Boundary provides a means of translating the boundary arrows in any IDEFO Diagrams into Predicate data forms through the translation rules of the boundary arrows.
- A23 Save Diagram provides a means of translating the IDEFO Diagram into a set of the predicate data forms and saving it into a file which user specifies interactively.

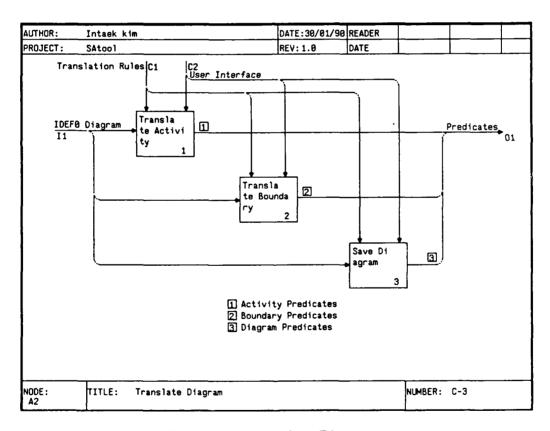


Figure 3.3. Translate Diagram

# IV. High Level Design

#### Introduction

The purpose of this chapter is to present and justify the preliminary software design for the IDEF<sub>0</sub> Diagram Translator (IDT) and the IDEF<sub>0</sub> Syntax Expert System (ISES). Throughout the remainder of this investigation, IDT and ISES refer to these two particular systems. Preliminary design is associated with the transformation of requirements into the software architecture. The transformation starts with several considerations of previous studies addressed in Chapter III. The modified screen layout and menu selection are then presented. In addition, the main functions of the IDEF<sub>0</sub> Diagram Translator, the IDEF<sub>0</sub> Syntax Expert System, and the associated components are introduced.

# Previous Study Considerations

Since the tool should interface with the SAtool, the hardware and the software to be used are already chosen. Thus, the Sun3 and the Sun4 workstations using the SunOS and the SunView window-based environment are required for this tool because the SAtool was developed through the Sun workstation. Also since the IDEF<sub>0</sub> validation tool was implemented with the C language in order to translate the IDEF<sub>0</sub> diagram into the predicate data forms, the C language is used to expand the capability of the translating process. This decision is reasonable because many portions of the validation tool and the SAtool could be directly reused. Appendix C represents a summary of the data structures which are used for the earlier version of SAtool and this thesis effort.

# Screen Layout and Menu System

In this thesis investigation, the screen layout of the SAtool and the validation tool should be modified by adding new menu items for checking IDEF<sub>0</sub> syntax.

Figure 4.1 is a picture of the actual screen layout used by the tool.

The actual screen is divided into five windows: the Input Window, the Message Window, the Selection Window, the Diagram Window, and the Data Dictionary Window in vertical order. The function of each window is the same as in the SAtool except the function of the Selection Window. The Selection Window, located directly below the Message Window is used for selecting the menu which contains the next desired operation. The six ovals arranged in left to right order are: RECALL DGM, EDIT DD, EDIT FPT, MISC FUNC, SAVE DGM, and CHECK SYNTAX. The function of RECALL DGM is to read a previously saved IDEF<sub>0</sub> diagram. The EDIT DGM function is to create and edit an IDEF<sub>0</sub> diagram according to its attribute menus. The function of EDIT DD is to create and edit data dictionary information. The function of EDIT FPT is to edit facing page text of an IDEF<sub>0</sub> diagram. The function of MISC FUNC is to save an IDEF<sub>0</sub> diagram, change directory, start new diagram, and exit SAtool. The function of SAVE DGM is to save the graphics information (.gph extension) and the data dictionary information (.dbs extension) in files under the name provided by the user. Finally, The CHECK SYNTAX function is to translate and save IDEF<sub>0</sub> diagrams as predicate data forms and to check IDEF<sub>0</sub> syntax. Figure 4.2 presents a picture of all menu selections to be used and their decompositions.

Since the functions of all menu selections are the same as those of the SAtool except Activity, Boundary, and Save (.pro) selections, the other detailed descriptions of the functions of the menu selections except for above three are available in reference (9). Further descriptions of Activity, Boundary, and Save (.pro) selection functions are discussed in the next section.

# IDEF<sub>0</sub> Diagram Translator

The translator for the IDEF<sub>0</sub> diagram is used to translate the IDEF<sub>0</sub> graphical features into predicate data forms. The requirements for the formalization criteria

SA TOOL											
INPUT:	DISABLED										
MESSAGE:	WELCOME, P	lease m	nake a	selec	tion.						
RECALL I	DOM ENT	DGM )	EDIT	00	EDIT	FPT	MISC FUNC	SAVE DOM	CHECK	SYNTAX	
AUTHOR:							DATE:	READER		Ì	
PROJECT:							REV:	DATE		<u> </u>	
NODE:	TITLE:						· · · · · · · · · · · · · · · · · · ·		NUMBER:		

Figure 4.1. Screen Layout

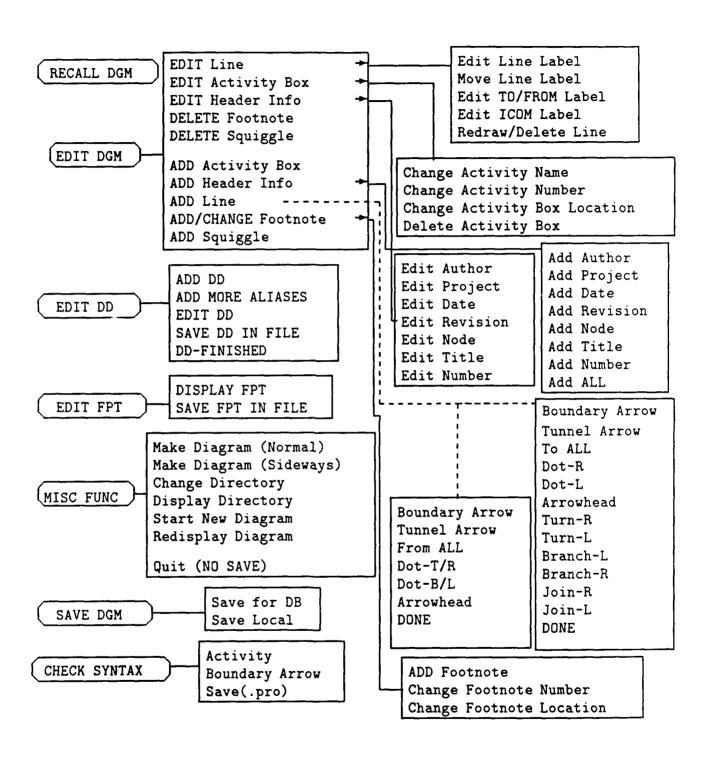


Figure 4.2. SAtool Menus

have been discussed in Chapter III. There are many ways to formalize the graphical language such as PSL/PSA, TAGS, SREM, and etc (12:163-209). Also, there are lots of different ways to represent knowledge for expert system, for example rules, semantic nets, frames, scripts, predicate logic, and others (13:63-102). In this thesis investigation, predicate logic is used to translate an IDEF<sub>0</sub> diagram into a set of predicate data forms. The predicate logic is often used as a means of knowledge representation in expert systems and is the basis for logic programming. Many specialists ragard it as the single most important knowledge representation method. The predicate logic also provides for symbols to represent objects and then allows these symbols to be used as components of statements. As shown in Figure 4.3, the IDEF<sub>0</sub> graphical notations used by the SAtool consist of two major constructs: activity box and arrow. The informations related to an activity box are ACTIVITY NAME, AC-TIVITY NUMBER, INPUT, OUTPUT, CONTROL, and MECHANISM. Although there are many types of arrows such as BRANCH, JOIN, SPREAD, BUNDLE, and others, the arrow type can be considered to be one of either the boundary arrow type connected on the parent IDEF<sub>0</sub> diagram, the tunnel arrow type, or the interboxes arrow type which is connected between two boxes in the same IDEF<sub>0</sub> diagram. Among the arrow types, only the boundary arrow type has the information of the relationships between an IDEF<sub>0</sub> diagram and its parent IDEF<sub>0</sub> diagram. This boundary arrow type can be considered as either boundary INPUT, OUTPUT, CONTROL, or MECHANISM. The other arrow types are related to an activity box in an IDEF<sub>0</sub> diagram. Thus, it is useful to translate and create separately the graphical informations for an activity box and the boundary arrows. Another reason for translating the IDEF<sub>0</sub> diagram into a set of predicate data forms and generating seperately its information is to reduce the size of the predicate data file for checking IDEF<sub>0</sub> syntax and to provide simplicity of the knowledge base as well. Therefore, the function of the IDEF<sub>0</sub> Diagram Translator is divided into three parts. The function menus of CHECK SYNTAX in Figure 4.2 show the functions of the IDEF<sub>0</sub> Diagram Transla-

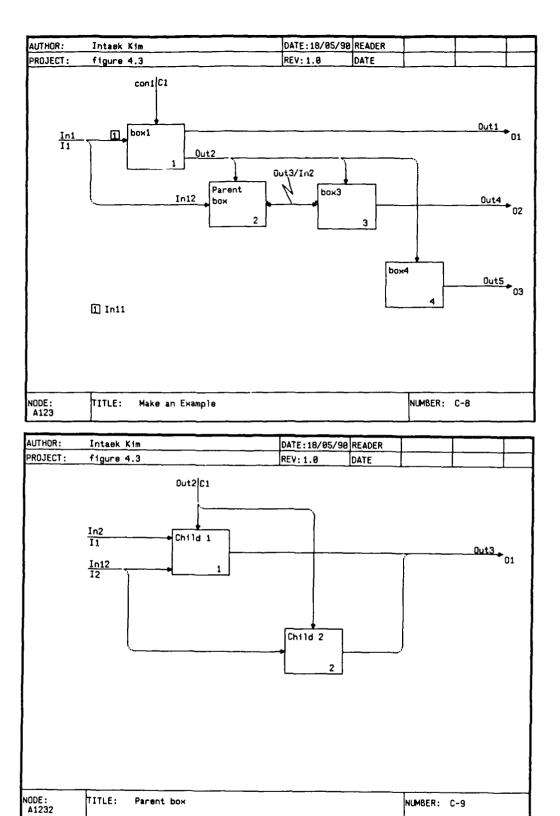


Figure 4.3. IDEF<sub>0</sub> Diagram with Parent

tor. The Activity selection menu is used to translate and create a file of the predicate data forms automatically from an activity box clicked on by the user in any IDEF<sub>0</sub> diagrams. The Boundary Arrow selection menu is used to translate and create automatically a file of the predicate data forms from the boundary arrow information in the current IDEF<sub>0</sub> diagram and the predicate data file of its parent IDEF<sub>0</sub> diagram. Finally, Save (.pro) selection menu is used to translate an IDEF<sub>0</sub> diagram into a set of predicate data forms and to save it to a file that the user specifies. This saved file is to be used for checking boundary IDEF<sub>0</sub> syntax with the next decomposed IDEF<sub>0</sub> diagram

# IDEF<sub>0</sub> Syntax Expert System Components

The IDEF<sub>0</sub> Syntax Expert System provides a means of checking the IDEF<sub>0</sub> syntax in any IDEF<sub>0</sub> diagrams using a created predicate data file. The components of ISES are the inference engine (control strategy) which is selected as BC3, the knowledge base (rule base) which consists of the IDEF<sub>0</sub> syntax rules, the data base (working memory) which is made up of a list of facts derived from the IDEF<sub>0</sub> diagram, and the user interface which supplies facts or other information to the ISES and transmits expert advise to the user as shown in Figure 2.3.

Inference Engine. An inference engine applies the knowledge to the solution of a specific problem. In general, the same control strategy can be used to reason out all kinds of actual problems because it is separated from the knowledge base for the particular application. One of the inference rules, called modus ponens, is used in producing a proof which allows a fact or truth to be inferred from two other statements. For example, if propositions P and P implies Q are true, then proposition Q is true. The inference engine repeatedly applies the method of modus ponens to the knowledge base to extract a specific value or solve a particular problem. During consultation of the IDEF<sub>0</sub> Syntax Expert System, the following functions should be performed:

- interface with users
- obtain and load the knowledge base
- apply the rules in the knowledge base and the facts in the working memory to achieve goals
- ullet display the messages about the results of checking IDEF $_0$  syntax.

In Chapter II, the backward chaining control structure has been represented. Backward chaining is often beneficial when there are many more facts than final goals (thus called goal-driven reasoning). Since the number of facts in any IDEF<sub>0</sub> diagram is many more than that of goals for IDEF<sub>0</sub> Syntax, backward chaining is useful for the ISES design.

Knowledge Base. The heart of the expert system is the knowledge base, which contains the problem-solving knowledge of the particular application. In the IDEF<sub>0</sub> Syntax Expert System, the knowledge base is represented in the form of if...then rules and is separated into two parts: the knowledge base of the IDEF<sub>0</sub> syntax for an activity box and the knowledge base of IDEF<sub>0</sub> syntax for boundary arrows. The former includes the IDEF<sub>0</sub> syntax rules and goals about an activity box focused on ACTIVITY NAME, ACTIVITY NUMBER, INPUT, OUTPUT, and MECHANISM. The latter is made up of the IDEF<sub>0</sub> syntax rules and goals about boundary arrows focused on boundary INPUT, OUTPUT, CONTROL, and MECHANISM. Since the IDEF<sub>0</sub> syntax has semantic meanings, it is difficult to represent completely the IDEF<sub>0</sub> syntax to the knowledge base. For example, ACTIVITY NAME should be used in the form of a verb, which needs rules as many as the number of verbs in the dictionary.

In the previous section,  $IDEF_0$  Diagram Translator, the reasons why two different predicate data files are generated separately have been discussed. Since the knowledge base should be consistent with the  $IDEF_0$  Diagram Translator, the

knowledge base for IDEF<sub>0</sub> Syntax Expert System should be separated into two sub-components. It is also necessary to separate it into two sub-components in order to reduce the complexity and redundency of the knowledge base because while checking IDEF<sub>0</sub> syntax for an activity box, the knowledge base for boundary arrows is unnecessary. The names of the two separated portions of the knowledge base are the *ActivitySArule* for the activity IDEF<sub>0</sub> syntax rules, and the *BoundarySArule* for the boundary arrow IDEF<sub>0</sub> syntax rules. A detailed discussion of IDEF<sub>0</sub> syntax rules follows in Chapter V.

Data Base (Working Memory). The data base contains a broad range of information about the current status of the problem being solved. The temporary output files of the IDEF<sub>0</sub> Diagram Translator become the initial data base for the IDEF<sub>0</sub> Syntax Expert System in accordance with checking IDEF<sub>0</sub> syntax of the activity box or the boundary arrow. While checking the IDEF<sub>0</sub> syntax, the data base also contains a list of rules that have been examined and fired. After checking IDEF<sub>0</sub>, the sequence of the rules fired can be given in order to explain the reasoning process.

User Interface. The user interface allows the user to communicate with the expert system and also provides the user with an insight into the problem-solving process carried out by the inference engine. There are several ways to communicate with the expert system such as questions and answers, menu choices, statements, and others. In the ISES, the user interface facility as a piece of software is contained in the inference engine and provides a means of asking questions and answering through the verbose and why trace operation. It also provides menu choices in order to check IDEF<sub>0</sub> syntax for activity box or boundary arrow.

# Testing Techniques

There are two common techniques to test software referred to as black box and white box testing (12:470).

Black box testing enables the software engineer to derive sets of input conditions that will fully test all functional requirements for a program. This attempts to find errors in the following categories:

- incorrect or missing functions,
- interface errors,
- errors in data structures,
- performance errors, and
- initialization and termination errors (12:484).

White box testing is a test method based on the control structure of the procedural design. This is used to derive the following test cases:

- guarantee that all independent paths within a module have been exercised at least once,
- exercise all logical decisions on their true and false sides,
- execute all loops at their boundaries and within their operational bounds, and
- exercise internal data structures to assure their validity (12:472).

As mentioned earlier, the black box testing method is useful at this point because it focuses on the functional requirements of the software. The functions defined in the requirements analysis phase are compared to the requirements specification to be sure that all requirements are satisfied. In the case of the IDEF<sub>0</sub> Diagram Translator, the black box test can be applied. For example, does the IDT translate IDEF<sub>0</sub> diagram into a set of predicate data forms which contains all syntactical information of the IDEF<sub>0</sub> diagram? In the case of the IDEF<sub>0</sub> Syntax Expert System, the black box test can also applied. For example, does the ISES contains all syntactical rules of IDEF<sub>0</sub> diagram comparision with the requirements specification.

# Summary

This chapter presented a high level software design for the IDEF<sub>0</sub> Diagram Translator and the IDEF<sub>0</sub> Syntax Expert System. To be consistent with the earlier version of IDT, several considerations of the earlier version are addressed. The screen layout and the menu system modified were presented. In addition to, the main functions of the IDEF<sub>0</sub> Diagram and the components of the IDEF<sub>0</sub> Syntax Expert are also addressed, and the test design was introduced. The next chapter presents low level design, implementation, and software test for the IDEF<sub>0</sub> Diagram Translator and the IDEF<sub>0</sub> Syntax Expert System.

# V. Low Level Design and Implementation

#### Introduction

This chapter discusses the low level design and implementation of an IDEF<sub>0</sub> Diagram Translator and an IDEF<sub>0</sub> Syntax Expert System specified in the previous design chapter. The IDT becomes a portion of the SAtool and translates IDEF<sub>0</sub> diagrams into the predicate data forms. These predicate data forms are used for checking IDEF<sub>0</sub> syntax and become the initial data base (working memory) of the IDEF<sub>0</sub> Syntax Expert System.

# IDEF<sub>0</sub> Diagram Translator Design

The flow diagram model of the IDEF<sub>0</sub> Diagram Translator is shown in Figure 5.1. There are three components in Figure 5.1: Translate Activity, Translate Boundary, and Save Diagram for the IDT. These three components accept an IDEF<sub>0</sub> diagram, User's Choice, and Parent Predicate as inputs and generate outputs such as Activity Predicate, Boundary Predicate, and Diagram Predicate applied by Translation Rules. The function of Translate Activity is to translate an activity box in any IDEF<sub>0</sub> diagram into the predicate data forms and produce a file of the predicate data forms. The file name of the predicate data forms for an activity box is CHECKBOX.PRO (temporary file). The function of Translate Boundary is to translate boundary information in any IDEF<sub>0</sub> diagram and the parent diagram related to the current IDEF<sub>0</sub> diagram into the predicate data forms and generate a file of the predicate data forms which is a temporary file. The file name of the predicate data forms for the boundary information is CHECKBOUNDARY.PRO (temporary file). The function of Save Diagram is to translate the IDEF<sub>0</sub> diagram into the predicate data forms and save into a file of \*.pro. The symbol \* is a name which the user enters. The \*.pro file is used to check IDEF<sub>0</sub> syntax about the boundary arrows. Appendix B shows the structure charts for the IDEF<sub>0</sub> Diagram Translator implementation.

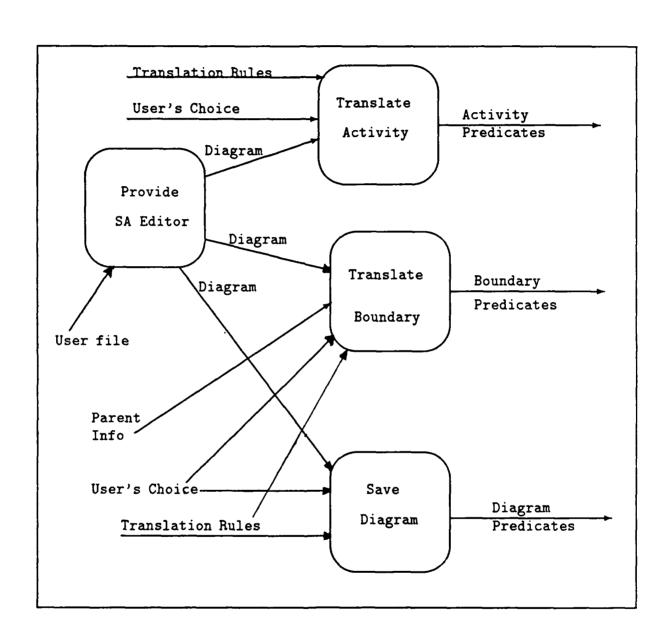


Figure 5.1. Flow Diagram for  $IDEF_0$  Diagram Translator

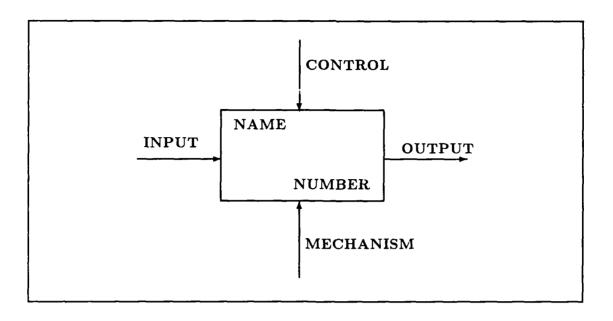


Figure 5.2. A Typical Activity Box

The earlier version of the SAtool was implemented in C. To reuse many modules of the earlier SAtool without modifications or with slight modifications, a decision was made to proceed using C language for the IDEF<sub>0</sub> Diagram Translator.

#### Translation Rules

Since the IDEF<sub>0</sub> Diagram Translator consists of three components, it is reasonable to discuss the translation rules according to Activity and Boundary.

Activity. A typical drawing model for an activity box is shown in Figure 5.2. The IDEF<sub>0</sub> Diagram Translator should produce Activity Predicate data forms which include all graphical information of an activity box. As discussed in Chapter IV, the information of an arbitrary activity box, which is based on NAME, NUMBER, INPUT, OUTPUT, CONTROL, and MECHANISM is shown in Figure 5.2. Among the information, the NAME is a key information because the other information depends on the NAME. For example, if the NAME is known among activity boxes, the other information can be extracted easily. Table 5.1 presents the translation

Term	Predicate	Triple
NAME	activityname(is, NAME)	[activityname, is, NAME]
NUMBER	NAME(number_is, NUMBER)	[NAME, number_is, NUMBER]
INPUT	NAME(input_is, LABEL)	[NAME, input_is, LABEL]
OUTPUT	NAME(output is, LABEL)	[NAME, output is, LABEL]
CONTROL	NAME(control is, LABEL)	[NAME, control is, LABEL]
MECHANISM	NAME(mechanism_is, LABEL)	[NAME, mechanism is, LABEL]
NUMBER OF INPUTS	NAME(has_input_number, COUNT)	[NAME, has_input_number, COUNT]
NUMBER OF OUTPUTS	NAME(has_control_number, COUNT)	[NAME, has_output_number, COUNT]
NUMBER OF CONTROLS	NAME(has_control_number, COUNT)	[NAME, has_control_number, COUNT]
NUMBER OF MECHANISM	NAME( has_mechanism_number, COUNT)	[NAME, has_mechanism_number, COUNT]

Table 5.1. Translation Rules for Activity Box

rules for an activity box. Column 1 in the table 5.1 presents the information items of an activity box. Column 2 shows predicate data forms of the items. Since the predicate data forms produced by IDEF<sub>0</sub> Diagram Translator become the initial data base of the IDEF<sub>0</sub> Syntax Expert System, each item of the predicate data form should be represented by a three-element list of the triple form: [Object, Attribute, Value]. Column 3 shows the triples of the items. The triple form is used as the actual input of the IDEF<sub>0</sub> Syntax Expert System. The activity box name, NAME, is translated into the predicate activityname(is, NAME), which means activity box is named as NAME. The predicate NAME(number\_is, NUMBER) for NUMBER means the number of activity box NAME is NUMBER. In the case of INPUT, it is translated into the predicate NAME(input\_is, LABEL), which means the input of

activity box NAME is LABEL. Similarly, in the case of OUTPUT, CONTROL, and MECHANISM, they are easily translated into the predicate data forms as shown in the table 5.1. In the case of NUMBER OF INPUTS, it is translated into the predicate NAME(has\_input\_number, COUNT), which means activity box NAME has COUNT INPUTS. This information is used for checking the number of inputs which is limited and the boundary arrow. Also, NUMBER OF OUTPUTS, NUMBER OF CONTROLS, and NUMBER OF MECHANISMS are translated similarly.

Boundary. The type of a boundary arrow is one of either input, output, control, or mechanism. These are the touched arrows of an activity box as well. Since the boundary arrows should be related to the current IDEF<sub>0</sub> diagram and an activity box of the parent IDEF<sub>0</sub> diagram, the predicate information of the activity box should have been saved already. Table 5.2 shows the translation rules for boundary arrows. BOUNDARY INPUT in column 1 is translated into the predicate boundary-input(is, LABEL) in column 2, which means LABEL is a boundary input. The triple [boundary\_input, is, LABEL] in column 3 represents the actual data base form for the IDEF<sub>0</sub> Syntax Expert System. BOUNDARY OUTPUT, BOUNDARY CONTROL, and BOUNDARY MECHANISM are translated similarly as shown in Table 5.2. NUMBER OF BOUNDARY INPUTS is translated into the predicate boundary-input(has\_number, COUNT), which means the number of boundary inputs is COUNT. In the case of NUMBER OF BOUNDARY OUTPUTS, NUMBER OF BOUNDARY CONTROLS, and NUMBER OF BOUNDARY MECHANISMS, they are each translated as shown in Table 5.2.

# IDEF<sub>0</sub> Syntax Expert System

The detailed structure of the IDEF<sub>0</sub> Syntax Expert System is shown in Figure 5.3. Since the user interface portion is contained in the inference engine as discussed in Chapter IV, this is not mentioned in this chapter. Thus, the IDEF<sub>0</sub> Syntax Expert System consists of three major components: Rule Base, Working Memory,

Term	Predicate	Triple
BOUNDARY INPUT	boundary_input (is, LABEL)	[ boundary_input , is, LABEL ]
BOUNDARY OUTPUT	boundary_output (is, LABEL)	[ boundary_output , is, LABEL ]
BOUNDARY CONTROL	boundary_control (is, LABEL)	[ boundary_control , is, LABEL ]
BOUNDARY MECHANISM	boundary_mechanism (is, LABEL)	[ boundary_mechanism , is, LABEL ]
NUMBER OF BOUNDARY INPUTS	boundary_input (has_number, COUNT)	[ boundary_input , has_number, COUNT ]
NUMBER OF BOUNDARY OUTPUTS	boundary_output (has_number, COUNT)	[ boundary_output , has_number, COUNT ]
NUMBER OF BOUNDARY CONTROLS	boundary_control (has_number, COUNT)	[ boundary_control , has_number, COUNT ]
NUMBER OF BOUNDARY MECHANISMS	boundary_mechanism (has_number, COUNT)	[ boundary_mechanism , has_number, COUNT ]

Table 5.2. Translation Rules for Boundary Arrows

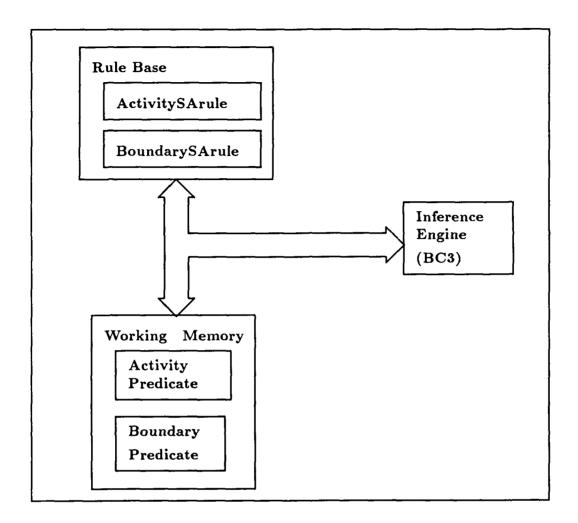


Figure 5.3. Structure of IDEF<sub>0</sub> Syntax Expert System

and Inference Engine as shown in Figure 5.3.

The detailed discussion of the Rule Base is given in the next section, Rule Base. The Working Memory (Data Base) is initially set up by the predicate output of the IDEF<sub>0</sub> Diagram Translator. The predicate output is a collection of fact tripples discussed in the previous section, Translation Rules. The inference engine, BC3, which directed problem-solving processes and acted as a rule interpreter, was available. BC3 is a shell for a backward chaining expert system. Since the backward chaining strategy is good when there are many more facts than final goals, BC3 is suitable for use as the inference engine of the IDEF<sub>0</sub> Syntax Expert System. Also,

since BC3 was originally used on the Zenth Z-248 workstations, in order to run it on the Sun workstations, BC3 should be modified. The modified BC3 is listed in appendix F.

BC3 was implemented in Prolog-1 which is a dialect of many Prolog languages and is for the personal computer. To reuse BC3 with slight modification, Quintus Prolog, which is available on the Sun workstations, was selected to implement the inference engine for the ISES.

## Rule Base

The inference engine (BC3) applies each element of the rule base to the solution of the specific domain. The specific domain is divided into two categories, one for an activity box and the other for the boundary arrows in any IDEF<sub>0</sub> diagram. Thus the rule base (IDEF<sub>0</sub> Syntax) is separated into two parts: Activity IDEF<sub>0</sub> Syntax and Boundary Arrow IDEF<sub>0</sub> Syntax.

Activity IDEF<sub>0</sub> Syntax. The Activity IDEF<sub>0</sub> Syntax focuses on an activity box in any IDEF<sub>0</sub> diagram as shown Figure 5.2. The following list in English sentences is extracted from Figure 5.2 for Activity IDEF<sub>0</sub> Syntax.

- An activey box must have a name.
- An activity box must have a number except for the top-most level activity box.
- An activity box must have at least a touched control arrow and a touched output arrow.
- If an activity box has touched arrows, the arrows must have their arrow labels.
- If an activity box lies in the top-most level, the box number must be empty.
- If an activity box is not in the top-most level, the box number must be within 1 to 6.

- In the case of CONTROL and OUTPUT, the number of touched arrows must be within 1 to 5.
- In the case of INPUT and MECHANISM, the number of touched arrows must be within 0 to 5.

Figures 5.4 and 5.5 show a more detailed Activity IDEF<sub>0</sub> Syntax according to above English sentences.

Column 1 in Figure 5.4 and 5.5 shows the cases of INPUT, OUTPUT, and NUMBER in the figure 5.2. In the case CONTROL and MECHANISM, if-then rules are similar to OUTPUT and INPUT respectively. Column 2 presents all the diagram types which the user could possibly draw. Column 3 shows if-then rules related to the possible drawings. As shown in Figure 5.4, Activity IDEF<sub>0</sub> Syntax focuses on whether NAME, INPUT, OUTPUT, CONTROL, MECHANISM, and NUMBER are correct or not. Thus the goals for Activity IDEF<sub>0</sub> Syntax rules become a list of triples about NAME, INPUT, OUTPUT, CONTROL, MECHANISM, and NUMBER.

Boundary IDEF<sub>0</sub> Syntax. The Boundary IDEF<sub>0</sub> Syntax is associated with boundary arrows of an IDEF<sub>0</sub> diagram and its parent IDEF<sub>0</sub> diagram. The following english sentences represent the Boundary IDEF<sub>0</sub> Syntax.

- There must be an activity box in the parent IDEF<sub>0</sub> diagram.
- The number of input, output, control, or mechanism arrow(s) of the parent activity box must be equal to that of the boundary input, output, control, or mechanism arrow(s).
- Each arrow of the parent activity box must have its label.
- Each boundary arrow must have its label.
- Each boundary arrow label must correspond with label of the parent activity box arrow.

Case	Diagram	If-then rules	
INPUT (MECHANISM)		If there is no input arrow then INPUT is correct on the activity box.	
	11 b	If there is at least a blank LABEL then there is a LABEL error on the activity box.	
	11, 1m,	If the number of input arrows is greater than 5 then the number of input arrows should be reduced. Otherwise, INPUT is correct.	
OUTPUT (CONTROL)		If there is no output arrow then there is an OUTPUT error on the activity box.	
	11 b	If there is at least a blank LABEL then there is a LABEL error on the activity box.	
	11 <sub>2</sub> 1n <sub>2</sub>	If the number of output arrows is greater than 5 that of output arrows should be reduced. Otherwise, OUTPUT is correct.	
NUMBER		If the activity box is the top-most level and NUMBER is empty then NUMBER of the activity box should be empty.	
	NUMBER	If the activity box is the top-most level and NUMBER is not empty then NUMBER of the activity box should be empty.	
		If the activity box is not the top- most level and NUMBER is greater than 0 and less than 7 then NUMBER of the activity box is correct. Otherwise, NUMBER of the activity box is beyond the limitation.	

Figure 5.4. Activity IDEF<sub>0</sub> Syntax

Case	Diagram	If-then rules	
NAME	NAME	If there is no activity box name then there is a NAME error on the activity box.	
		If there is a NAME on the activity box then NAME is correct.	

Figure 5.5. Activity IDEF<sub>0</sub> Syntax(continued)

- In the case of (boundary) CONTROL and (boundary) OUTPUT, the number of arrows must be greater than or equal to 1 and less than 6.
- In the case of (boundary) INPUT and (boundary) MECHANISM, the number of arrows must be within 0 to 5.

Figure 5.6 represents the more detailed Boundary IDEF<sub>0</sub> Syntax. Column 1 in Figure 5.6 and 5.7 shows the cases of Boundary Input and Boundary Output. Ifthen rules for Boundary Mechanism and Boundary Control are similar to Boundary Input and Boundary Output respectively. Column 2 shows the models of a parent activity box which is possibly drawn focused on INPUT(OUTPUT) arrow(s). Column 3 shows the models of Boundary Input(Output) arrow(s) which is possibly drawn. Column 4 presents if-then rules about the Boundary IDEF<sub>0</sub> Syntax. Boundary IDEF<sub>0</sub> Syntax focuses on whether the boundary arrows in any IDEF<sub>0</sub> diagram correspond with the arrows on the parent activity box. Thus the goals for Boundary IDEF<sub>0</sub> Syntax rules is a list of triples about Boundary Input, Boundary Output, Boundary Control, and Boundary Mechanism.

# Software Test

The software testing methods for IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System were performed using three steps: unit testing, integration test-

CASE	PARENT	CHILD	If-then rules		
PARENT	THERE IS NO PARENT BOX INFORMATION.	DON'T CARE	If there is no information about the parent activity box then can not check.		
BOUNDARY INPUT (MECHANISM)	DON'T CARE		If there is at least one blank boundary arrow LABEL then there is a LABEL error.		
		DON'T CARE	If there is at least one blank LABEL on the parent activity box then parent Input arrow has no LABEL.		
		_li	If the number of Input arrows of the parent box is greater than or less than that of the boundary Input arrows then there is an error of the number of parent Input and boundary Input arrows.		
			If there is no Input arrow of parent activity box and boundary arrow then boundary Input is correct.		
			If LABELs of Input arrows of the parent activity box and boundary Input arrows are all matched then boundary Input is correct.  If there is at least one LABEL which is mismatched then there is an error of		
			mismatched LABEL.		
	DON'T CARE		If the number of boundary Input arrows is greater than 5 then the number of those should be reduced.		

Figure 5.6. Boundary IDEF<sub>0</sub> Syntax

CASE	PARENT	CHILD	If-then rules
BOUNDARY OUTPUT (CONTROL)	DON'T CARE	-l1	If there is at least one blank boundary arrow LABEL then there is a LABEL error.
	11	DON'T CARE	If there is at least one blank LABEL on the parent activity box then parent Output arrow has no LABEL.
	DON'T CARE		If there is no boundary Output arrow then there must be at least a boundary Output arrow.
		DON'T CARE	If there is no Output arrow then there must be at least an Output arrow on parent activity box.
	li l	li i	If the number of Output arrows of the parent box is greater than or less than that of the boundary Output arrows then there is an error of the number of parent Output and boundary Output arrows.
		li i	If LABELs of Input arrows of the parent activity box and boundary Input arrows are all matched then boundary Input is correct.  If there is at least one LABEL which is mismatched then there is an error of
	DON'T CARE	li ln	mismatched LABEL.  If the number of boundary Output arrows is greater than 5 then the number of those should be reduced.

Figure 5.7. Boundary IDEF<sub>0</sub> Syntax(Continued)

ing, and validation testing (12:502). These methods use the white box testing methods discussed in Chapter IV.

The Unit testing step focuses on each module individually to be sure that it functions properly as a unit (12:502). The test considerations are module interface, data structure, boundary conditions, basis path through the control structure, and error handling paths (12:503). Unit test considerations are applied to IDT and ISES individually. For IDEF<sub>0</sub> Diagram Translator, IDEF<sub>0</sub> diagram is prepared. Predicate data files are also prepared for IDEF<sub>0</sub> Syntax Expert System.

The Integration testing step is applied to take unit-tested modules and construct a complete program structure to ensure that the interfaces between modules are correct (12:507). Bottom-up integration is used because IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System are lower-level than SAtool. First, testing for IDEF<sub>0</sub> Diagram Translator was performed and then the whole system of SAtool was examined. Also testing of IDEF<sub>0</sub> Syntax Expert System was applied separated because ISES is separated with IDT and SAtool.

The Validation testing step is performed to provide final assurance that the software meets the mentioned requirements. This focuses on the Are we building the right product? (12:499). This step was used to examine whether the IDT produced predicate data forms correctly and the ISES checked completely IDEF<sub>0</sub> syntax.

# Summary

This chapter described the low level design and implementation of IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System based on the requirements of the tool. The translation rules and the components of ISES are also represented. Finally, testing methodology applied in this investigation are discussed.

# VI. Conclusions and Recommendations

# Introduction

The objective of this thesis investigation was to design and implement an application of expert system for checking IDEF<sub>0</sub> syntax of IDEF<sub>0</sub> diagrams as derived from the SAtool. This chapter presents the conclusions and directions for possible future research.

## Conclusions

This investigation is classified into two major categories: IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System. The work for the IDEF<sub>0</sub> Diagram Translator was performed in two phases. During the first phase, the formulation of graphical features of the IDEF<sub>0</sub> diagram was derived through predicate calculus representation, since the predicate calculus is a convenient representation for facts and rules of inference. The formal definition of the IDEF<sub>0</sub> graphical features does not have completeness but consistency because the IDEF<sub>0</sub> graphical language contains semantic meanings. The second phase included the development of the  ${\rm IDEF_0}$ Diagram Translator which translates the IDEF<sub>0</sub> graphical features in the IDEF<sub>0</sub> diagram into a set of predicate data forms. The predicate data forms focus on an activity box and associated arrows and on the boundary arrows in any IDEF<sub>0</sub> diagram. Predicate data forms become the data base (working memory) for the IDEF<sub>0</sub> Syntax Expert System. The IDEF<sub>0</sub> Syntax Expert System consists of the inference engine, the knowledge base, the data base, and the user interface. The inference engine applies the knowledge to the solution of a specific domain. To check IDEF<sub>0</sub> syntax in any IDEF<sub>0</sub> diagram, the backward chaining control strategy is useful because there are many more facts than final goals. The knowledge base was identified with emphasis on the activity box and on the boundary arrows in any IDEF<sub>0</sub> diagram. The knowledge base structure is easy to extend to new IDEF<sub>0</sub> syntax rules independently of other rules and to change independently of other rules. Each segment of the knowledge base defines a small and relatively independent piece of knowledge.

# Recommendations

Based on the results of this study and the observations made during it, this section presents some recommendations for future research which could lead to enhance the capability of the ISES.

Activity Currently, the tool can check the IDEF<sub>0</sub> syntax of only a single activity box and associated arrows in any IDEF<sub>0</sub> diagrams. The relationship among activity boxes and arrows in composite IDEF<sub>0</sub> diagrams could be defined to enhancements of the ISES's capability. For example,

- The name of an activity box should not be the same as that of other activity boxes in any IDEF<sub>0</sub> diagrams.
- The number of an activity box should not be the same as that of other activity boxes in any IDEF<sub>0</sub> diagram.
- The line label on an activity box should not be the same as that on other activity boxes in any IDEF<sub>0</sub> diagram.

Boundary Since the ISES can check syntax of the boundary arrows except the tunnel arrow, add the IDEF<sub>0</sub> syntax rules about the tunnel arrow. This issue implies IDEF<sub>0</sub> syntax check of the multilevel IDEF<sub>0</sub> diagrams.

Integrate the translation process with the syntax checking process to be more user friendly. This issue needs to address how the C language should interface with Quintus Prolog or some other prolog.

Apply the structure of the knowledge-based IDEF<sub>0</sub> syntax system to incorporate the design knowledge of a specific software application. Since the design knowledge provides a means of abstracting software design into resuable modules,

the design knowledge using the IDEF<sub>0</sub> methodology can be reused for a similar software design. The knowledge-based IDEF<sub>0</sub> syntax system uses IDEF<sub>0</sub> model segments to represent design modules which are combined and refined to generate an entire IDEF<sub>0</sub> model.

# Summary

This chapter presented the conclusions derived from the design and implementation of an application of expert system for checking IDEF<sub>0</sub> syntax of IDEF<sub>0</sub> diagrams drawn from the SAtool and the recommendations for future research.

# Appendix A. Requirements Analysis Diagrams

This appendix contains the requirements analysis IDEF<sub>0</sub> diagrams for the IDEF<sub>0</sub> Diagram Translator. These diagrams are not exactly one-to-one correspondence with the implementation modules, but are close.

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#### A-0 Provide SAtool

Abstract: Provide SAtool provides a means of mechanism by which the user is able to draw Activity IDEFO Diagrams. From these diagrams, Facing Page Text and Data Dictionaries for Activities and Data and Predicate Data forms are generated.

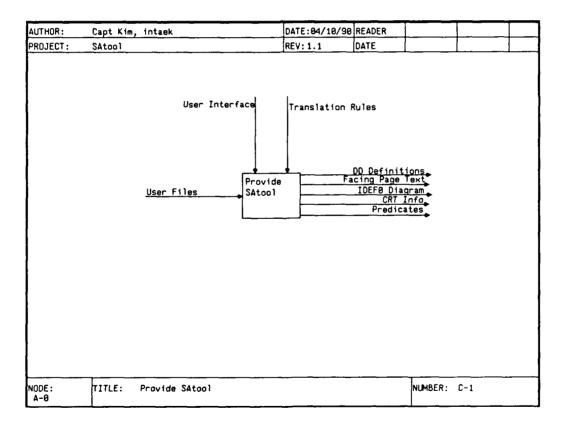


Figure A.1. Provide SAtool (Node A-0)

#### A0 Provide SAtool

Abstract: Provide SAtool provides the user a mechanism by which the user is able to draw Activity IDEFO Diagrams. From these diagrams, Facing Page Text, Data Dictionaries for Activities and Data, and Predicate Data forms are generated.

Al Provide SA Editor provides a means of drawing Activity IDEFO Diagrams. From these diagrams, Facing Page Text and Data Dictionaries for Activities and Data are generated.

A2 Translate Diagram provides a means of translating IDEFO Diagrams into Predicate Data Forms.

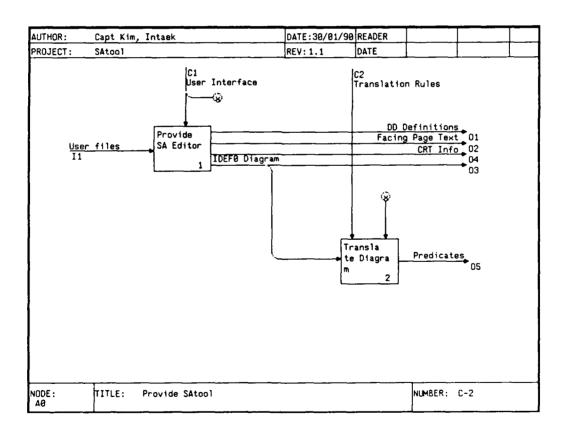


Figure A.2. Provide SAtool (Node A0)

# A2 Translate Diagram

Abstract: Translate Diagram provides the user a means of translating IDEFO Diagrams into Predicate Data Forms.

- A21 Translate Activity provides a means of translating an activity box in any IDEFO Diagrams into a set of predicate data forms through the translation rules of the activity box.
- A22 Translate Boundary provides a means of translating the boundary arrows in any IDEFO Diagrams into Predicate data forms through the translation rules of the boundary arrows.
- A23 Save Diagram provides a means of translating the IDEFO Diagram into a set of the predicate data forms and saving it into a file which user specifies interactively.

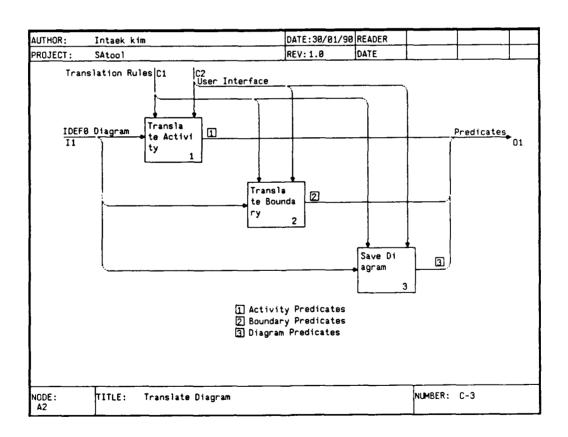


Figure A.3. Translate Diagram (Node A2)

# A21 Translate Activity

Abstract: Translate Activity provides a means of translating an activity box in any IDEFO Diagrams into a set of predicate data forms through the translation rules of the activity box.

- A211 'find clicked box' module provides a means of finding an activity box which user specifies using the mouse in the IDEFO diagram.
- A212 'save header info' module provides a means of saving the head information of the IDEFO diagram which is needed to check Boundary IDEFO Syntax.
- A213 'save box arrow info' module provides a means of grasping the information of the activity box and the arrows which are attached on the activity box.

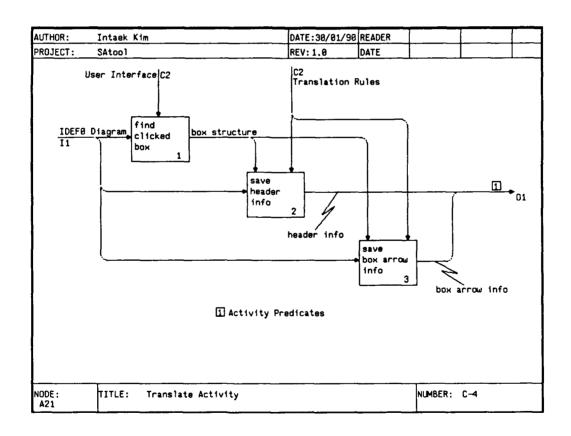


Figure A.4. Translate Activity (Node A21)

#### A213 save box arrow info

Abstract: 'save box arrow info' module provides a means of grasping the information of the activity box and the arrows which are attached on the activity box.

A2131 'get box info' module provides a means of holding and saving the activity box name and the number into a file in forms of predicate using the translation rules.

A2132 'get arrow info' module provides the information of the arrows which are touched an activity box in forms of predicate.

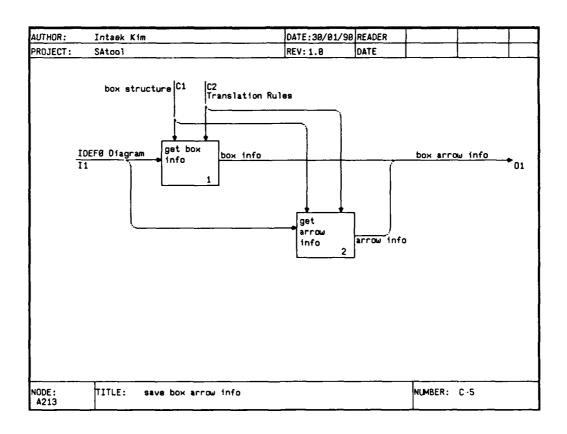


Figure A.5. Save Box Arrow Info (Node A213)

A2132 get arrow info

Abstract: 'get arrow info' module provides the information of the arrows which are touched an activity box in forms of predicate.

- A21321 'get inputs' module provides a means of grasping the predicate data forms of all kind of input arrows which are touched on an activity box.
- A21322 'get outputs' module provides a means of grasping the predicate data forms of all kind of output arrows which are touched on an activity box.
- A21323 'get controls' module provides a means of grasping the predicate data forms of all kind of control arrows which are touched on an activity box.
- A21324 'get mechanisms' module provides a means of grasping the predicate data forms of all kind of mechanism arrows which are touched on an activity box.

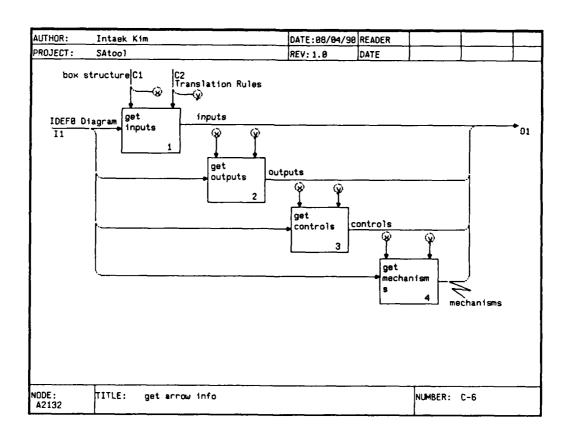


Figure A.6. Get Arrow Info (Node A2132)

# A21321 get inputs

Abstract: 'get inputs' module provides a means of grasping the predicate data forms of all kind of input arrows which are touched on an activity box.

A213211 'get single head input' module provides a means grasping the single headed input arrows' information translated in the predicate data forms which are touched on an activity box.

A213212 'get double head input' module provides a means of grasping the predicate data forms of the double headed input arrows which are touched on an activity box.

A213213 'get doublehead in/w slash' module provides a means of grasping the predicate data forms of the double headed input arrows with the slash which are touched on an activity box.

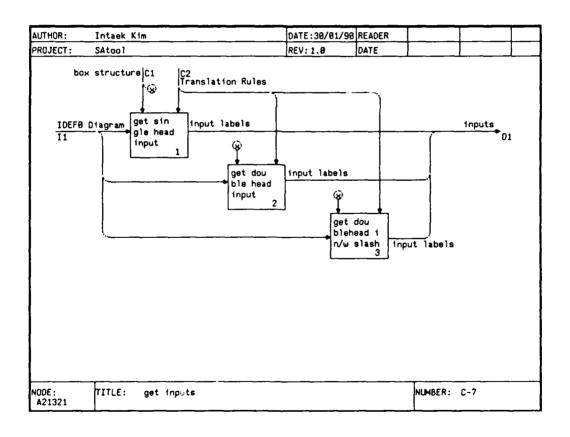


Figure A.7. Get Inputs (Node A21321)

# A21322 get outputs

Abstract: 'get outputs' module provides a means of grasping the predicate data forms of all kind of output arrows which are touched on an activity box.

A213221 'get single head output' module provides a means of grasping the predicate data forms of the output arrows with a single head which are touched on an activity box.

A213222 'get double head out/con' module provides a means of grasping the predicate data forms of the output arrows with the double head one for the output and the other for the control arrow of another box.

A213223 'get double head out/in' module provides a means of grasping the predicate data forms of the output arrows with the double head one for the output and the other for the input arrow of another box.

A213224 'get double head output' module provides a means of grasping the predicate data forms of the output arrow which leaves the right side of an activity box and there is a double head.

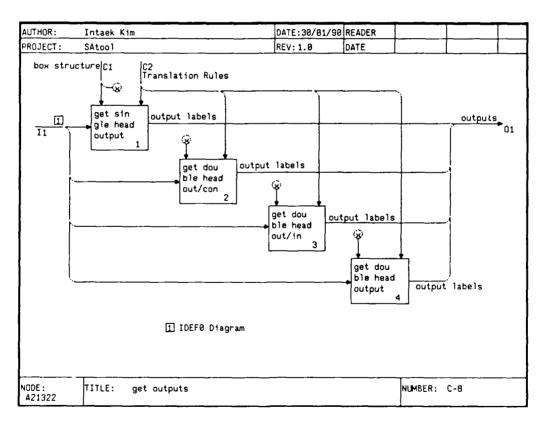


Figure A.8. Get Outputs (Node A21322)

#### A21323 get controls

Abstract: 'get controls' module provides a means of grasping the predicate data forms of all kind of control arrows which are touched on an activity box.

A213231 'get single head control' module provides a means of grasping the predicate data forms of the control line which comes to the upper side of an activity box and there is a single headed arrow.

A213232 'get double head control' module provides a means of grasping the predicate data forms of the control line which comes to the upper side of an activity box and there is a double headed arrow.

A213233 'get double head con/slash' module provides a means of grasping the predicate data forms of the control line which comes to the upper side of an activity box and there is a double headed arrow with a slash.

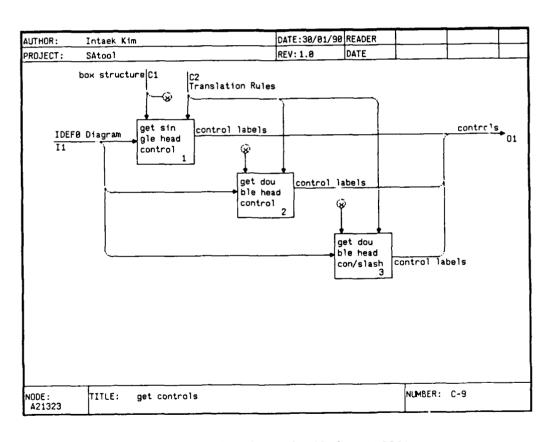


Figure A.9. Get Controls (Node A21323)

#### A22 Translate Boundary

Abstract: Translate Boundary provides a means of translating the boundary arrows in any IDEFO Diagrams into Predicate data forms through the translation rules of the boundary arrows.

- A221 'get parent box' module provides a means of grasping the predicates for the parent activity box and producing the parent information.
- A222 'save null boundary' module provides a means of grasping the predicates if there is no the boundary arrow in according with the input, output, control, or mechanism of the IDEFO diagram.
- A223 'save boundary info' module provides a means of grasping the predicates of the boundary arrows if there is at least one boundary arrow in the IDEFO diagram.

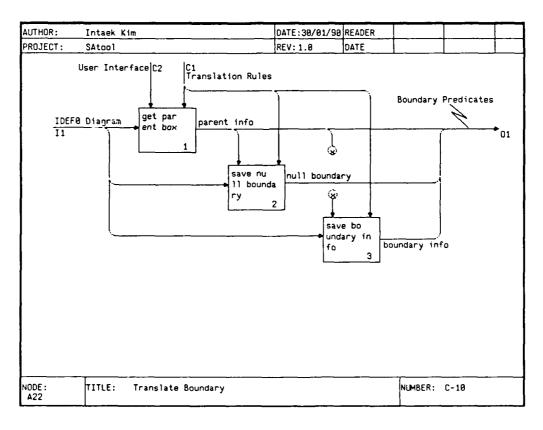


Figure A.10. Translate Boundary (Node A22)

# A223 save boundary info

Abstract: 'save boundary info' module provides a means of grasping the predicates of the boundary arrows if there is at least one boundary arrow in the IDEFO diagram.

A2231 'search boundary lines' module provides a means of searching for the boundary lines for the IDEFO diagram and producing a linked list of line structure as the output.

A2232 'get boundary line labels' module provides a means of grasping the line labels of the boundary lines in the IDEFO diagram.

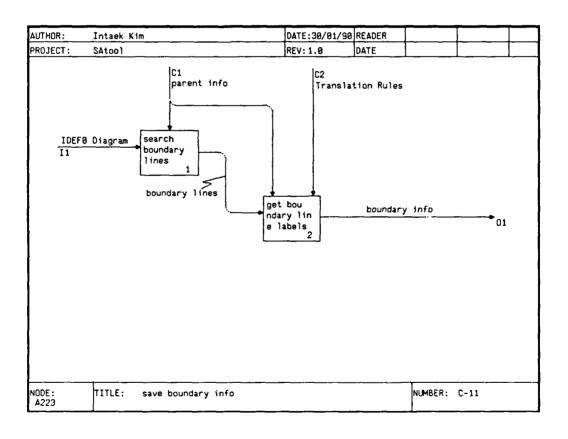


Figure A.11. Save Boundary Info (Node A223)

## A23 Save Diagram

Abstract: Save Diagram provides a means of translating the IDEFO Diagram into a set of the predicate data forms and saving it into a file which user specifies interactively.

A231 The function of 'get file name' module is to get a file name from the user in order to save the predicate data forms for the IDEFO diagram into it.

A232 The function of 'store predicates' module is to save the predicates for the IDEFO diagram into a file which the user specifies.

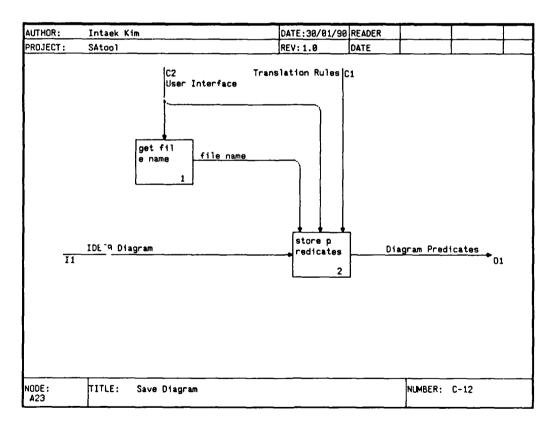


Figure A.12. Save Diagram (Node A23)

## A232 store predicates

Abstract: The function of 'store predicates' module is to save the predicates for the IDEFO diagram into a file which the user specifies.

A2321 'save header info' has the same function of module A212. See A212 description.

A2322 'traverse boxes' module function is to traversing every boxes in the IDEFO diagram.

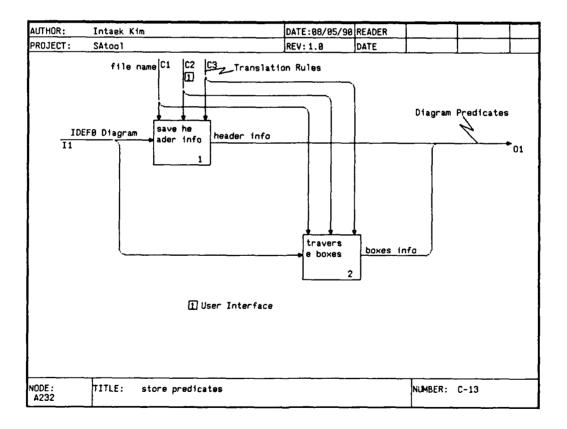


Figure A.13. Store Predicates (Node A232)

## A2322 traverse boxes

Abstract: 'traverse boxes' module function is to traversing every boxes in the IDEFO diagram.

A23221 'get a box' module function is to get the information for an activity box in the IDEFO diagram.

A23222 'save box arrow info' module function is the same as module A213.

A23223 'get boxes arrows' module function is to gether the predicates of every activity box and arrow in the IDEFO diagram.

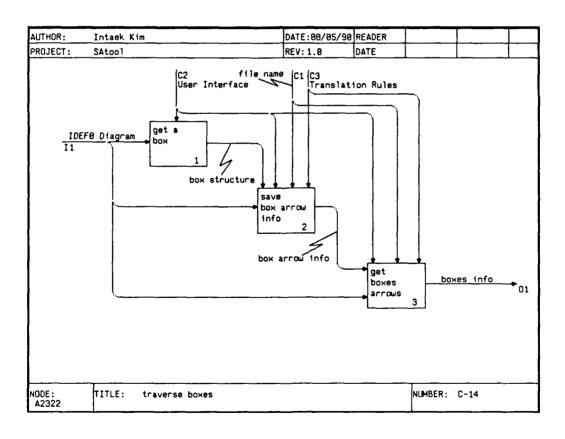


Figure A.14. Traverse Boxes (Node A2322)

# Appendix B. Structured Chart

This appendix contains the detailed design structure charts for the IDEF<sub>0</sub> Diagram Translator implementation. The detailed design is concerned with the requirements analysis diagrams in appendix A. There is a close, but not exactly one-to-one, correspondence between the design modules and the implementation modules.

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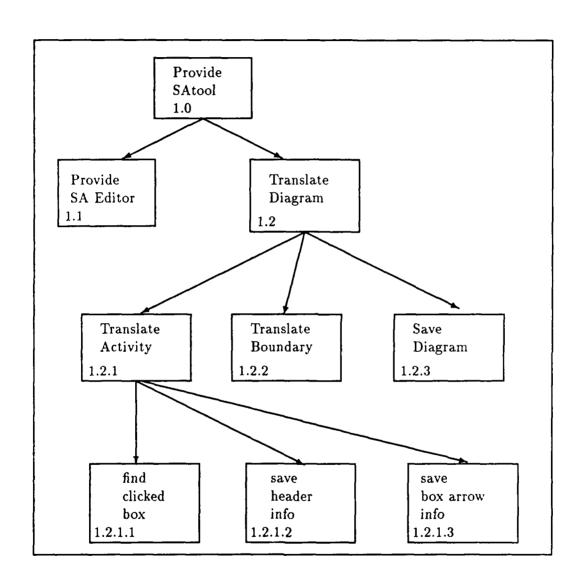


Figure B.1. Provide SAtool(Module 1.0)

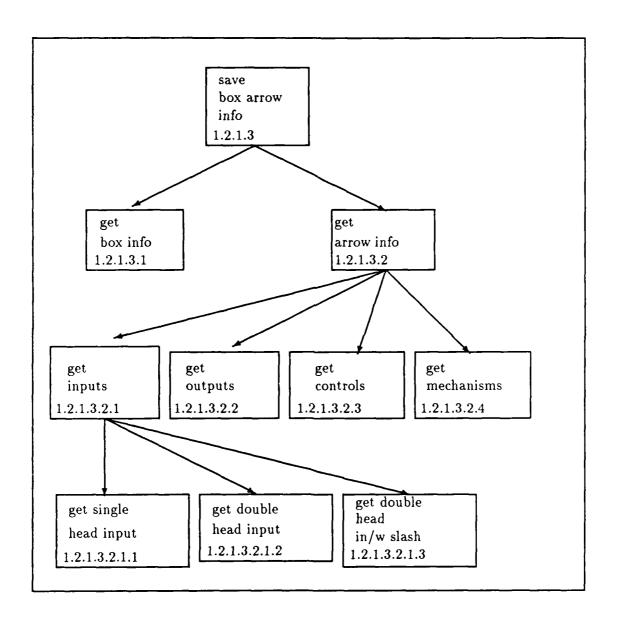


Figure B.2. Save Box Arrow Info(Module 1.2.1.3)

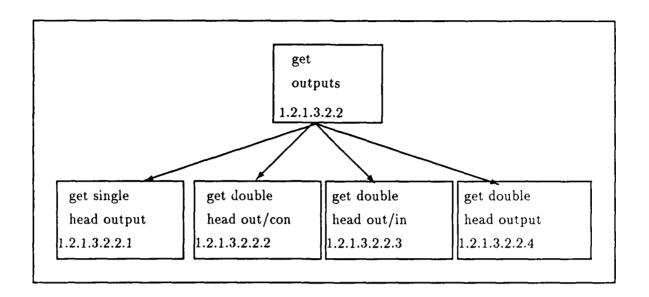


Figure B.3. Get Outputs(Module 1.2.1.3.2.2)

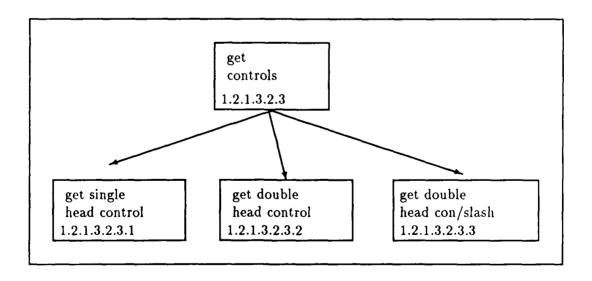


Figure B.4. Get Controls(Module 1.2.1.3.2.3)

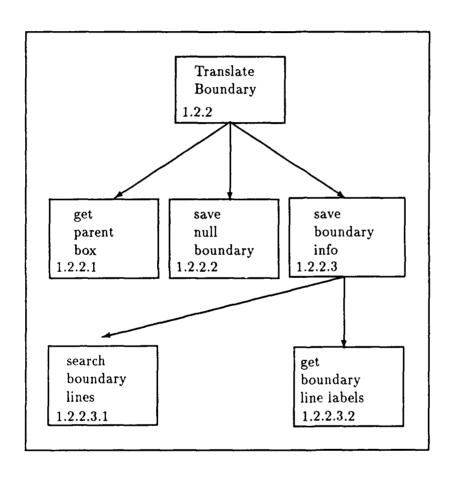


Figure B.5. Translate Boundary(Module 1.2.2)

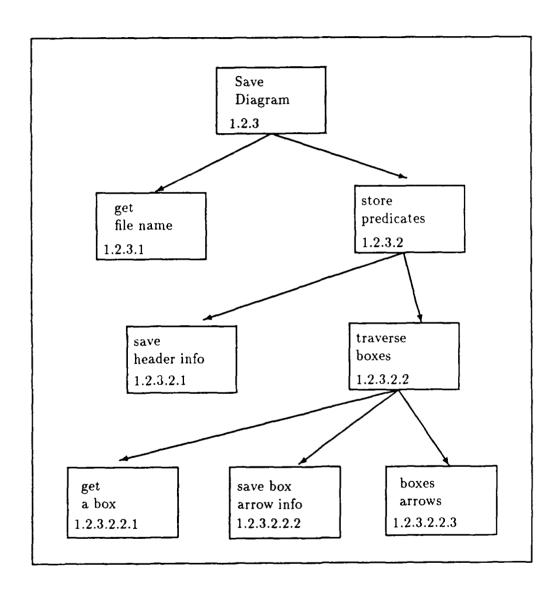


Figure B.6. Save Diagram(Module 1.2.3)

# Appendix C. Data Structures of SAtool

### Introduction

The purpose of this appendix is to discuss of the Data Structures of SAtool developed by Steven E. Johnson (9). A discussion of the data structures of the SAtool is needed because this thesis work should interface with the SAtool and use the IDEF<sub>0</sub> diagrams and files generated by the SAtool. The SAtool allows users to interactively create and edit IDEF<sub>0</sub> diagrams and to automatically produce the data dictionary information.

### Data Structure

Five primary data structures were designed to hold all the graphics and data dictionary information: the box structure, the line structure, the squiggle line structure, the header structure, and the footnote structure (9:4-11 - 4-14).

The box structure contains the information which is necessary to locate, name, and enumerate activity boxes (9:4-11). The activity boxes in the IDEF<sub>0</sub> diagram are hooked by the linked list. The box structure uses two C pointers one for the next box structure and the other for pointing an activity data dictionary structure (9:4-11).

The line structure consists of the fields which are necessary to lacate, label, draw the lines, enumerate the lines to identify them, store the ICOM labels, and store the TO/FROM ALL labels (9:4-11 - 4-12). The line structure uses two numbers to identify the type of each end of the line (ie. arrowhead, tunnel, dot, turn right, or branch left, etc.) and uses C pointers to store the lines in binary trees with the root nodes (9:4-12). Figure C.1 shows four groups of the lines and the corresponding linked list structure is presented in Figure C.2 (9:4-12 - 4:13). The tree arrangement in Figure A.2 maps to how the line segments actually connect to one another and C pointer supports the simple recursive functions used to traverse the binary trees

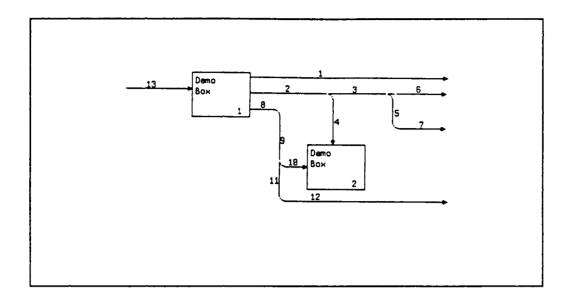


Figure C.1. Example Group of lines (9:4-12)

(9:4-12). The line structure uses another C pointer to point to a data divtionary information for a data element.

The squiggle line structure contains the information which is necessary to locate and to identify each squiggle line in the IDEF<sub>0</sub> diagram (9:4-13). The squiggle line structure uses a C pointer to store the squiggle lines for a particular IDEF<sub>0</sub> diagram in a single linked list (9:4-13 - 4-14).

The header structure consists of seven fields which are needed to draw, locate. and classify AUTHOR, DATE, PROJECT, REV, NODE, TITLE, and NUMBER of an IDEF<sub>0</sub> diagram(9:4-14 - 4-15). A single C pointer is used to save the header information since each IDEF<sub>0</sub> diagram only has one header (9:4-14).

Finally, the footnote structure contains the information which is needed to draw, locate and identify a matched pair of footnote labels (9:4-14). A C pointer is defined to point another footnote structure since the footnote structures for a IDEF<sub>0</sub> diagram are stored in a single linked list (9:4-14).

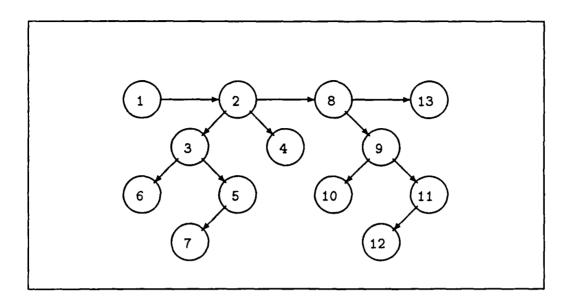


Figure C.2. The Linked List for Lines (9:4-13)

# Summary

In this appendix, the data structure of the SAtool which is necessary to perform this thesis investigation was addressed from Johnson's effort. This information was used throughout this thesis effort.

# Appendix D. User's Manual

User's Manual introduces the basics of the ISES. The purpose of this manual provides a broad understanding of the ISES's operation, then provides a more detailed example for learning to use the ISES.

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### Introduction

The IDEF<sub>0</sub> Syntax Expert System(ISES) is an interactive syntax check system. It provides a means for checking IDEF<sub>0</sub> syntax in any IDEF<sub>0</sub> diagrams drawn by SAtool and, producing error messages, error recovery, and editing suggestions. ISES allows the user to select a menu for drawing an IDEF<sub>0</sub> diagram and checking IDEF<sub>0</sub> syntax. The functions of the main menus include:

- RECALL DGM read in a previously saved IDEF<sub>0</sub> diagram.
- EDIT DGM edit an IDEF<sub>0</sub> diagram according to its attribute menus.
- EDIT DD edit a data dictionary.
- EDIT FPT edit a facing page text.
- MISC FUNC save a diagram, change directory, exit SAtool, etc.
- SAVE DGM save a graphics information and a data dictionary information.
- CHECK SYNTAX check IDEF<sub>0</sub> syntax.

The first six menus are for drawing IDEF<sub>0</sub> diagrams, generating Data Dictionary information, and Facing Page Text and the last one is for ISES to check IDEF<sub>0</sub> syntax. A detailed guide for drawing IDEF<sub>0</sub> diagram is available in the user's manual of Johnson's thesis (9). This User's Manual focuses on the CHECK SYNTAX part. ISES runs on Sun3<sup>TM</sup> and Sun4<sup>TM</sup> workstations using the SunOS<sup>TM 1</sup> and the SunView<sup>TM</sup> window-based environment. This manual explains how CHECK SYNTAX can be used to check IDEF<sub>0</sub> syntax. Some previous familiarity with IDEF<sub>0</sub> syntax and SAtool is required. Though not necessary, some familiarity with SunOS and SunView is helpful. Users should be thoroughly familiar with the concepts presented in this manual before using ISES.

<sup>&</sup>lt;sup>1</sup>SunOS<sup>TM</sup> is a trademark of Sun Microsystems, Inc.

#### The Mouse

To move the cursor, place the mouse on its pad and move it in the desired direction. During the execution of SAtool, User is required to click the mouse button. All mouse button inputs should be clicked on the proper location in IDEF<sub>0</sub> diagram, otherwise, the mouse button inputs are ignored.

## • Right Button

The right button is used almost to abort operation of menu selected.

#### • Middle Button

The middle button is used to exit SAtool (see Exit SAtool).

## • Left Button

The left mouse button is used to select one of menus and to start a menu operation.

#### How to draw lines well

Almost of the unnoticed errors are produced in the field of drawing lines. They provide a menas of generating unsuitable predicate data forms.

## 1. Boundary lines

- All boundary lines should have their ICOM labels.
- All arrows of Inputs, Mechanisms, and Controls must be touched on any box. The segment of lines inside an activity box is truncated automatically.
- All output lines should be begun inside an activity box.

## 2. Inter activity box lines

Every starting and ending point of the line segments should be begun and ended inside an activity box excepting the branch, join lines, and TO/FROM lines.

## Getting Started

Set your UNIX path variable to include the ISES executable directory.

- Enter "suntools"
   enter SunView and SunWindow environment.
- Enter "SAtool"
   enter the IDEF<sub>0</sub> Diagram Translator environment.

### 3. The Main Menu

Menus are displayed as the oval forms on the screen. Move the cursor to one the following choices to select:

- RECALL DGM
- EDIT DGM
- EDIT DD
- EDIT FPT
- MISC FUNC
- SAVE DGM
- CHECK SYNTAX

## 4. IDEF<sub>0</sub> diagram

By selecting one of the first five menus, The user is able to draw a new  $IDEF_0$  diagram or update the previous  $IDEF_0$  diagram.

## 5. IDEF<sub>0</sub> Syntax

After drawing an IDEF<sub>0</sub> diagram, select CHECK SYNTAX oval by clicking the left mouse button on it. Now, three attribute submenus are displayed as the rectangular forms. Move the cursor to on, of the following choices to select:

• Activity

- Boundary
- Save(.pro)

IDEF<sub>0</sub> syntax consists of Activity and Boundary IDEF<sub>0</sub> syntax.

## 6. Activity IDEF<sub>0</sub> syntax

After clicking the left button on Activity rectangular box of the submenus,

- (a) Move the cursor inside a box to be checked and click the left mouse button (Right ABORT).
- (b) Enter the Prolog environment using another window.

## 7. Boundary IDEF<sub>0</sub> Syntax

NOTE: User must have the predicate file of the parent IDEF<sub>0</sub> diagram.

- (a) After clicking the left button on Boundary rectangular box of menus,
- (b) move the cursor inside the input window and enter the file name with the parent activity box information (Right-ABORT).
- (c) Enter the Prolog environment using another window.

## 8. Saving the predicate file

- (a) After clicking the left mouse button on "save(.pro)" rectangular box of menus,
- (b) move the cursor inside the input window, and then enter the file name for the current IDEF<sub>0</sub> diagram. This file is a set of predicate data forms translated from the graphical information in the IDEF<sub>0</sub> diagram. It is used to check Boundary IDEF<sub>0</sub> syntax. The extension of the predicate file is a .pro.

## 9. Prolog Environment

Enter "prolog". This invokes the Prolog interpreter.

(a) Enter "['ISES']." - consult ISES.

Now, the following message is showed:

```
*/
             WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
/*
                                                              */
/*
                                                              */
/* I.
       Type start.
                        to begin a new session.
                                                              */
/* II. Answer all questions using lower case and ending with*/
/*
        a period.
                                                              */
/*
                                                              */
/* III. Type
                 halt. to exit Prolog session.
                                                              */
                                                              */
```

(b) Enter "start." to begin checking IDEF<sub>0</sub> syntax of a clicked box in the current IDEF<sub>0</sub> diagram. Then, the message:

Question: Do you want verbose operation (y./n.)? is displayed. Enter "y." or "n.". In sace of "y.", the list of rules fired are shown and in case of "n.", only the IDEF<sub>0</sub> syntax messages are displayed. (See Examples).

(c) After then, the following message is shown on the screen:

Question: Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS or to have HELP MESSAGES ?

```
To check ACTIVITY BOX -> Enter a.
To check BOUNDARY ARROWS -> Enter b.
To have HELP MESSAGES -> Enter h.
Choice:
If checking Activity IDEF$_{0}$ syntax, enter "a.", checking Boundary IDEF$_{0}$ syntax, enter "b.", or wishing Help Message, enter "h.".
```

## 10. IDEF<sub>0</sub> Syntax Message

According to selection above description, the resulting messages of the IDEF<sub>0</sub> syntax checking procedure are displayed (see Examples).

## 11. Trace

The message, Question: Do you wish to see how this answer was arrived at (y./n.)? is followed by the resulting messages. Enter "y." or "n.". "y." means that the trace regarding the IDEF<sub>0</sub> syntax message derived is displayed and "n." skips (see Examples).

## 12. Save Working Memory

Then, the message, Question: Do you wish to save the current working memory in a file (y./n.)? is displayed. Enter "y." or "n.". In the case of "y.", the current working memory is saved in a file which is specified by the user and of "n.", the current working memory is erased automatically.

## 13. Exit Prolog Environment

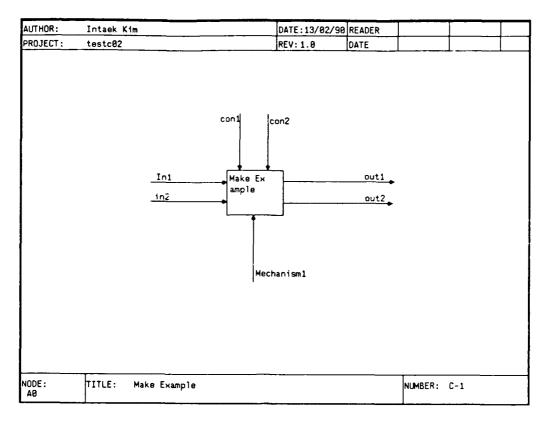
By entering "halt." or "ctrl c", prolog session is ended.

## 14. Exit SAtool

To exit SAtool, click the left mouse button on the "MISC FUNC" oval of the main menus. And then click the left mouse button on the "QUIT" under the "MISC FUNC" oval. Finally, click the middle button of the mouse.

## Examples

This section presents the actual demonstrations for checking process of the correct IDEF<sub>0</sub> diagram, however, the checking process about the IDEF<sub>0</sub> diagram with errors is the same as the correct case.



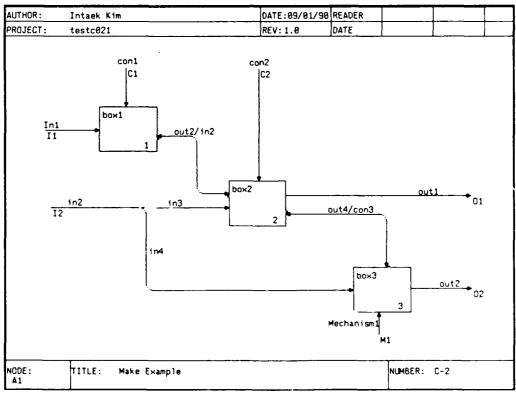


Figure D.1. Test  $\mathrm{IDFF}_0$  Diagram

```
Quintus Prolog Release 2.4.2 (Sun-4, SunOS 4.0)
Copyright (C) 1988, Quintus Computer Systems, Inc. All rights reserved.
1310 Villa Street, Mountain View, California (415) 965-7700
| ?- ['ISES'].
[consulting /usr2/eng/ikim/SAtoolExpert/ISES...]
  /*
                                                       */
  /*
            WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
                                                       */
  /*
                                                       */
  /* I. Type start. to begin a new session.
                                                       */
  /* II. Answer all questions using lower case, ending with
                                                       */
  /*
        a period.
                                                       */
                                                       */
  /* III. Type halt. to exit prolog session.
                                                       */
  /*
                                                       */
  [ISES consulted 1.367 sec 19,008 bytes]
yes
| ?- start.
Question: Do you want verbose operation(y./n.)? n.
Question: Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS
         or to have HELP MESSAGES ?
       To check ACTIVITY BOX -> Enter a.
       To check BOUNDARY ARROWS -> Enter b.
       To have HELP MESSAGE -> Enter h.
Choice : a.
/******** IDEFO Syntax Messages ***********/
          --> CORRECT: Activity Name is OK.
Name
Input
          --> CORRECT: Input is OK.
          --> CORRECT: Output is OK.
Output
```

ares> prolog

```
Control --> CORRECT: Control is OK.

Mechanism --> CORRECT: Mechanism is OK.

Number --> CORRECT: Activity number is OK.
```

This activity must be the top most level box.

Question: Do you wish to see how this answer was arrived at(y./n.)? n.

Question: Do you wish to save the current working memory in a file(y./n.)? n.

```
*/
          WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
                                                      */
/*
                                                     */
/* I. Type start. to begin a new session.
                                                     */
                                                      */
/* II. Answer all questions using lower case, ending with
                                                     */
/*
                                                      */
      a period.
/*
                                                     */
/* III. Type halt. to exit prolog session.
                                                      */
                                                      */
ves
! ?- start.
Question: Do you want verbose operation(y./n.)? y.
 Question: Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS
         or to have HELP MESSAGES ?
       To check ACTIVITY BOX -> Enter a.
       To check BOUNDARY ARROWS -> Enter b.
       To have HELP MESSAGE -> Enter h.
Choice : a.
 /************* IDEFO Syntax Messages **********/
                        --> ERROR: No Activity Name.
Trying rule1:: [Name, ,
                  Each box must have an activity name]
Trying rule2:: [Name, , --> CORRECT: Activity Name is OK]
Proved rule2:: [Name, , --> CORRECT: Activity Name is OK]
          --> CORRECT: Activity Name is OK.
                        --> CORRECT: No Input Arrows, however,
Trying rule3:: [Input, ,
                      Input is OK]
Trying rule4:: [Input, ,
                       --> ERROR: No Input Label
                Each Input arrow must have an input label]
Trying rule5:: [Input, , --> RECOMMEND:
                You would better reduce the number of Input arrows
                from 0 to 5]
Trying rule6:: [Input, , --> CORRECT: Input is OK]
Proved rule6:: [Input, , --> CORRECT: Input is OK]
```

Trying rule7:: [Output, , --> ERROR: You should have at least one output arrow] Trying rule8:: [Output, , --> ERROR: No Output Label. Each Output Arrow should have an output Label] --> RECOMMEND: Trying rule9:: [Output, , You would better reduce the number of Output arrows from 1 to 5] Trying rule10:: [Output, , --> CORRECT: Output is OK] Proved rule10:: [Output, , --> CORRECT: Output is OK] Output --> CORRECT: Output is OK. Trying rule11:: [Control, , --> ERROR: You should have at least one control arrow] Trying rule12:: [Control, , --> ERROR: No Control Label. Each Control Arrow should have a control Label] Trying rule13:: [Control, , --> RECOMMEND: You would better reduce the number of Control arrows from 1 to 5] Trying rule14:: [Control, , --> CORRECT: Control is OK] Proved rule14:: [Control, , --> CORRECT: Control is OK] Control --> CORRECT: Control is OK. Trying rule15:: [Mechanism, ,--> ERROR: No Mechanism Label. Each Mechanism Arrow should have a mechanism Label] Trying rule16:: [Mechanism, ,--> CORRECT: No Mechanism Arrows, however, Mechanism is OK] Trying rule17:: [Mechanism, ,--> RECOMMEND: You would better reduce the number of Mechanism arrows from 0 to 5] Trying rule18:: [Mechanism, ,--> CORRECT: Mechanism is OK] Proved rule18:: [Mechanism, ,--> CORRECT: Mechanism is OK] --> CORRECT: Mechanism is OK. Mechanism Trying rule19:: [Number, , --> CORRECT: Activity number is OK. This activity must be the top most level box] Proved rule19:: [Number, , --> CORRECT: Activity number is OK. This activity must be the top most level box] Number --> CORRECT: Activity number is OK.

--> CORRECT: Input is OK.

Input

This activity must be the top most level box.

```
Question: Do you wish to see how this answer
          was arrived at(y./n.)? y.
                   --> CORRECT: Activity Name is OK]
GOAL::
        [Name. .
GOAL::
        [Input.,
                    --> CORRECT: Input is OK]
GOAL:: [Output, , --> CORRECT: Output is OK]
GOAL:: [Control, , --> CORRECT: Control is OK]
GOAL:: [Mechanism, ,--> CORRECT: Mechanism is OK]
        [Number, , --> CORRECT: Activity number is OK.
GOAL::
                    This activity must be the top most level box]
rule19:: [Number. , --> CORRECT: Activity number is OK.
         This activity must be the top most level box] Was Derived From
       [title, is, Make Example] AND
       [Make Example, number_is, 0] AND
       [0, ==, 0]
SOLVED:: [0.==.0]
TOLD:: [Make Example, number_is, 0]
       [title.is.Make Example]
TOLD::
rule18:: [Mechanism, ,--> CORRECT: Mechanism is OK] Was Derived From
       [activityname, is, Make Example]
KNOWN:: was_told:: [activityname,is,Make Example]
rule14:: [Control, , --> CORRECT: Control is OK] Was Derived From
       [activityname, is, Make Example]
KNOWN:: was_told:: [activityname,is,Make Example]
rule10:: [Output, , --> CORRECT: Output is OK] Was Derived From
       [activityname, is, Make Example]
KNOWN:: was_told:: [activityname,is,Make Example]
rule6:: [Input, . --> CORRECT: Input is OK] Was Derived From
       [activityname, is, Make Example]
KNOWN:: was_told:: [activityname,is,Make Example]
                   --> CORRECT: Activity Name is OK] Was Derived From
rule2:: [Name, .
       [activityname, is, Make Example] AND
       [Make Example, \==,]
SOLVED:: [Make Example, \==,]
TOLD:: [activityname, is, Make Example]
    /* After this session, all working memory elements will */
    /* be erased except for elements being protected by
                                                        */
         keep statements in the knowledge base.
                                                        */
```

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Question: Do you wish to save the current working memory

in a file(y./n.)? y.

Please supply a filename: 'example.wm'.

```
/*
/*
         WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
                                                  */
/*
                                                  */
/* I. Type start. to begin a new session.
                                                  */
/* II. Answer all questions using lower case, ending with
                                                  */
/*
     a period.
/*
/* III. Type halt. to exit prolog session.
                                                  */
/*
                                                  */
ves
! ?- start.
Question: Do you want verbose operation(y./n.)? n.
Question: Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS
        or to have HELP MESSAGES ?
       To check ACTIVITY BOX -> Enter a.
       To check BOUNDARY ARROWS -> Enter b.
       To have HELP MESSAGE -> Enter h.
Choice : b.
 /************ IDEFO Syntax Messages **********/
Boundary Input
               --> CORRECT: Boundary Input is OK.
Boundary Output --> CORRECT:
                   Boundary Output is OK.
Boundary Control --> CORRECT: Boundary
                    Control is OK.
Boundary Mechanism --> CORRECT: Boundary
                    Mechanism is OK.
```

Question: Do you wish to see how this answer

## was arrived at(y./n.)? n.

Question: Do you wish to save the current working memory in a file(y./n.)? n.

```
*/
/*
          WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
                                                       */
                                                       */
/*
/* I. Type start. to begin a new session.
                                                       */
/*
                                                       */
/* II. Answer all questions using lower case, ending with
                                                       */
      a period.
                                                       */
                                                       */
/* III. Type halt. to exit prolog session.
                                                       */
yes
! ?- start.
Question: Do you want verbose operation(y./n.)? y.
 Question: Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS
         or to have HELP MESSAGES ?
        To check ACTIVITY BOX -> Enter a.
        To check BOUNDARY ARROWS -> Enter b.
        To have HELP MESSAGE -> Enter h.
Choice : b.
 /************ IDEFO Syntax Messages ***********/
Trying rule2:: [Boundary Input, , --> !!! THIS IS A FATAL ERROR !!!]
Trying rule1:: [boundarysarule,is,stalled]
Trying rule6:: [Boundary Input, , --> ERROR: No boundary input label]
Trying rule7:: [Boundary Input, ,--> ERROR: Parent Input has no label]
Trying rule8:: [Boundary Input, ,--> ERROR: The number of Input
        arrow(s) of Parent Activity box is greater than that of
        Boundary Input arrow(s) -- Must have the same number]
Trying rule9:: [Boundary Input, , --> ERROR: The number of Input
        arrow(s) of Parent Activity box is less than that of
        Boundary Input arrow(s) -- Must have the same number]
Trying rule10:: [Boundary Input, ,
                                 --> RECOMMEND:
        You would better reduce the number of arrows to six
        below
Trying rule11:: [Boundary Input, , --> CORRECT: Boundary Input is OK]
Trying rule13:: [Boundary Input, , --> CORRECT: Boundary Input is OK]
```

```
Trying rule12:: [case_of_boundary_in,is,1]
Trying rule14:: [Boundary Input, . --> CORRECT: Boundary Input is OK]
Trying rule12:: [case_of_boundary_in,is,2]
Proved rule12:: [case_of_boundary_in,is,2]
Proved rule14:: [Boundary Input, . --> CORRECT: Boundary Input is OK]
                    --> CORRECT: Boundary Input is OK.
Boundary Input
Trying rule3:: [Boundary Output, .
There is nothing about Parent activity]
Trying rule1:: [boundarysarule.is.stalled]
Trying rule19:: [Boundary Output, ,--> ERROR: No boundary output label]
Trying rule20:: [Boundary Output, , --> ERROR:
Parent Output has no label]
Trying rule21:: [Boundary Output, ,--> ERROR: No boundary output arrow.
                        Should have at least one boundary output arrow]
Trying rule22:: [Boundary Output, , --> ERROR: No parent output arrow.
                       Should have at least one parent output arrow]
Trying rule23:: [Boundary Output, ,
                                     --> ERROR: The number of Output
        arrow(s) of Parent Activity box is greater than that of
        Boundary Output arrow(s) -- Must have the same number]
Trying rule24:: [Boundary Output, , --> ERROR: The number of Output
        arrow(s) of Parent Activity box is less than that of
        Boundary Output arrow(s) -- Must have the same number]
Trying rule25:: [Boundary Output, ,
                                      --> RECOMMEND:
        You would better reduce the number of arrows to six
Trying rule27:: [Boundary Output, , --> CORRECT:
                        Boundary Output is OK]
Trying rule26:: [case_of_boundary_out,is,1]
Trying rule28:: [Boundary Output, , --> CORRECT:
                        Boundary Output is OK]
Trying rule26:: [case_of_boundary_out,is,2]
Proved rule26:: [case_of_boundary_out,is,2]
Proved rule28:: [Boundary Output, , --> CORRECT:
                        Boundary Output is OK]
Boundary Output
                    --> CORRECT:
                        Boundary Output is OK.
Trying rule4:: [Boundary Control, , -->
Maybe you have tried to check syntax with
               a file without PARENT ACTIVITY BOX information]
Trying rule1:: [boundarysarule, is, stalled]
Trying rule33:: [Boundary Control, ,--> ERROR: No boundary control label]
Trying rule34:: [Boundary Control, ,--> ERROR: Parent Control has no label]
Trying rule35:: [Boundary Control, , --> ERROR: No boundary control arrow.
```

```
Should have at least one boundary control arrow]
Trying rule36:: [Boundary Control, , --> ERROR: No parent control arrow.
                        Should have at least one parent control arrow]
Trying rule37:: [Boundary Control, , --> RECOMMEND:
         You would better reduce the number of arrows to six
         belowl
Trying rule39:: [Boundary Control, , --> CORRECT: Boundary
                         Control is OK]
Trying rule38:: [case_of_boundary_con,is,1]
Trying rule40:: [Boundary Control, , --> CORRECT: Boundary
                         Control is OK]
Trying rule38:: [case_of_boundary_con,is,2]
Proved rule38:: [case_of_boundary_con,is,2]
Proved rule40:: [Boundary Control, , --> CORRECT: Boundary
                         Control is OK]
                     --> CORRECT: Boundary Control is OK.
Boundary Control
Trying rule5:: [Boundary Mechanism, ,--> PLEASE START AGAIN !!!]
Trying rule1:: [boundarysarule, is, stalled]
Trying rule45:: [Boundary Mechanism, ,--> ERROR:
 No boundary mechanism label]
Trying rule46:: [Boundary Mechanism, ,--> ERROR:
 Parent Mechanism has no label]
Trying rule47:: [Boundary Mechanism, ,--> ERROR: The number of
         Mechanism arrow(s) of Parent Activity box is greater than that
         of Boundary Mechanism arrow(s) -- Must have the same
         number]
Trying rule48:: [Boundary Mechanism, ,--> ERROR: The number of
         Mechanism arrow(s) of Parent Activity box is less than
         that of Boundary Mechanism arrow(s) -- Must have the same
         number]
Trying rule49:: [Boundary Mechanism, ,--> RECOMMEND:
         You would better reduce the number of arrows to six
         below]
Trying rule50:: [Boundary Mechanism, ,--> CORRECT: Boundary
                         Mechanism is OK]
Trying rule52:: [Boundary Mechanism, ,--> CORRECT: Boundary
                         Mechanism is OK]
Trying rule51:: [case_of_boundary_mech,is,1]
Proved rule51:: [case_of_boundary_mech,is,1]
Proved rule52:: [Boundary Mechanism, ,--> CORRECT: Boundary
                         Mechanism is OK]
                    --> CORRECT: Boundary
Boundary Mechanism
```

Mechanism is OK.

Question: Do you wish to see how this answer was arrived at (y./n.)? y. GOAL:: [Boundary Input. . --> CORRECT: Boundary Input is OK] GOAL:: [Boundary Output, , --> CORRECT: Boundary Output is OK] [Boundary Control, , --> CORRECT: Boundary GOAL:: Control is OK] GOAL:: [Boundary Mechanism, ,--> CORRECT: Boundary Mechanism is OK] rule52:: [Boundary Mechanism, ,--> CORRECT: Boundary Mechanism is OK] Was Derived From [case\_of\_boundary\_mech,is,1] AND [child\_title,is,Make Example] AND [Make Example,mechanism\_is,Mechanism1] AND [boundary\_mechanism1, is, Mechanism1] TOLD:: [boundary\_mechanism1,is,Mechanism1] TOLD:: [Make Example, mechanism\_is, Mechanism1] KNOWN:: was\_told:: [child\_title,is,Make Example] rule51:: [case\_of\_boundary\_mech,is,1] Was Derived From [boundary\_mechanism, has\_number, 1] AND [child\_title,is,Make Example] AND [Make Example, has\_mechanism\_number, 1] TOLD:: [Make Example, has\_mechanism\_number.1] KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_mechanism, has\_number, 1] rule40:: [Boundary Control, , --> CORRECT: Boundary Control is OK] Was Derived From [case\_of\_boundary\_con,is,2] AND [boundary\_control1, is, con1] AND [boundary\_control2.is.con2] AND [child\_title, is, Make Example] AND [Make Example,control\_is,con1] AND [Make Example,control\_is,con2] TOLD:: [Make Example,control\_is,con2] TOLD:: [Make Example,control\_is,con1] KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_control2, is, con2] TOLD:: [boundary\_control1, is, con1] rule38:: [case\_of\_boundary\_con,is,2] Was Derived From [boundary\_control, has\_number, 2] AND [child\_title,is,Make Example] AND

[Make Example, has\_control\_number, 2]

TOLD:: [Make Example, has\_control\_number, 2] KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_control, has\_number, 2] rule28:: [Boundary Output, , --> CORRECT: Boundary Output is OK] Was Derived From [case\_of\_boundary\_out,is,2] AND [boundary\_output1, is.out1] AND [boundary\_output2,is,out2] AND [child\_title, is, Make Example] AND [Make Example,output\_is,out1] AND [Make Example,output\_is,out2] TOLD:: [Make Example, output\_is, out2] TOLD:: [Make Example,output\_is,out1] KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_output2, is, out2] TOLD:: [boundary\_output1,is,out1] rule26:: [case\_of\_boundary\_out,is,2] Was Derived From [boundary\_output,has\_number,2] AND [child\_title, is, Make Example] AND [Make Example, has\_output\_number, 2] TOLD:: [Make Example, has\_output\_number, 2] KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_output, has\_number, 2] --> CORRECT: Boundary Input is OK] rule14:: [Boundary Input, , Was Derived From [case\_of\_boundary\_in,is,2] AND [boundary\_input1, is, in2] AND [boundary\_input2, is, In1] AND [child\_title, is, Make Example] AND [Make Example, input\_is, in2] AND [Make Example, input\_is, In1] TOLD:: [Make Example, input\_is, In1] [Make Example, input\_is, in2] TOLD:: KNOWN:: was\_told:: [child\_title,is,Make Example] TOLD:: [boundary\_input2, is, In1] TOLD:: [boundary\_input1, is, in2] rule12:: [case\_of\_boundary\_in,is,2] Was Derived From [boundary\_input, has\_number, 2] AND [child\_title.is, Make Example] AND [Make Example, has\_input\_number, 2] TOLD:: [Make Example, has\_input\_number, 2] TOLD:: [child\_title, is, Make Example] TOLD:: [boundary\_input, has\_number, 2]

Question: Do you wish to save the current working memory in a file(y./n.)? n.

```
/*
                                    */
    WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
/*
                                    */
                                    */
/*
/* I. Type start. to begin a new session.
                                    */
                                    */
/* II. Answer all questions using lower case, ending with
                                    */
                                    */
/* a period.
                                    */
/* III. Type halt. to exit prolog session.
                                    */
                                    */
```

yes | ?- halt. ares>

### Appendix E. Programmer's Guide

 $\label{lem:programmers} Programmer's \ Guide \ introduces \ several \ topics \ of \ interest \ to \ ISES \ programmers \ and \ developers.$ 

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#### Introduction

The focus of this thesis effort was to design and implement an application of expert system formulation for checking IDEF<sub>0</sub> syntax of IDEF<sub>0</sub> diagrams as derived from SAtool. The work in this thesis is divided into two major ctegories: IDEF<sub>0</sub> Diagram Translator and IDEF<sub>0</sub> Syntax Expert System. The IDEF<sub>0</sub> Diagram Translator translates the IDEF<sub>0</sub> diagrams into a set of predicate forms and the predicate forms file is used as the data base of the IDEF<sub>0</sub> Syntax Expert System. The IDEF<sub>0</sub> Syntax Expert System checks the IDEF<sub>0</sub> syntax of IDEF<sub>0</sub> diagrams. The objective of this appendix is to specify the procedure for generating the executable files and to outline some basic concepts of the translator and expert system.

#### Software Documentation

The existing source code is fully documented in AFIT System Development Documentation Guidelines and Standards (8). The following list shows the file header of the source codes.

- DATE: Date of current version number.
- VERSION: Current version number.
- TITLE: Title for this file.
- FILENAME: File name for the module.
- DESCRIPTION: Description of the module's function.
- AUTHOR: Name of one who responsible for this file.
- PROJECT: Name of the software project of which this file is a part.
- OPERATING SYSTEM: Name and version number of operating system under which this file was written.
- LANGUAGE: Name of language used for source code.
- FILE PROCESSING: How the file is used.

- CONTENTS: Modules contained in the file.
- HISTORY: List of major changes to the file.

The following list presents the subroutine header of the source codes.

- DATE: Date of the module.
- VERSION: Current version number.
- NAME: Module name.
- MODULE NUMBER: Module number of current program.
- DESCRIPTION: Text description of the module's function.
- ALGORITHM: Algorithm used.
- PASSED VARIABLES: Variables passed to the module.
- RETURNS: Value returned by the module.
- GLOBAL VARIABLES USED: Variables read by the module.
- GLOBAL VARIABLES CHANGED: Variables changed by the module.
- FILES READ: Files read by the module.
- FILES WRITTEN: Files written by the module.
- HARDWARE INPUT: I/O ports read.
- HARDWARE OUTPUT: I/O ports read.
- MODULES CALLED: Other procedures called.
- CALLING MODULES: What modules call.
- AUTHOR: One who wrote this module.
- HISTORY: List of major changes to the module.

```
OBJECTS = main.o dataddict.o messages.o boxfunctions.o
headerfunctions.o editboxfunc.o
miscfunctions.o ddsearchfuncs.o
endfuncs.o find.o morelinefuncs.o
linelabel.o moreddfuncs.o ddsearchfuncs.o
savefuncs.o
fptfuncs.o sqglefuncs.o fnotefuncs.o
moresave.o screendump.o readfuncs.o
session.o syntaxfuncs.o

HEADERS = globals.h
ALL = sad
CFLAGS = -O
LIBS = -lsuntool -lsunwindow -lpixrect -lm
sad : $(OBJECTS)
cc $(CFLAGS) $(OBJECTS) $(LIBS) -o SAtool
```

Figure E.1. Contents of Makefile

#### Make File

The file of the IDEF<sub>0</sub> Diagram Translator is included in the files of the SAtool because the IDEF<sub>0</sub> Diagram Translator was coded as a part of the SAtool under the SunOS<sup>TM</sup>. The file name of source code for the IDEF<sub>0</sub> Diagram Translator is syntaxfuncs.c. The executable file was produced by using the UNIX make utility. Figure E.1 shows the contents of the Makefile file. The system command make causes to be compiled and linked all together and generated the executable file, SAtool.

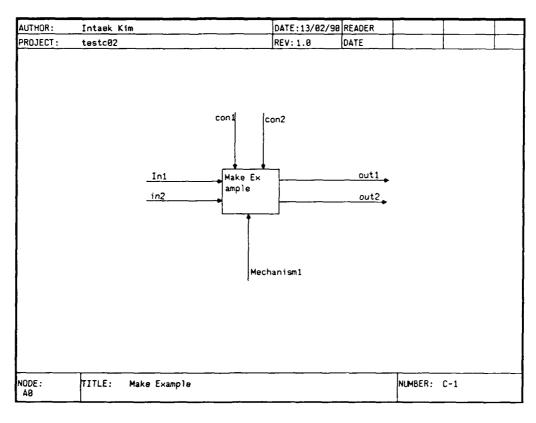
#### Files produced by IDT

\*.pro This file contains a set of predicate data forms into which the IDEF<sub>0</sub> Diagram Translator translates the IDEF<sub>0</sub> diagram. The symbol \* is a file name which the user specifies. The extension of the file is added automatically. This file is used to check the IDEF<sub>0</sub> syntax of boundary arrows in any IDEF<sub>0</sub> diagram. Figure E.3 shows an example of the predicate data file translated from the above IDEF<sub>0</sub>

diagram shown in Figure E.2. Also, Figure E.4 shows an example of the predicate data file from the below IDEF<sub>0</sub> diagram shown in Digure E.2.

CHECKBOX.PRO This file is a temporary file which is created and overwrited automatically. Also this file is used for checking the IDEF<sub>0</sub> syntax of an activity box which user specifies in any IDEF<sub>0</sub> diagram. Figure E.5 represents the predicate data file of the activity box, box1, which is clicked by the user in Figure E.2.

CHECKBOUNDARY.PRO This file contains the predicate data forms of the boundary arrows and its parent activity box and is a temporary file which is created and overwrited automatically. Also this file becomes the data base (working memory) of the IDEF<sub>0</sub> Syntax Expert System. Figure E.6 shows the predicate data file of boundary arrows in Figure E.2.



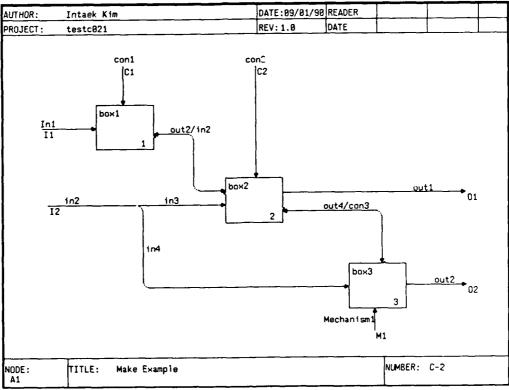


Figure E.2. Example of IDEF<sub>0</sub> Diagram

```
confirmed([title, is, 'Make Example']).
confirmed([node, is, 'A0']).
confirmed([activityname, is, 'Make Example']).
confirmed(['Make Example',number_is,0]).
confirmed(['Make Example',input_is,'In1']).
confirmed(['Make Example',input_is,'in2']).
confirmed(['Make Example', has_input_number, 2]).
confirmed(['Make Example',output_is,'out1']).
confirmed(['Make Example',output_is,'out2']).
confirmed(['Make Example', has_output_number, 2]).
confirmed(['Make Example',control_is,'con1']).
confirmed(['Make Example',control_is,'con2']).
confirmed(['Make Example', has_control_number, 2]).
confirmed(['Make Example', mechanism_is,'Mechanism1']).
confirmed(['Make Example', has_mechanism_number, 1]).
```

Figure E.3. Predicate Data File Produced by Save (.pro) (parent)

```
confirmed([title, is, 'Make Example']).
confirmed([node, is, 'A1']).
confirmed([activityname, is, 'box1']).
confirmed(['box1',number_is,1]).
confirmed(['box1',input_is,'in2']).
confirmed(['box1',input_is,'In1']).
confirmed(['box1', has_input_number, 2]).
confirmed(['box1',output_is,'out2']).
confirmed(['box1', has_output_number, 1]).
confirmed(['box1',control_is,'con1']).
confirmed(['box1', has_control_number, 1]).
confirmed(['box1',mechanism_is,null]).
confirmed([activityname, is, 'box2']).
confirmed(['box2',number_is,2]).
confirmed(['box2',input_is,'out2']).
confirmed(['box2',input_is,'in3']).
confirmed(['box2', has_input_number, 2]).
confirmed(['box2',output_is,'in2']).
confirmed(['box2',output_is,'out1']).
confirmed(['box2',output_is,'out4']).
confirmed(['box2', has_output_number, 3]).
confirmed(['box2',control_is,'con2']).
confirmed(['box2',control_is,'con3']).
confirmed(['box2', has_control_number, 2]).
confirmed(['box2',mechanism_is,null]).
confirmed([activityname, is, 'box3']).
confirmed(['box3',number_is,3]).
confirmed(['box3',input_is,'in4']).
confirmed(['box3', has_input_number, 1]).
confirmed(['box3',output_is,'out2']).
confirmed(['box3',output_is,'con3']).
confirmed(['box3', has_output_number, 2]).
confirmed(['box3',control_is,'out4']).
confirmed(['box3', has_control_number, 1]).
confirmed(['box3',mechanism_is,'Mechanism1']).
confirmed(['box3', has_mechanism_number, 1]).
```

Figure E.4. Predicate Data File Produced by Save (.pro) (child)

```
confirmed([title, is, 'Make Example']).
confirmed([node, is, 'A1']).
confirmed([activityname, is, 'box1']).
confirmed(['box1',number_is,1]).
confirmed(['box1',input_is,'in2']).
confirmed(['box1',input_is,'In1']).
confirmed(['box1', has_input_number, 2]).
confirmed(['box1',output_is,'out2']).
confirmed(['box1',output_number, 1]).
confirmed(['box1',control_is,'con1']).
confirmed(['box1',has_control_number, 1]).
confirmed(['box1',mechanism_is,null]).
```

Figure E.5. Predicate Data File Produced by Activity

```
confirmed([title, is, 'Make Example']).
confirmed([node, is, 'A0']).
confirmed([activityname, is, 'Make Example']).
confirmed(['Make Example',number_is,0]).
confirmed(['Make Example',input_is,'In1']).
confirmed(['Make Example',input_is,'in2']).
confirmed(['Make Example', has_input_number, 2]).
confirmed(['Make Example',output_is,'out1']).
confirmed(['Make Example',output_is,'out2']).
confirmed(['Make Example', has_output_number, 2]).
confirmed(['Make Example',control_is,'con1']).
confirmed(['Make Example',control_is,'con2']).
confirmed(['Make Example', has_control_number, 2]).
confirmed(['Make Example',mechanism_is,'Mechanism1']).
confirmed(['Make Example', has_mechanism_number, 1]).
confirmed([child_title, is, 'Make Example']).
confirmed([boundary_control1, is, 'con1']).
confirmed([boundary_output1, is, 'out1']).
confirmed([boundary_output2, is, 'out2']).
confirmed([boundary_input1, is, 'in2']).
confirmed([boundary_control2, is, 'con2']).
confirmed([boundary_input2, is, 'In1']).
confirmed([boundary_mechanism1, is, 'Mechanism1']).
confirmed([boundary_input, has_number, 2]).
confirmed([boundary_output, has_number, 2]).
confirmed([boundary_control, has_number, 2]).
confirmed([boundary_mechanism, has_number, 1]).
```

Figure E.6. Predicate Data File Produced by Boundary

### Appendix F. Source Code

The purpose of this appendix represents the source code which was implemented during this thesis investigation. This appendix contains two main source codes: one for IDEF<sub>0</sub> Diagram Translator, the other for IDEF<sub>0</sub> Syntax Expert System.

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### IDEF<sub>0</sub> Diagram Translator

The purpose of this section is to provide the source code documentation of the IDEF<sub>0</sub> Diagram Translator. The documentation conformed to the software engineering standards in AFIT's System Development Documentation Guidelines and Standards draft #4 (8).

```
***********
    DATE:
             3 Feb 1990
    VERSION: 1.0
    TITLE:
             IDEFO Diagram Translator
    FILENAME: syntaxfuncs.c
    DESCRIPTION:
         This file provides a means of translating the IDEFO diagram
          into a set of predicate forms and generating the predicates
         files as the data base of IDEFO Syntax Expert System.
    PROJECT: AI
    OPERATING SYSTEM: UNIX 4.3
    LANGUAGE:
    FILE PROCESSING: Must compile with SAtool.
    CONTENTS: check_input_abox(), single_headed_input(),
              double_headed_input(), double_headed_input_with_slash(), *
              check_output_abox(), single_headed_output(),
              double_headed_output_with_control(),
              double_headed_output_with_input(), double_headed_output()*
              check_control_abox(), single_headed_control(),
              double_headed_control_with_slash(),
              double_headed_control(), check_mechanism_abox(),
              single_headed_mechanism(),search_labels_touched_abox(),
              get_labels_for_abox(), save_arrow_info_of_abox(),
              create_temp_box_info(), find_clicked_box(),
              check_button_for_activity(), check_activity(),
              create_temp_boundary_info(), check_parent_box_file(),
              check_button_for_boundary(), check_boundary(),
              save_null_boundary(), search_NLR_boundary_line_info(),
              search_boundary_info(), save_boundary_line_info(),
              save_header_info(), traverse_boxes(), store_predicates(),*
              overwrite_predicates(), get_filename_for_predicates()
              save_predicates()
    AUTHOR: Intack Kim
    HISTORY:
              10/Jan/90 : Modify the print out format
             02/Feb/90 : Add save_header_info()
             04/Feb/90 : Add save_boundary_info()
#include <stdio.h>
#include <suntool/sunview.h>
#include <suntool/canvas.h>
#include <suntool/panel.h>
```

```
**********************
    DATE: 15 Feb 1990
    VERSION: 1.0
 * NAME: check_input_abox()
    MODULE NUMBER:
    DESCRIPTION:
              This module provides a means of generating a linked
              list, Line_labels, with all information labels of input *
              arrows attatched on an activity box.
    ALGORITHM:
    PASSED VARIABLES: tem_line
    RETURNS: None
    GLOBAL VARIABLES USED: Line_labels
    GLOVAL VARIABLES CHANGED: Line_labels
    FILES READ: None
    FILES WRITTEN: None
    HARDWARE INPUT: None
    HARDWARE OUTPUT: None
    MODULES CALLED: signle_headed_input(), double_headed_input()
                   double_headed_input_with_slash()
    CALLING MODULES: search_labels_touched_abox()
    AUTHOR: Intack Kim
    HISTORY:
    ABSTRACT DATA TYPE :
    ORDER OF :
 **************
void
check_input_abox(tem_line)
struct line_struct *tem_line;
extern struct text_line_struct *Line_labels;
extern add_to_inputs_tree();
char buf[DESCRIPTION_LINE_LENGTH+1];
   single_headed_input(tem_line);
   double_headed_input(tem_line);
   double_headed_input_with_slash(tem_line);
   return;
   }
```

```
/********************************
 * DATE: 25 Feb 1990
 * VERSION: 1.0
 * NAME: single_headed_input()
 * MODULE NUMBER:
 * DESCRIPTION: This module provides a means of hooking the line
               labels of input arrows with single head to a linked
               list, Line_label.
 * ALGROTHM:
 * PASSED VARIABLES: tem_line
 * RETURNS: None
 * GLOBAL VARIABLES USED: box,line,Line_labels
 * GLOBAL VARIABLES CHANGED: Line_labels
 * FILES READ: None
 * FILES WRITTEN: None
 * HARDWARE INPUT: None
 * HARDWARE CUTPUT: None
 * MODULES CALLED: get_closest_label(), get_branchnode(),
                  find_TO_ALL_branchnode(), find_branchnode(),
                  add_to_inputs_tree()
 * CALLING MODULES: check_input_abox()
 * AUTHOR: Intack Kim
 * HISTORY:
 * ABSTRACT DATA TYPE:
 * ORDER OF:
 single_headed_input(tem_line)
 struct line_struct *tem_line;
  {
 extern get_closest_label(),get_branchnode(),find_TO_ALL_branchnode();
 extern find_branchnode(),add_to_inputs_tree();
 extern struct box_struct *box;
 extern struct line_struct *line;
struct line_struct *original_line;
       buf[DESCRIPTION_LINE_LENGTH +1];
char
int
       original_line_type;
   if( /* check if tem_line is touched on the left side of */
/* box with 10 deviation.
       (tem_line->end_position.x >= box->swcorner.x-10) &&
       (tem_line->end_position.x <= box->swcorner.x)
       (tem_line->end_position.y <= box->swcorner.y)
       (tem_line->end_position.y >= box->swcorner.y-BOX_HT) &&
      ( /* check if tem_line has an end arrow */
```

```
(tem_line->end_activity_num & (ARROW_HEAD)) != 0) &&
       ((tem_line->end_activity_num & (DOT_B_L|DOT_T_R)) == 0)) {
    /* This tem_line is an input arrow for temp_box with some
                                                                     */
    /* toralence, i.e. not require the line must be touched box
                                                                     */
    /* so as to be an input arrow(single headed arrow).
                                                                     */
      original_line_type = get_branchnode(tem_line);
      original_line = (struct line_struct *)
      find_branchnode(original_line_type);
      if(get_closest_label(original_line,buf,tem_line->struct_type)
  == MYTRUE)
{ /* tem_line has a line label */
Line_labels = (struct text_line_struct *)
add_to_inputs_tree(Line_labels,buf);
return:
}
    /* No line label on tem_line - Check if tem_line is a FROM_ALL */
    /* type line. If so, search TO_ALL and its label
                                                                    */
      if(original_line->start_activity_num == FROM_ALL) {
      original_line = (struct line_struct *)
      find_TO_ALL_branchnode(original_line->to_from_label);
      original_line_type = get_branchnode(original_line);
      original_line = (struct line_struct *)
      find_branchnode(original_line_type);
       if(get_closest_label(original_line,buf,line->struct_type)
 == MYTRUE)
 {
 Line_labels = (struct text_line_struct *)
         add_to_inputs_tree(Line_labels,buf);
 return;
 }
        }
    strcpy(buf,"");
    Line_labels = (struct text_line_struct *)
    add_to_inputs_tree(Line_labels,buf);
   }
  return;
 }
```

```
* DATE: 16 Feb 1990
* VERSION: 1.0
* NAME : double_headed_input()
* MODULE NUMBER:
* DESCRIPTION: This module provides a linked list(Line_labels) which *
               has the labels of the input arrows touched on an
               activity box with double head.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: box, line
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get2_closest_label(), find_branchnode(),
                  find_TO_ALL_branchnode(), get_branchnode(),
                  add_to_inputs_tree()
* CALLING MODULES: check_input_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
double_headed_input(tem_line)
struct line_struct *tem_line;
 {
extern get2_closest_label(),find_branchnode(),find_TO_ALL_branchnode();
extern get_branchnode(),add_to_inputs_tree();
extern struct box_struct *box;
extern struct line_struct *line;
struct line_struct *original_line;
     buf[DESCRIPTION_LINE_LENGTH +1];
char
int
      original_line_type;
   if((tem_line->end_position.x >= box->swcorner.x-10) &&
      (tem_line->end_position.x <= box->swcorner.x)
      (tem_line->end_position.y <= box->swcorner.y)
      (tem_line->end_position.y >= box->swcorner.y-BOX_HT) &&
      ( /* tem_line is an input arrow with a double headed arrow */
                                                                  */
        /* in some extend(toralence).
       (tem_line->end_activity_num & (ARROW_HEAD)) != 0) &&
```

```
((tem_line->end_activity_num & (DOT_T_R)) != 0)) {
       original_line_type = get_branchnode(tem_line);
       original_line = (struct line_struct *)
 find_branchnode(original_line_type);
       if(/* Exists the line label on tem_line */
  get2_closest_label(original_line,buf,tem_line->struct_type)
          == MYTRUE) {
          Line_labels = (struct text_line_struct *)
  add_to_inputs_tree(Line_labels,buf);
          return;
          }
      /* No line label on tem_line - Check if tem_line is a FROM_ALL */
      /* type line. If so, search TO_ALL and its label.
      if(original_line->start_activity_num == FROM_ALL) {
       original_line = (struct line_struct *)
find_TO_ALL_branchnode(original_line->to_from_label);
       original_line_type = get_branchnode(original_line);
       original_line = (struct line_struct *)
find_branchnode(original_line_type);
       if(get2_closest_label(original_line,buf,line->struct_type)
           == MYTRUE) {
          Line_labels = (struct text_line_struct *)
  add_to_inputs_tree(Line_labels,buf);
  return;
  }
         }
      strcpy(buf,"");
      Line_labels = (struct text_line_struct *)
     add_to_inputs_tree(Line_labels,buf);
     return;
     }
```

```
* DATE: 16 Feb 1990
 * VERSION: 1.0
 * NAME : double_headed_input_with_slash()
 * MODULE NUMBER:
 * DESCRIPTION: This module provides a means of adding the line
                labels with a slash of input arrows with double head *
                to Line_labels structure
 * ALGORITHM:
 * PASSED VARIABLES: tem_line
 * RETURNS: None
 * GLOBAL VARIABLES USED: box
 * GLOBAL VARIABLES CHANGED: Line_labels
 * FILES READ: None
 * FILES WRITTEN: None
 * HARDWARE INPUT: None
 * HARDWARE OUTPUT: None
 * MODULES CALLED: get3_first_label(), add_to_inputs_tree()
 * CALLING MODULES: check_input_abox()
 * AUTHOR: Intack Kim
 * HISTORY:
 * ABSTRACT DATA TYPE:
 * ORDER OF:
double_headed_input_with_slash(tem_line)
struct line_struct *tem_line;
extern get3_first_label(),add_to_inputs_tree();
extern struct box_struct *box;
char buf[DESCRIPTION_LINE_LENGTH +1];
   if((tem_line->start_position.x < box->swcorner.x+BOX_WIDTH+15) &&
       (tem_line->start_position.x >= box->swcorner.x+BOX_WIDTH) &&
       (tem_line->start_position.y <= box->swcorner.y)
                                                                 22
       ((tem_line->start_activity_num & (ARROW_HEAD)) !=0)
       ((tem_line->start_activity_num & (DOT_B_L)) != 0))
       {
if(get3_first_label(tem_time,buf) == MYTRUE) {
 Line_labels = (struct text_line_struct *)
 add_to_inputs_tree(Line_labels,buf);
 return;
 }
        strcpy(buf,"");
Line_labels = (struct text_line_struct *)
```

```
add_to_inputs_tree(Line_labels,buf);
}
return;
}
```

```
/*********************************
* DATE: 17 Feb 1990
* VERSION: 1.0
* NAME: check_output_abox()
* MODULE NUMBER:
 * DESCRIPTION:
              This module provides a means of generating a linked
              list, Line_labels, with all information labels of
              output arrows attatched on an activity box.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
 * HARDWARE OUTPUT: None
* MODULES CALLED: single_headed_output(), double_headed_output()
                 double_headed_output_with_input()
                 double_headed_output_with_control()
* CALLING MODULES: search_labels_touched_abox()
* AUTHOR: Intack Kim
* HISTORY:
 * ABSTRACT DATA TYPE:
* ORDER OF:
*****************
check_output_abox(tem_line)
struct line_struct *tem_line;
   extern struct text_line_struct *Line_labels;
    single_headed_output(tem_line);
    double_headed_output_with_control(tem_line);
    double_headed_output_with_input(tem_line);
    double_headed_output(tem_line);
    return;
  }
```

```
* DATE: 17 Feb 1990
 * VERSION: 1.0
 * NAME: single_headed_output()
 * MODULE NUMBER:
 * DESCRIPTION:
              This module provides a means of adding line labels of *
              output arrows with single head.
 * ALGORITHM:
 * PASSED VARIABLES: tem_line
 * RETURNS: None
 * GLOBAL VARIABLES USED: Line_labels
 * GLOBAL VARIABLES CHANGED: Line_labels
 * FILES READ: None
 * FILES WRITTEN: None
 * HARDWARE INPUT: None
 * HARDWARE OUTPUT: None
 * MODULES CALLED: get_first_label(), add_to_inputs_tree()
 * CALLING MODULES: check_output_abox()
 * AUTHOR: Intack Kim
 * HISTORY:
 * ABSTRACT DATA TYPE:
 * ORDER OF:
 *********************
 single_headed_output(tem_line)
 struct line_struct *tem_line;
    extern get_first_label(),add_to_inputs_tree();
   char buf[DESCRIPTION_LINE_LENGTH +1];
   if(/* Output leaves the right side of a box - single headed arrow */
       (tem_line->start_position.x < box->swcorner.x+BOX_WIDTH+15) &&
       (tem_line->start_position.x >= box->swcorner.x+BOX_WIDTH)
       (tem_line->start_position.y <= box->swcorner.y)
                                                                  22
       (tem_line->start_position.y >= box->swcorner.y-BOX_HT)
                                                                  & &
       ((tem_line->start_activity_num & (ARROW_HEAD)) == 0))
       if(get_first_label(tem_line,buf) == MYTRUE) {
         Line_labels = (struct text_line_struct *)
 add_to_inputs_tree(Line_labels,buf);
         return:
 }
       strcpy(buf,"");
Line_labels = (struct text_line_struct *)
```

```
add_to_inputs_tree(Line_labels,buf);
}
return;
}
```

```
* DATE: 17 Feb 1990
* VERSION: 1.0
* NAME: double_headed_output_with_control()
* MODULE NUMBER:
* DESCRIPTION:
              This module provides a means of adding the line labels*
              of the output arrows with double heads which become
              controls of another activity box to a linked list.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: line, Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get_branchnode(), find_branchnode(),
                  get3_closest_label(), find_TO_ALL_branchnode()
                  add_to_inputs_tree()
* CALLING MODULES: check_output_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
double_headed_output_with_control(tem_line)
struct line_struct *tem_line;
 -{
  extern get_branchnode(),find_branchnode(),get3_closest_label();
  extern find_TO_ALL_branchnode();
  extern struct line_struct *line;
  struct line_struct *original_line;
  char buf[DESCRIPTION_LINE_LENGTH+1];
  int original_line_type;
  if(/* Output with a double headed control arrow - need to
      /* search the tem_line for the closest label with a slash */
      /* and get the label after the slash.
                                                                 */
      (tem_line->end_position.y > box->swcorner.y-BOX_HT-15)
                                                                 社社
      (tem_line->end_position.y <= box->swcorner.y-BOX_HT)
                                                                 金金
      (tem_line->end_position.x <= box->swcorner.x+BOX_WIDTH)
                                                                 数数
      (tem_line->end_position.x >= box->swcorner.x)
```

```
22
       ((tem_line->end_activity_num & (ARROW_HEAD)) != 0)
       ((tem_line->end_activity_num & (DOT_B_L)) != 0))
original_line_type = get_branchnode(tem_line);
       original_line = (struct line_struct *)
find_branchnode(original_line_type);
       if(get3_closest_label(original_line,buf,tem_line->struct_type)
 == MYTRUE) {
 Line_labels = (struct text_line_struct *)
 add_to_inputs_tree(Line_labels,buf);
 return:
 }
        /* No line label on tem_line - Check if tem_line is a
                                                                    */
/* FROM_ALL type output. If so, search TO_ALL and its label.*/
if(original_line->start_activity_num == FROM_ALL) {
  original_line = (struct line_struct *)
 find_TO_ALL_branchnode(original_line->to_from_label);
          original_line_type = get_branchnode(original_line);
 original_line = (struct line_struct *)
 find_branchnode(original_line_type);
  if(/*have found the branchnode of the TO_ALL line segment*/
    get3_closest_label(original_line,buf,line->struct_type))
     Line_labels = (struct text_line_struct *)
     add_to_inputs_tree(Line_labels,buf);
     return:
     }
         strcpy(buf,"");
Line_labels = (struct text_line_struct *)
add_to_inputs_tree(Line_labels,buf);
      return;
```

```
* DATE: 17 Feb 1990
* VERSION: 1.0
* NAME : double_headed_output_with_input()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of adding the line
               labels of the output with double heads and the slash*
               to a linked list, Line_labels.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: line
* GLOBAL VARIABLES CHANGED: Line_labels
* FILE READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get_branchnode(), find_branchnode(),
                 get3_closest_label(), add_to_inputs_tree()
                 find_TO_ALL_branchnode()
* CALLING MODULES: check_output_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
double_headed_output_with_input(tem_line)
struct line_struct *tem_line;
 {
  extern get_branchnode(),find_branchnode(),get3_closest_label();
  extern add_to_inputs_tree(),find_TO_ALL_branchnode();
  extern struct line_struct *line;
  struct line_struct *original_line;
  int original_line_type;
  char buf[DESCRIPTION_LINE_LENGTH +1];
  if(/* tem_line is an output with a double headed input */
     /* arrow on a box */
     (tem_line->end_position.x >= box->swcorner.x -15)
                                                        金金
     (tem_line->end_position.x <= box->swcorner.x)
                                                        22
     (tem_line->end_position.y <= box->swcorner.y)
     (tem_line->end_position.y >= box->swcorner.y -BOX_HT)&&
     ((tem_line->end_activity_num & (ARROW_HEAD)) != 0)
     ((tem_line->end_activity_num & (DOT_T_R)) != 0))
```

```
original_line_type = get_branchnode(tem_line);
      original_line = (struct line_struct *)
                find_branchnode(original_line_type);
       if(/* Exists label on the original line */
  get3_closest_label(original_line,buf,tem_line->struct_type)
  == MYTRUE) {
 Line_labels = (struct text_line_struct *)
 add_to_inputs_tree(Line_labels,buf);
 return;
 }
       if(/* No line label on tem_line - Check if
                                                       */
 /* tem_line is a type of FROM_ALL
  /* line. If so, Search TO_ALL and its label. */
 original_line->start_activity_num == FROM_ALL)
  original_line = (struct line_struct *)
find_TO_ALL_branchnode(original_line->to_from_label);
           original_line_type = get_branchnode(original_line);
  original_line = (struct line_struct *)
           find_branchnode(original_line_type);
   if(/* Find the branchnode of the TO_ALL line segment */
      get3_closest_label(original_line,buf,line->struct_type)
     ≈= MYTRUE)
     {
              Line_labels = (struct text_line_struct *)
      add_to_inputs_tree(Line_labels,buf);
              return;
      }
          }
          strcpy(buf,"");
 Line_labels = (struct text_line_struct *)
add_to_inputs_tree(Line_labels,buf);
        }
       return;
      }
```

```
* DATE: 17 Feb 1990
* VERSION: 1.0
* NAME: double_headed_output()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of adding the line
               labels of the output with double heads to a linked
               list, Line_labels.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HAREWARE OUTPUT: None
* MODULES CALLED: get2_first_label(), add_to_inputs_tree()
* CALLING MODULES: check_output_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
double_headed_output(tem_line)
struct line_struct *tem_line;
  extern get2_first_label(),add_to_inputs_tree();
  char buf[DESCRIPTION_LINE_LENGTH +1];
  if(/* Output line leaves the right side of a box and there is */
     /* a double headed arrow.
     (tem_line->start_position.x <= box->swcorner.x + BOX_WIDTH + 10) &&
     (tem_line->start_position.x >= box->swcorner.x + BOX_WIDTH)
                                                                       金金
     (tem_line->start_position.y <= box->swcorner.y)
     (tem_line->start_position.y >= box->swcorner.y - BOX_HT)
                                                                       28
     ((tem_line->start_activity_num & (ARROW_HEAD)) != 0)
                                                                       &&
     ((tem_line->start_activity_num & (DOT_B_L|DOT_T_R)) != 0))
      if(get2_first_label(tem_line,buf) == MYTRUE) {
         Line_labels = (struct text_line_struct *)
                       add_to_inputs_tree(Line_labels,buf);
         return;
}
```

```
/*********************************
* DATE: 19 Feb 1990
* VERSION: 1.0
* NAME: check_control_abox()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of generating a linked *
              list(Line_labels) which has the information of labels*
              of the control arrows attatched on an activity box.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: None
* GLOBAL VARIABLES CHANGED: None
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: single_headed_control(), double_headed_control() *
                 double_headed_control_with_slash()
* CALLING MODULES: search_labels_touched_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
************************************
void
check_control_abox(tem_line)
struct line_struct *tem_line;
 {
  extern struct text_line_struct *Line_labels;
   single_headed_control(tem_line);
   double_headed_control(tem_line);
   double_headed_control_with_slash(tem_line);
   return;
  }
```

```
* DATE: 19 Feb 1990
* VERSION: 1.0
* NAME: single_headed_control()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of adding the line
              labels of the control arrows with a single head to
              the linked list, Line_labels.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: line, Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get_closest_label(), add_to_inputs_tree()
                  get_branchnode(), find_branchnode(),
                  find_TO_ALL_branchnode()
* CALLING MODULES: check_control_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
single_headed_control(tem_line)
struct line_struct *tem_line;
 {
  extern get_closest_label(),add_to_inputs_tree(),get_branchnode();
  extern find_branchnode(),find_TO_ALL_branchnode();
  extern struct line_struct *line;
  struct line_struct *original_line;
  int original_line_type;
  char buf[DESCRIPTION_LINE_LENGTH +1];
  if(/* Control line comes to the upper side of a box and there
     /* is a single headed arrow.
                                                                  */
     (tem_line->end_position.x <= box->swcorner.x + BOX_WIDTH) &&
     (tem_line->end_position.x >= box->swcorner.x)
     (tem_line->end_position.y >= box->swcorner.y - BOX_HT - 10) &&
     (tem_line->end_position.y <= box->swcorner.y - BOX_HT)
                                                                 22
     ((tem_line->end_activity_num & (ARROW_HEAD)) != 0)
                                                                 22
     ((tem_line->end_activity_num & (DOT_B_L|DOT_T_R)) == 0))
```

```
{
     original_line_type = get_branchnode(tem_line);
     original_lime = (struct line_struct *)
                      find_branchnode(original_line_type);
     if(/* found a label on the line */
        get_closest_label(original_line,buf,tem_line->struct_type)
        == MYTRUE) {
        Line_labels = (struct text_line_struct *)
                        add_to_inputs_tree(Line_labels,buf);
return:
     if(/* No label on tem_line - Check if a FROM_ALL exists. */
/* If so, search TO_ALL and get a label for this line. */
original_line->start_activity_num == FROM_ALL) {
original_line = (struct line_struct *)
              find_TO_ALL_branchnode(original_line->to_from_label);
        original_line_type = get_branchnode(original_line);
original_line = (struct line_struct *)
                         find_branchnode(original_line_type);
        if(/* found the branchnode of the TO_ALL line segment */
  get_closest_label(original_line,buf,line->struct_type)
  == MYTRUE) {
  Line_labels = (struct text_line_struct *)
                          add_to_inputs_tree(Line_labels,buf);
          return;
  }
        strcpy(buf,"");
Line_labels = (struct text_line_struct *)
                       add_to_inputs_tree(Line_labels,buf);
     return;
```

```
* DATE: 19 Feb 1990
* VERSION: 1.0
* NAME : double_headed_control_with_slash()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of adding the labels
             of the control arrows with double heads and the
             slash to the linked list, Line_labels.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get2_closest_label(), add_to_inputs_tree(),
                get_branchnode(), find_TO_ALL_branchnode(),
                find_branchnode()
* CALLING MODULES: check_control_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
*****************
double_headed_control_with_slash(tem_line)
struct line_struct *tem_line;
 extern get2_closest_label(),add_to_inputs_tree(),get_branchnode();
 extern find_TO_ALL_branchnode(),find_branchnode();
 struct line_struct *original_line;
 char buf[DESCRIPTION_LINE_LENGTH +1];
 int original_line_type;
 if(/* Have a control with a double headed arrow */
    (tem_line->end_position.y >= box->swcorner.y-BOX_HT-15) &&
    (tem_line->end_position.x <= box->swcorner.x+BOX_WIDTH) &&
    (tem_line->end_position.x >= box->swcorner.x) &&
    ((tem_line->end_activity_num & (ARROW_HEAD)) != 0) &&
    ((tem_line->end_activity_num & (DOT_B_L)) != 0))
     original_line_type = get_branchnode(tem_line);
```

```
original_line = (struct line_struct *)
                        find_branchnode(original_line_type);
       if(get2_closest_label(original_line,buf,tem_line->struct_type)
== MYTRUE) {
Line_labels = (struct text_line_struct *)
                        add_to_inputs_tree(Line_labels,buf);
        return;
}
       if(/* No label on tem_line - Check if tem_line is a FROM_ALL */
 /* type line. If so, search the TO_ALL and get a label
                                                            */
                                                            */
 /* from TO_ALL line.
 original_line->start_activity_num == FROM_ALL) {
 original_line = (struct line_struct *)
                 find_TO_ALL_branchnode(original_line->to_from_label);
          original_line_type = get_branchnode(original_line);
 original_line = (struct line_struct *)
                           find_branchnode(original_line_type);
 if(get2_closest_label(original_line,buf,line->struct_type)
     == MYTRUE) {
    Line_labels = (struct text_line_struct *)
                            add_to_inputs_tree(Line_labels,buf);
             return;
    }
       strcpy(buf,"");
Line_labels = (struct text_line_struct *)
                       add_to_inputs_tree(Line_labels,buf);
     }
     return;
     }
```

```
* DATE: 20 Feb 1990
 * VERSION: 1.0
 * NAME: double_headed_control()
* MODULE NUMBER:
 * DESCRIPTION: This module provides a means of adding the line
              labels of the control arrows with double head to
              the linked list, Line_labels.
 * ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: Nnone
 * HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get3_first_label(), add_to_inputs_tree()
 * CALLING MODULES: check_control_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
 * ORDER OF:
 **********************
double_headed_control(tem_line)
struct line_struct *tem_line;
  {
   extern get3_first_label(),add_to_inputs_tree();
   char buf[DESCRIPTION_LINE_LENGTH +1];
   if((tem_line->start_position.x <= box->swcorner.x+BOX_WIDTH+10) &&
      (tem_line->start_position.x >= box->swcorner.x+BOX_WIDTH)
      (tem_line->start_position.y <= box->swcorner.y)
                                                                  28
      (tem_line->start_position.y >= box->swcorner.y-BOX_HT)
      ((tem_line->start_activity_num & (ARROW_HEAD)) != 0)
                                                                  数数
      ((tem_line->start_activity_num & (DOT_T_R)) != 0))
if(get3_first_label(tem_line,buf) == MYTRUE) {
  Line_labels = (struct text_line_struct *)
                         add_to_inputs_tree(Line_labels,buf);
          return;
  }
        strcpy(buf,"");
        Line_labels = (struct text_line_struct *)
```

```
add_to_inputs_tree(Line_labels,buf);
}
return;
}
```

```
* DATE: 21 Feb 1990
* VERSION: 1.0
* NAME: check_mechanism_abox()
* MODULE NUMBER:
* DESCRIPTION: This module provides means of generating a linked
            list, Line_labels, with all information labels of the *
            mechanism arrows attatched on an activity box.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: Line_labels
* GLOBAL VARIABLES CHANGED: None
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: single_headed_mechanism()
* CALLING MODULES: search_labels_touched_abox()
* AUTHOR: Intack Kim
* HISTORY:
* ABSTRACT DATA TYPE:
* ORDER OF:
check_mechanism_abox(tem_line)
struct line_struct *tem_line;
  extern struct text_line_struct *Line_labels;
  single_headed_mechanism(tem_line);
  return;
 }
```

```
/*****************************
* DATE: 21 Feb 1990
* VERSION: 1.0
* NAME: single_headed_mechanism()
* MODULE NUMBER:
* DESCRIPTION: This module provides a means of adding the line
              labels of the mechanism arrows with a single head to
              a linked list, Line_labels.
* ALGORITHM:
* PASSED VARIABLES: tem_line
* RETURNS: None
* GLOBAL VARIABLES USED: line, Line_labels
* GLOBAL VARIABLES CHANGED: Line_labels
* FILES READ: None
* FILES WRITTEN: None
* HARDWARE INPUT: None
* HARDWARE OUTPUT: None
* MODULES CALLED: get_branchnode(), find_branchnode(),
                  get_closest_label(), find_TO_ALL_branchnode()
                  add_to_inputs_tree()
* CALLING MODULES: check_mechanism_abox()
* AUTHOR: Intack Kim
* HISTORY:
 * ABSTRACT DATA TYPE:
 * ORDER OF:
                  ************
single_headed_mechanism(tem_line)
struct line_struct *tem_line;
   extern get_branchnode(),find_branchnode(),get_closest_label();
   extern find_TO_ALL_branchnode(),add_to_inputs_tree();
   extern struct line_struct *line;
   struct line_struct *original_line;
   int original_line_type;
   char buf[DESCRIPTION_LINE_LENGTH +1];
   if((tem_line->end_position.x <= box->swcorner.x+BOX_WIDTH) &&
      (tem_line->end_position.x >= box->swcorner.x)
                                                              数数
      (tem_line->end_position.y >= box->swcorner.y)
                                                              22
      (tem_line->end_position.y <= box->swcorner.y+15))
       original_line_type = get_branchnode(tem_line);
       original_line = (struct line_struct *)
                        find_branchnode(original_line_type);
```

```
if(get_closest_label(original_line,buf,tem_line->struct_type)
     == MYTRUE) {
     Line_labels = (struct text_line_struct *)
                    add_to_inputs_tree(Line_labels,buf);
     return:
   if(original_line->start_activity_num == FROM_ALL) {
     original_line = (struct line_struct *)
          find_TO_ALL_branchnode(original_line->to_from_label);
     original_line_type = get_branchnode(original_line);
     original_line = (struct line_struct *)
                      find_branchnode(original_line_type);
     if(get_closest_label(original_line,buf,line->struct_type)
        == MYTRUE) {
       Line_labels = (struct text_line_struct *)
                       add_to_inputs_tree(Line_labels,buf);
       return;
       }
   strcpy(buf,"");
   Line_labels = (struct text_line_struct *)
                  add_to_inputs_tree(Line_labels,buf);
   }
return;
}
```

```
DATE: 15 Feb 1990
    VERSION: 1.0
   NAME: search_labels_touched_abox()
   MODULE NUMBER:
    DESCRIPTION:
                This module provides a means of searching the line
                label in accordance with the ICOM code.
    ALGORITHM:
    PASSED VARIABLES: tem_line
    RETURNS: None
    GLOBAL VARIABLES USED: Line_labels
    GLOVAL VARIABLES CHANGED: None
   FILES READ: None
    FILES WRITTEN: None
   HARDWARE INPUT: None
    HARDWARE OUTPUT: None
    MODULES CALLED: check_input_abox(), check_output_abox(),
                     check_control_abox().
                     check_mechanism_abox(),
    CALLING MODULES: get_labels_for_abox()
   AUTHOR: Intack Kim
    HISTORY:
    ABSTRACT DATA TYPE :
     ORDER OF :
  search_labels_touched_abox(tem_line,indicator)
   struct line_struct *tem_line;
   char indicator;
     extern struct text_line_struct *Line_labels;
    if(tem_line == NULL) return;
     else {
      search_labels_touched_abox(tem_line->left,indicator);
      search_labels_touched_abox(tem_line->right,indicator);
      switch(indicator) {
      case 'I':
       check_input_abox(tem_line);
break;
```

```
DATE: 15 Feb 1990
    VERSION: 1.0
   NAME: get_labels_for_abox
   MODULE NUMBER:
    DESCRIPTION:
         This module provides a means of constructing a linked list *
         (Line_labels) made of the line labels in accordance with
         the variable, indicator.
    ALGORITHM:
    PASSED VARIABLES: indicator
   RETURNS: None
   GLOBAL VARIABLES USED: line_rootnode, Line_labels
   GLOVAL VARIABLES CHANGED: Line_labels
   FILES READ: None
   FILES WRITTEN: None
   HARDWARE INPUT: none
   HARDWARE OUTPUT: None
   MODULES CALLED: search_labels_touched_abox()
   CALLING MODULES: save_arrow_info_of_abox()
   AUTHOR: Intack Kim
   HISTORY:
    ABSTRACT DATA TYPE :
    ORDER OF :
***********************************
struct text_line_struct *
get_labels_for_abox(indicator)
 char indicator;
   extern struct line_struct *line_rootnode;
   extern struct text_line_struct *Line_labels;
   struct line_struct *temporary_line;
   Line_labels = NULL;
   temporary_line = line_rootnode;
   while(temporary_line != NULL) {
     search_labels_touched_abox(temporary_line,indicator);
     temporary_line = temporary_line->next;
   return(Line_labels);
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: save_arrow_info_of_abox()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of saving all information of
         arrows(ICOM) attatched on a box to a file.
   ALGORITHM:
   PASSED VARIABLES: fp, tem_box
  RETURNS: None
   GLOBAL VARIABLES USED: box
   GLOVAL VARIABLES CHANGED:
  FILES READ: None
  FILES WRITTEN: *.pro, CHECKBOX.PRO, or CHECKBOUNDARY.PRO
  HARDWARE INPUT: None
  HARDWARE OUTPUT: None
  MODULES CALLED: get_labels_for_abox(), fprintf(), itoa()
  CALLING MODULES: create_temp_box_info(), traverse_boxes()
  AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
******************
void
save_arrow_info_of_abox(fp,tem_box)
FILE
        *fp;
 struct box_struct *tem_box;
  extern itoa();
  extern struct box_struct *box;
  struct text_line_struct *ICOM_labels;
  char buf[DESCRIPTION_LINE_LENGTH+1];
  int number_of_input = 0;
  int number_of_output = 0;
  int number_of_control = 0;
  int number_of_mechanism = 0;
/******** NAME ********/
  fprintf(fp,"confirmed(['%s',is,in_data]).\n",
```

```
tem_box->name.text_string);
 /*********** NUMBER ********/
   itoa(tem_box->number,buf);
   fprintf(fp, "confirmed(['%s', number_is, %s]).\n",
   tem_box->name.text_string,buf);
   box = tem_box;
 ICOM_labels = (struct text_line_struct *)get_labels_for_abox('I');
   if (ICOM_labels == NULL) {
     fprintf(fp, "confirmed(['%s',input_is,null]).\n",
     tem_box->name.text_string);
   else {
    while(ICOM_labels != NULL) {
     fprintf(fp,"confirmed(['%s',input_is,'%s']).\n",
             tem_box->name.text_string, ICOM_labels->text_line);
     ICOM_labels = ICOM_labels->next;
     number_of_input += 1;
     }
    itoa(number_of_input,buf);
    fprintf(fp, "confirmed(['%s', has_input_number, %s]).\n",
            tem_box->name.text_string,buf);
    }
 /**************** OUTPUTS *******************/
   ICOM_labels = (struct text_line_struct *)get_labels_for_abox('0');
   if (ICOM_labels == NULL) {
     fprintf(fp,"confirmed(['%s',output_is,null]).\n",
             tem_box->name.text_string);
     }
   else {
    while(ICOM_labels != NULL) {
     fprintf(fp, "confirmed(['%s',output_is,'%s']).\n",
             tem_box->name.text_string,ICOM_labels->text_line);
     ICOM_labels = ICOM_labels->next:
     number_of_output += 1;
     }
    itoa(number_of_output,buf);
    fprintf(fp,"confirmed(['%s', has_output_number, %s]).\n",
            tem_box->name.text_string, buf);
   }
```

```
ICOM_labels = (struct text_line_struct *)get_labels_for_abox('C');
  if (ICOM_labels == NULL) {
    fprintf(fp, "confirmed(['%s',control_is,null]).\n",
           tem_box->name.text_string);
    }
  else {
   while(ICOM_labels != NULL) {
    fprintf(fp,"confirmed(['%s',control_is,'%s']).\n",
           tem_box->name.text_string,ICOM_labels->text_line);
    ICOM_labels = ICOM_labels->next;
    number_of_control += 1;
    }
   itoa(number_of_control,buf);
   fprintf(fp,"confirmed(['%s', has_control_number, %s]).\n",
          tem_box->name.text_string.buf);
   }
ICOM_labels = (struct text_line_struct *)get_labels_for_abox('M');
  if (ICOM_labels == NULL) {
   fprintf(fp, "confirmed(['%s',mechanism_is,null]).\n",
          tem_box->name.text_string);
   }
  else {
   while(ICOM_labels != NULL) {
    fprintf(fp, "confirmed(['%s', mechanism_is, '%s']). \n",
           tem_box->name.text_string,ICOM_labels->text_line);
    ICOM_labels = ICOM_labels->next;
    number_of_mechanism += 1;
    }
   itoa(number_of_mechanism,buf);
   fprintf(fp, "confirmed(['%s', has_mechanism_number, %s]).\n",
          tem_box->name.text_string,buf);
   }
  return;
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
  NAME: create_temp_box_info()
   MODULE NUMBER:
   DESCRIPTION:
               This module provides a means of creating the
         temporary file, CHECKBOX.PRO, which contains a set of
         predicate forms for an activity box.
   ALGORITHM:
   PASSED VARIABLES: found_box
   RETURNS: None
   GLOBAL VARIABLES USED: None
   GLOVAL VARIABLES CHANGED: None
   FILES READ: None
  FILES WRITTEN: CHECKBOX.PRO
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: fopen(), put_message(), disable_input_window(),
                   save_header_info(), save_arrow_info_of_abox(),
                   fclose(), printf()
   CALLING MODULES: find_clicked_box()
  AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
                      ****************
void
create_temp_box_info(found_box)
 struct box_struct *found_box;
  extern put_message(),disable_input_window();
 FILE *fp;
  if((fp = fopen("CHECKBOX.PRO","w")) == NULL) {
   put_message(1,"Unable to open the CHECKBOX.PRO file -- ABORT.");
   disable_input_window();
   }
  else {
   disable_input_window();
   save_header_info(fp);
```

```
save_arrow_info_of_abox(fp,found_box);
if(fclose(fp) != 0) printf("FILE CLOSE FAILED\n");
}
return;
}
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: find_clicked_box()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of checking if there is a box *
          within the cordinate(x,y) clicked by the user mouse.
   ALGORITHM:
   PASSED VARIABLES: x,y
   RETURNS: None
   GLOBAL VARIABLES USED: box_rootnode
   GLOVAL VARIABLES CHANGED: None
  FILES READ: None
   FILES WRITTEN: None
  HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: put_message(), create_temp_box_info(),
                 my_window_set(), null_proc()
  CALLING MODULES: check_button_for_activity()
  AUTHOR: Intack Kim
  HISTORY:
   ABSTRACT DATA TYPE :
   CRDER OF :
void
find_clicked_box(x,y)
 int x,y;
  extern put_message(),my_window_set(),null_proc();
  extern struct box_struct *box_rootnode;
  struct box_struct *bbox;
  bbox = box_rootnode;
  while (bbox != NULL) {
  if(x >= bbox->swcorner.x && x <= bbox >swcorner.x + BOX_WIDTH &&
     y >= bbox->swcorner.y - BOX_HT && y <= bbox->swcorner.y)
      put_message(1,"Enter the prolog environment
     using another window.");
      create_temp_box_info(bbox);
```

```
my_window_set(null_proc);
    return;
}
bbox = bbox->next;
}
put_message(1,"Box not found - Try again(|L| to select)");
return;
}
```

```
/********************************
    DATE: 15 Feb 1990
    VERSION: 1.0
    NAME: check_button_for_activity()
    MODULE NUMBER:
    DESCRIPTION:
         This module provides a means of checking the button clicked*
         by user mouse if the clicked button is right or left.
    ALGORITHM:
    PASSED VARIABLES: window, event, arg(Sun variable)
    RETURNS: None
    GLOBAL VARIABLES USED: window, event, arg
    GLOVAL VARIABLES CHANGED: None
    FILES READ: None
    FILES WRITTEN: None
    HARDWARE INPUT: None
    HARDWARE OUTPUT: None
    MODULES CALLED: event_id(), find_clicked_box(), event_x(),
                   event_y(), my_window_set(), null_proc(),
                   put_message()
    CALLING MODULES: check_activity()
    AUTHOR: Intack Kim
    HISTORY:
    ABSTRACT DATA TYPE :
    ORDER OF :
****************
check_button_for_activity(window, event, arg)
 Window window;
 Event *event;
 caddr_t arg;
  extern my_window_set();
  extern put_message(), null_proc();
  /* Check for left or right button */
  switch(event_id(event)) {
    case MS_LEFT:
      if(event_is_up(event))
         find_clicked_box(event_x(event), event_y(event));
```

```
break;

case MS_RIGHT:
    my_window_set(null_proc);
    put_message(1,"ABORT -- Make another selection.");
    break;

default:
    break;
}
return;
}
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: check_activity()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of finally producing a file
         contained a set of predicate forms for an activity box.
   ALGORITHM:
   PASSED VARIABLES: None
   RETURNS: None
   GLOBAL VARIABLES USED: header_rootnode, box_rootnode
   GLOVAL VARIABLES CHANGED: None
   FILES READ: None
   FILES WRITTEN: None
    HARDWARE INPUT: None
   HARDWARE OUTPUT: None
    MODULES CALLED: put_message(), strcmp(), my_move_cursor(),
                  my_window_set()
   CALLING MODULES: make_windows()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
    ORDER OF :
void
check_activity()
    extern put_message(), my_window_set(), my_move_cursor();
    extern struct box_struct *box_rootnode;
    extern struct header_struct *header_rootnode;
    if(box_rootnode == NULL) {
      put_message(1,"FATAL: Can't check this empty diagram
     -- Make another selection.");
      return:
    if(strcmp(header_rootnode->title.text_string,"") == 0) {
     put_message(1,"NO TITLE: Please enter the TITLE
    using EDIT DGM and then RETRY!!!");
    return;
```

```
my_move_cursor(INIT_LOC_X,INIT_LOC_Y);
my_window_set(check_button_for_activity);
put_message(1,"Move cursor inside activity box
and click left button - Right to ABORT.");
return;
}
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: create_temp_boundary_info()
   MODULE NUMBER:
   DESCRIPTION:
          This module provides a means of saving all information of
         boundary arrows of the IDEFO diagram into the tempory
         file CHECKBOUNDARY.PRO.
   ALGORITHM:
   PASSED VARIABLES: parentfile
   RETURNS: None
   GLOBAL VARIABLES USED: header_rootnode
   GLOVAL VARIABLES CHANGED: None
   FILES READ: parentfile
   FILES WRITTEN: CHECKBOUNDARY.PRO
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: fopen(), put_message(), disable_input_window()
                    getc(), putc(), fclose(), search_boundary_info(),*
                    save_null_boundary(), printf()
   CALLING MODULES: check_parent_box_file()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
void
create_temp_boundary_info(parentfile)
 char parentfile[];
 {
   extern put_message(),disable_input_window();
  FILE *parentfp, *childfp;
   extern struct header_struct *header_rootnode;
   extern int number_of_boundary_input,number_of_boundary_output;
   extern int number_of_boundary_control,number_of_boundary_mechanism;
   int ch:
  number_of_boundary_input = 0;
  number_of_boundary_output = 0;
  number_of_boundary_control = 0;
```

```
number_of_boundary_mechanism = 0;
 if((parentfp = fopen(parentfile,"r")) == NULL ||
    (childfp = fopen("CHECKBOUNDARY.PRO","w")) == NULL) {
   put_message(1,"Unable to open the predicate file(s) -- ABORT");
    disable_input_window();
    }
 else {
    disable_input_window();
    while((ch = getc(parentfp)) != EOF)
     putc(ch,childfp);
   fclose(parentfp);
   search_boundary_info(childfp);
    save_null_boundary(childfp);
    if(fclose(childfp) != 0) printf("FILE CLOSE FAILED\n");
   }
return;
}
```

```
******************
   DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: check_parent_box_file()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of checking if there is the
         file which the user specifies in the current directory.
   ALGORITHM:
   PASSED VARIABLES: None
   RETURNS: None
   GLOBAL VARIABLES USED: None
   GLOVAL VARIABLES CHANGED: None
   FILES READ: None
   FILES WRITTEN: None
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: strcpy(), panel_get_value(), fix_input(),
                   strcmp(), put_message(), disable_input_window(), *
                   strcat(), check_filename(),
                   create_temp_boundary_info()
   CALLING MODULES: check_button_for_boundary()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
                      ****************
check_parent_box_file()
{
 extern put_message(), disable_input_window(),check_filename();
 extern my_move_cursor(),my_window_set(),fix_input();
 char name[FILE_NAME_LENGTH + 1], name2[FILE_NAME_LENGTH + 5];
 int file_type_indicator;
 strcpy(name,(char *)panel_get_value(input_item));
 fix_input(name);
 if(strcmp(name,"") ==0) {
   put_message(1,"ABORT: No file name received
  -- Make another selection.");
   disable_input_window();
   return(PANEL_NONE);
```

```
}
strcpy(name2,name);
strcat(name2,".pro");
file_type_indicator = check_filename(name2);
switch (file_type_indicator) {
 case -1:
   disable_input_window();
  put_message(1,"ABORT: File does not exist
  -- Make another selection.");
   break;
 case -3:
   disable_input_window();
   put_message(1,"ABORT: File is a directory
  -- Make another selection.");
  break;
 case -2: /* READ ONLY */
 case 0: /* READ/WRITE.. */
   put_message(1,"Enter prolog environment using another window.");
   create_temp_boundary_info(name2);
   break;
 default:
   put_message(1,"Unknown condition -- Make another selection.");
  disable_input_window();
  break;
 }
return(PANEL_NONE);
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
  NAME: check_button_for_boundary()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of conforming if the user
         would like to be continue to check IDEFO syntax for the
         boundary arrows in any IDEFO diagram.
   ALGORITHM:
   PASSED VARIABLES: window, event, arg(Sun variables)
  RETURNS: None
   GLOBAL VARIABLES USED: None
  GLOVAL VARIABLES CHANGED: None
  FILES READ: None
  FILES WRITTEN: None
  HARDWARE INPUT: None
  HARDWARE OUTPUT: None
  MODULES CALLED: event_is_up(), event_id(), enable_input_window(),*
                   panel_set(), check_parent_box_file(),
                   put_message(), my_window_set(), null_proc()
  CALLING MODULES: check_boundary()
  AUTHOR: Intack Kim
  HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
check_button_for_boundary(window, event, arg)
 Window window;
 Event
         *event;
 caddr_t arg;
 extern put_message(),enable_input_window(),my_window_set();
 extern null_proc();
 if(event_is_up(event)) return;
 switch event_id(event)
  {
    case MS_LEFT:
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: check_boundary()
   MODULE NUMBER:
   DESCRIPTION:
          This module provides a means of producing all information
          of a set of predicate forms of the boundary arrows in the
          IDEFO diagram so as to check IDEFO syntax of boundary
          arrows.
   ALGORITHM:
   PASSED VARIABLES: None
  RETURNS: None
   GLOBAL VARIABLES USED: header_rootnode, line_rootnode
   GLOVAL VARIABLES CHANGED: None
   FILES READ: None
   FILES WRITTEN: None
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: disable_input_window(), put_message(), strcmp(), *
                    my_move_cursor(), my_window_set()
   CALLING MODULES: make_windows()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
check_boundary()
  extern put_message(),my_window_set();
  extern my_move_cursor(),disable_input_window();
  extern struct box_struct *box_rootnode;
  extern struct line_struct *line_rootnode;
  extern struct header_struct *header_rootnode;
  if(box_rootnode == NULL) {
   disable_input_window();
   put_message(1,"FATAL: Can't check this empty diagram
   -- Make another selection.");
   return:
```

```
if(line_rootnode == NULL) {
  put_message(1,"FATAL: Can't check this diagram
 -- Make another selection.");
  disable_input_window();
  return;
  }
if(strcmp(header_rootnode->title.text_string,"") == 0) {
  put_message(1,"NO TITLE: Please enter the TITLE using EDIT DGM
 and then RETRY!!!");
  disable_input_window();
  return;
  }
put_message(1,"CHECK BOUNDARY ARROW : |L| to check - |R| to ABORT");
my_move_cursor(INIT_LOC_X,INIT_LOC_Y);
my_window_set(check_button_for_boundary);
return;
}
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
  NAME: save_null_boundary()
  MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of saving the information
         which contains that there is no boundary arrow in
         accordance with ICOM.
  ALGORITHM:
  PASSED VARIABLES: fp(file pointer)
* RETURNS: None
   GLOBAL VARIABLES USED: number_of_boundary_input.
                          number_of_boundary_output,
                          number_of_boundary_control,
                          number_of_boundary_mechanism
  GLOVAL VARIABLES CHANGED: None
   FILES READ: None
  FILES WRITTEN: fp
   HARDWARE INPUT: None
  HARDWARE OUTPUT: None
  MODULES CALLED: fprintf(), itoa()
   CALLING MODULES: create_temp_boundary_info()
  AUTHOR: Intack Kim
   HISTORY:
  ABSTRACT DATA TYPE :
   ORDER OF :
                    *****************
save_null_boundary(fp)
FILE
      *fp;
₹
 extern int number_of_boundary_input,number_of_boundary_output;
 extern int number_of_boundary_control,number_of_boundary_mechanism;
 extern itoa();
 char buf[DESCRIPTION_LINE_LENGTH+1];
 if (number_of_boundary_input == 0)
    fprintf(fp,"confirmed([boundary_input, is, null]).\n");
 if (number_of_boundary_output == 0)
    fprintf(fp,"confirmed([boundary_output, is, null]).\n");
 if (number_of_boundary_control == 0)
```

```
fprintf(fp,"confirmed([boundary_control, is, null]).\n");
if (number_of_boundary_mechanism == 0)
    fprintf(fp,"confirmed([boundary_mechanism, is, null]).\n");
itoa(number_of_boundary_input, buf);
fprintf(fp,"confirmed([boundary_input, has_number, %s]).\n",buf);
itoa(number_of_boundary_output,buf);
fprintf(fp,"confirmed([boundary_output, has_number, %s]).\n",buf);
itoa(number_of_boundary_control,buf);
fprintf(fp,"confirmed([boundary_control, has_number, %s]).\n",buf);
itoa(number_of_boundary_mechanism,buf);
fprintf(fp,"confirmed([boundary_mechanism, has_number, %s]).\n",buf);
return(MYTRUE);
}
```

```
DATE: 15 Feb 1990
  VERSION: 1.0
  NAME: search_NLR_boundary_line_info()
  MODULE NUMBER:
  DESCRIPTION:
              This module provides a means of working through the
              tree of line structure in left and right direction.
  ALGORITHM:
  PASSED VARIABLES: line_info, fp
  RETURNS: None
  GLOBAL VARIABLES USED: None
  GLOVAL VARIABLES CHANGED: None
  FILES READ: None
  FILES WRITTEN: None
  HARDWARE INPUT: None
  HARDWARE OUTPUT: None
  MODULES CALLED: save_boundary_line_info(),
                   search_NLR_boundary_line_info()
 CALLING MODULES: search_boundary_info()
 AUTHOR: Intack Kim
  HISTORY:
  ABSTRACT DATA TYPE :
  ORDER OF :
void
search_NLR_boundary_line_info(line_info,fp)
struct line_struct *line_info;
FILE
      *fp;
 extern int number_of_boundary_input,number_of_boundary_output;
 extern int number_of_boundary_control,number_of_boundary_mechanism;
 if(line_info == NULL) return;
 else
    save_boundary_line_info(line_info,fp);
    search_NLR_boundary_line_info(line_info->left,fp);
    search_NLR_boundary_line_info(line_info->right,fp);
    }
return:
```

}

```
|*******************************
    DATE: 15 Feb 1990
    VERSION: 1.0
    NAME: search_boundary_info()
    MODULE NUMBER:
    DESCRIPTION:
               This module provides a means of working through the
               tree of the line structure in next direction.
   ALGORITHM:
    PASSED VARIABLES: fp
    RETURNS: None
    GLOBAL VARIABLES USED: line_rootnode
    GLOVAL VARIABLES CHANGED: None
    FILES READ: None
    FILES WRITTEN: None
    HARDWARE INPUT: None
    HARDWARE OUTPUT: None
    MODULES CALLED: fprintf(), search_NLR_boundary_line_info()
    CALLING MODULES: create_temp_boundary_info()
* AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
    ORDER OF :
                        ***********
search_boundary_info(fp)
  FILE
           *fp;
   extern struct line_struct *line_rootnode;
   struct line_struct *line_info;
   extern int number_of_boundary_input,number_of_boundary_output;
   extern int number_of_boundary_control,number_of_boundary_mechanism;
   fprintf(fp,"confirmed([child_title, is, '%s']).\n",
           header_rootnode->title.text_string);
   line_info = line_rootnode;
   while(line_info != NULL)
      search_NLR_boundary_line_info(line_info,fp);
      line_info = line_info->next;
      }
   return(MYTRUE);
```

}

```
/********************
    DATE: 4 Feb 1990
    VERSION: 1.0
   NAME : save_boundary_line_info()
   MODULE NUMBER :
    DESCRIPTION:
       This module is to save the information of the boundary
       arrows in the IDEFO diagram.
    ALGORITHM :
    PASSED VARIABLES: line_info, fp (File Pointer)
    RETURNS: None
   GLOBAL VARIABLES USED : None
    GLOBAL VARIABLES CHANGED: None
   FILES READ : None
   FILES WRITTEN : fp
   HARDWARE INPUT : None
   HARDWARE OUTPUT : None
   MODULES CALLED : itoa(), fprintf()
   CALLING MODULES : search_NLR_boundary_line_info()
   AUTHOR: Intack Kim
    HISTORY:
    ABSTRACT DATA TYPE:
    ORDER OF:
 save_boundary_line_info(line_info,fp)
  struct line_struct *line_info;
  FILE
           *fp;
   extern itoa();
   extern int number_of_boundary_input,number_of_boundary_output;
   extern int number_of_boundary_control,number_of_boundary_mechanism;
   char buf[DESCRIPTION_LINE_LENGTH+1];
   switch(line_info->start_ICOM[0])
      case 'I':
       number_of_boundary_input += 1;
       itoa(number_of_boundary_input,buf);
       fprintf(fp,"confirmed([boundary_input%s, is, '%s']).\n",
               buf,line_info->label.text_string);
```

```
break;
  case 'C':
   number_of_boundary_control += 1;
   itoa(number_of_boundary_control,buf);
   fprintf(fp, "confirmed([boundary_control%s, is, '%s']).\n",
           buf,line_info->label.text_string);
   break;
  case 'M':
  number_of_boundary_mechanism += 1;
   itoa(number_of_boundary_mechanism,buf);
   fprintf(fp,"confirmed([boundary_mechanism%s, is, '%s']).\n",
           buf ,line_info->label.text_string);
  break;
  default:
   if (line_info->end_ICOM[0] == '0')
      number_of_boundary_output += 1;
       itoa(number_of_boundary_output,buf);
       fprintf(fp,"confirmed([boundary_output%s, is, '%s']).\n",
               buf,line_info->label.text_string);
      break;
  }
return(MYTRUE);
}
```

```
DATE: 2 Feb 1990
  VERSION: 1.0
 NAME : save_header_info()
 MODULE NUMBER :
 DESCRIPTION:
       This module is to save the information of "NODE" and
       "TITLE" in the IDEFO diagram.
 ALGORITHM:
 PASSED VARIABLES : fp (File pointer)
 RETURNS : None
  GLOBAL VARIABLES USED : header_rootnode
 GLOBAL VARIABLES CHANGED: None
 FILES READ : None
 FILES WRITTEN : fp
 HARDWARE INPUT : None
 HARDWARE DUTPUT : None
 MODULES CALLED : fprintf()
 CALLING MODULES : store_predicates()
 AUTHOR : Intack Kim
 HISTORY:
 ABSTRACT DATA TYPE:
  ORDER OF:
save_header_info(fp)
 FILE
          *fp;
 {
  extern put_message(),disable_input_window();
  extern struct header_struct *header_rootnode;
  fprintf(fp,"confirmed([title, is, '%s']).\n",
          header_rootnode->title.text_string);
  fprintf(fp,"confirmed([node, is, '%s']).\n",
          header_rootnode->node.text_string);
  return(MYTRUE);
 }
```

```
DATE: 20 Feb 1990
   VERSIOIN: 1.0
   NAME:
                  traverse_boxes()
   MODULE NUMBER:
   DESCRIPTION:
             This module provides a means of traversing all the
             activity boxes in a diagram and of passing the
             information of each box to store_diagram().
   ALGORITHM:
   PASSED VARIABLES:
                         fp
  RETURNS: None
   GLOBAL VARIABLES USED: box_rootnode
   GLOBAL VARIABLES CHANGED: None
   FILES READ: None
  FILES WRITTEN: None
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: save_arrow_info_of_abox()
   CALLING MODULES: store_predicates()
  AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE:
   ORDER OF:
*****************
 traverse_boxes(fp)
  FILE
        *fp;
   extern struct box_struct *box_rootnode;
   struct box_struct *temp_box;
   temp_box = box_rootnode;
   while(temp_box != NULL)
    save_arrow_info_of_abox(fp,temp_box);
    temp_box = temp_box->next;
    }
   return;
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: store_predicates()
   MODULE NUMBER:
   DESCRIPTION:
         This module provides a means of saving all information of
         arrows(ICOM) attatched on the activity boxes in the IDEFO
         diagram.
   ALGORITHM:
   PASSED VARIABLES: file_name
  RETURNS: None
  GLOBAL VARIABLES USED: None
   GLOVAL VARIABLES CHANGED: None
  FILES READ: None
  FILES WRITTEN: None
  HARDWARE INPUT: None
  HARDWARE OUTPUT: None
   MODULES CALLED: fopen(), put_message(), disable_input_window(),
                   save_header_info(), traverse_boxes(), fclose()
  CALLING MODULES: overwrite_predictes()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
****
void
store_predicates(file_name)
  char file_name[];
  {
   extern put_message(),disable_input_window();
   FILE
          *fp;
   if ((fp = fopen(file_name,"w")) == NULL) {
       put_message(1,"Unable to open the file for predicates
                       -- ABORT.");
       disable_input_window();
   else {
         disable_input_window();
         save_header_info(fp);
```

```
traverse_boxes(fp);
  if(fclose(fp) != 0) printf("FILE CLOSE FAILED\n");
  }
  return;
}
```

```
DATE: 15 Feb 1990
   VERSION: 1.0
   NAME: overwrite_predicates()
   MODULE NUMBER:
   DESCRIPTION:
        This module provides a means of overwriting the predicate *
        forms into the existing flie which is specified by the user*
   ALGORITHM:
   PASSED VARIABLES: window, event, arg(Sun variables)
   RETURNS: None
   GLOBAL VARIABLES USED: last_file_name
   GLOVAL VARIABLES CHANGED: None
   FILES READ: None
   FILES WRITTEN: file pointer passed to a .pro file
   HARDWARE INPUT: None
   HARDWARE DUTPUT: None
   MODULES CALLED: event_is_up(), event_id(), strcpy(),
                 store_predicates(), put_message(), null_proc(),
                 my_window_set()
   CALLING MODULES: get_filename_for_predicates()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE :
   ORDER OF :
*******************
overwrite_predicates(window, event, arg)
  Window window;
  Event *event;
  caddr_t arg;
   extern null_proc(), put_message(),my_window_set();
   char file_name[FILE_NAME_LENGTH + 5];
   if (!event_is_up(event)) return;
   switch(event_id(event))
     case MS_LEFT:
       strcpy(file_name,last_file_name);
       store_predicates(file_name);
```

```
put_message(1,"OVERWRITE DONE -- Make another selection.");
   my_window_set(null_proc);
   break;

case MS_RIGHT:
   my_window_set(null_proc);
   put_message(1,"ABORT overwrite -- Make another selection.");
   break;
   }
   return;
}
```

```
DATE:
             10 Jan 1990
   VERSION: 1.0
   NAME: get_filename_for_predicates()
   MODULE NUMBER:
   DESCRIPTION:
        The purpose of this module is to get the file name from the
        user in which to save the predicates file.
   ALGORITHM:
   PASSED VARIABLES: None
                   PANEL_NONE (Sunview variable)
   RETURNS:
   GLOBAL VARIABLES USED: None
   GLOBAL VARIABLES CHANGED: None
   FILES READ: None
   FILES WRITTEN: None
   HARDWARE INPUT: None
   HARDWARE OUTPUT: None
   MODULES CALLED: put_message(),fix_input(),disable_input_window(),*
                    check_filename(),my_window_set(),stordicates(),
                    my_move_cursor();
   CALLING MODULES: save_predicates()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE:
   ORDER OF:
Panel_setting
get_filename_for_predicates()
  {
   extern put_message(),disable_input_window();
   extern fix_input(), my_window_set(),my_move_cursor();
   extern check_filename();
   char name[FILE_NAME_LENGTH+1], name2[FILE_NAME_LENGTH+5];
   int file_type_indicator;
   /* get the user input(file name) */
   strcpy(name,(char *)panel_get_value(input_item));
   fix_input(name); /* Remove blanks and replace \n to \0 */
   if(strcmp(name,"")==0)
   {
    put_message(1,"OPERATION ABORTED -- NO FILE NAME RECEIVED
```

```
-- Make another selection");
   disable_input_window();
  return(PANEL_NONE);
strcpy(name2, name);
strcat(name2,".pro");
/* Checks file name of "name2" is what type of file. */
file_type_indicator = check_filename(name2);
switch(file_type_indicator)
  case -1:
   store_predicates(name2);
   put_message(1,"SAVE DONE -- Make another selection.");
   break;
  case -2:
   disable_input_window();
   put_message(1,"Can't overwrite file -- READ ONLY
 -- Make another selection.");
   break;
  case -3:
   disable_input_window();
   put_message(1,"File is a DIRECTORY -- Make another selection.");
  break;
  case 0:
  put_message(1,"FILE EXISTS - |L| to overwrite, - |R| to ABORT.");
   my_move_cursor(INIT_LOC_X,INIT_LOC_Y);
   strcpy(last_file_name,name2);
   my_window_set(overwrite_predicates);
   break:
  default:
  put_message(1,"Unknown condition -- Make another selection");
   disable_input_window();
  break:
  }
return(PANEL_NONE);
}
```

```
DATE:
            10 Jan 1990
    VERSION: 1.0
   NAME: save_predicates()
   MODULE NUMBER:
    DESCRIPTION:
       This module provides a means of asking user for continuing to *
    save a predicate file(.pro) or to abort this function.
    ALGORITHM:
   PASSED VARIABLES:
   GLOBAL VARIABLES USED: *box_rootnode, *header_rootnode
   GLOBAL VARIABLES CHANGED: None
   FILES READ:
   FILES WRITTEN:
   HARDWARE INPUT:
   HARDWARE OUTPUT:
   MODULES CALLED: get_filename_for_predicates()
    CALLING MODULES: make_windows()
   AUTHOR: Intack Kim
   HISTORY:
   ABSTRACT DATA TYPE:
    ORDER OF:
**********************
 void
 save_predicates()
  {
   extern put_message(), enable_input_window();
   extern disable_input_window();
   extern struct box_struct *box_rootnode;
   extern struct header_struct *header_rootnode;
   if(box_rootnode == NULL) {
    put_message(1,"FATAL: Can't save this empty diagram
    -- Make another selection.");
     disable_input_window();
     return:
   if(strcmp(header_rootnode->title.text_string,"") == 0) {
     disable_input_window();
     put_message(1,"NO TITLE: Please enter the TITLE
```

## IDEF<sub>0</sub> Syntax Expert System

This section presents the source code documentations for the IDEF<sub>0</sub> Syntax Expert System and contains the inference engine(ISES) and two rule bases: Activity IDEF<sub>0</sub> Syntax Rules and Boundary IDEF<sub>0</sub> Syntax Rules.

## Inference Engine

```
* DATE: 25 Feb. 1990
 * VERSION: 1.1
 * NAME: BC3
 * DESCRIPTION: The purpose of this module is to provide an
                inference engine for checking IDEFO syntax.
                BC3 provides a means of a shell for backward
                chaining control strategy.
 * OPERATING SYSTEM: UNIX 4.3
 * LANGUAGE: Quintus Prolog
 * CONTENTS: *
 * AUTHOR: DR. Frank M. Brown
 * HISTORY: Version 1.0 - MS-DOS version(DR. Frank M. Brown)
            Version 1.1 - UNIX 4.3 version(Intaek Kim)
       ******************
/*
                                                               */
                              BC3
/*
                                                               */
/ ±
                                                               */
         A shell for backward-chaining expert systems.
                                                               */
/*
/*
                                                               */
/* Each item of knowledge is represented by a triple, i.e.,
                                                               */
   a three-element list of the form [Object, Attribute, Value]. */
/*
/*
                                                               */
/* An associated rule-base supplies the following data:
                                                               */
/*
                                                               */
/* 1. A goals-statement, in the form of a list of triples to
                                                               */
/*
       be solved in sequence. The solved triples are printed
                                                               */
/*
       by the shell.
                                                               */
/* 2. A collection of if-then rules for triples.
                                                               */
/* 3. A collection of 'fact' triples, i.e., triples asserted
                                                              */
                                                               */
/*
       as known a priori.
/* 4. A collection of 'askable' triples, indicating the forms */
       of triples whose values may be obtained from the user.
                                                               */
/* 5. A collection of 'keep' triples, indicating the form of
       the triples not to be erased from working memory at
                                                               */
/*
                                                               */
/*
       the beginning of a new session.
                                                               */
/*
/* Each item of knowledge stored in working memory is of the
                                                               */
   form confirmed([Obj,Attr,Val]) or denied([Obj,Attr,Val]).
                                                               */
/*
/*
                                                               */
/* To use the system, load BC3, load the appropriate rule-
                                                               */
```

```
/* base and type 'start.' Because BC3's operator-defini-
                                                    */
/* tions are used by the rule-bases, BC3 must load first.
                                                    */
                                                    */
/*----*/
                                                    */
/* The operators defined below enable the rules in the know- */
/* ledge-base to be expressed in a form more readable than
                                                    */
/* the standard (prefix) form.
                                                    */
                                                    */
/*----*/
?-op(250, xfx, ::).
?- op(245, xfx, then).
?- op(240, fx, if).
?- op(235, xfx, derived_from).
?- op(230, xfy, or).
?- op(225, xfy, and).
?- op(220, fy, not).
/*----*/
/*
                                                    */
/* The procedure 'start' begins by erasing from working mem- */
/* ory all 'confirmed' and 'denied' clauses, except those
                                                    */
/* clauses protected by 'keep' from erasure. The list of
                                                    */
/* goal-triples is then read from the rule-base and solved in */
/* turn by 'solve'. A trace is maintained of the back-
/* chaining search-tree generated in solving the goals. When */
/* the last of the goal-triples is solved, the values of all */
/* goals, except those solved by asking the user directly,
                                                    */
/* are displayed; the trace is also displayed, if requested,
                                                    */
/* as a "how" explanation of the solution.
                                                    */
                                                    */
/*
/*-----/
erase_working_memory :-
                             /* Erase all working-mem- */
 ( confirmed(Triple),
                             /* ory elements not pro- */
   not(keep::Triple),
   retract(confirmed(Triple))
                              /* tected by 'keep' state-*/
                              /* ments in the knowledge-*/
   denied(Triple),
                              /* base.
                                                    */
   not(keep::Triple),
   retract(denied(Triple))),
 fail.
```

```
erase_working_memory.
start :-
  ask_about_verbose,
  fail.
start :-
  ask_about_checking_type,
  fail.
start :-
  retractall(why_trace(_)),
                                   /* Erase the "why" trace. */
  goals:: Goals,
                                   /* Find the goal-triples, */
  prefix(Goals, PrefixedGoals),
                                   /* prefix each of them
                                   /* with the word 'goal',
  show_head_message,
  solve(Goals,[],PartTrace),
                                   /* satisfy all of the
                                   /* goals and then put the */
  !,nl,
  append(PrefixedGoals, PartTrace, Trace),
                                   /* list of goals at the
                                                             */
                                   /* front of the "how"
  ask_about_trace(Trace),
                                   /* trace. Supply a "how" */
  ask_about_saving_working_memory,
                                   /* explanation on request.*/
  erase_working_memory,
  start_message.
                                   /* If all triples can't
start :-
                                   /* be solved, announce it.*/
  write('I can''t solve this problem.'),nl,
  start_message.
/*----*/
/*
                                                             */
/* The predicate 'solve(Goals,Trace,New_trace)' means that
                                                             */
/* Goals is a list of goals (expressed as triples), and that
/* Trace and New_trace are, respectively, the trace-lists be- */
/* and after solution of the goal at the head of the goal-
/* list. The procedure 'solve' solves each of the goals in
                                                             */
/* turn. The first step in solving a goal is to erase the
/* "why" trace and to initialize it with that goal. Thus each */
/* goal is solved with a separate "why" trace. As each rule
/* is encountered in descending through the search-tree for a */
/* given goal, that rule is added to the front of the "why"
                                                             */
/* trace.
                                                             */
                                                             */
/*
```

```
solve([],Trace,Trace).
solve([Goal|Others],Trace,NewTrace) :-
 retractall(why_trace(_)), /* Initialize the "why"
                                                             */
  asserta(why_trace([goal::Goal])), /* trace.
                                                             */
  is_known(Goal,Trace,Trace1),
  ( confirmed(Goal).!
                                 /* Write each triple as
   /* it's solved, but don't */
nl,write_triple(Goal),nl ), /* write a triple that's */
 solve(Others,Tracei,NewTrace).
                                   /* been told explicitly */
                                   /* by the user.
write_triple([Obj,Attr,Val]) :-
 writelist([Obj,'',Attr,'',Val,'.']).
/*----*/
/*
                                                            */
/* The 'is_known' procedure maintains a trace of the path of */
/* the solution-tree leading to the triple currently under
                                                            */
/* consideration. 'is_known(Triple,Trace,NewTrace)' means
                                                            */
/* that if reasoning to a certain point has been recorded in */
/* the list 'Trace', then the additional triple 'Triple' is
                                                            */
/* known via reasoning recorded by the list 'NewTrace'.
                                                            */
                                                            */
/* A triple is not known if it has been denied by the user.
is_known(Triple,Trace,Trace) :-
 denied(Triple),
  ŀ,
 fail.
/* A triple is known if it is already logged in the trace.
is_known([0,A,V],Trace,[in_trace::[Tag,[0,A,V]]|Trace]) :-
 member(Tag::[0,A,V],Trace),
 Tag \== confirmed_not,
 !.
/* A triple is known if it has been confirmed by the user.
                                                            */
is_known(Triple,Trace,[was_told::Triple|Trace]) :-
```

```
confirmed(Triple).
```

```
/* A triple is known if it is a fact in the rule-base.
                                                                */
is_known(Triple,Trace,[fact::Triple|Trace]) :-
  fact:: Triple.
/* A triple [X,P,Y] is known if the Prolog goal P(X,Y) suc-
                                                                */
/* ceeds, either because P is a built-in predicate, or because */
/* the rule-base has prolog-code defining P. The triple
                                                                */
/* [2,member,[1,2]], for example, is converted into the goal
                                                                */
/* member(2,[1,2]), which is then executed by Prolog. To keep
                                                                */
/* non-Prolog-programmers out of trouble, the triple [X,is,Y]
                                                                */
/* is trapped so that it will not be executed as an arithmetic */
/* statement. The triple [X,:=,Y] is interpreted as Prolog's
                                                                */
/* arithmetic or assignment goal, X is Y.
                                                                */
                                                                */
/* This kind of triple, which runs off and does a computation, */
/* is called a "procedural attachment", or "demon."
                                                                */
is_known([Obj,Attr,Val],Trace,[solved::[Obj,Attr,Val]|Trace]) :-
                               Attr must be a legal functor.
  atom(Attr),
                         /*
    Attr == :=, !,
                         /* Interpret ':=' as Prolog's 'is'.
    Obj is Val
                                                                */
    not (check_reserved_words(Attr)),
 /* Interpret everything else, except */
    T = .. [Attr,Obj,Val],/* reserved words for rule-base, as a */
 /* functor on a two-place
    T, !
                         /* predicate to be solved as a goal. */
  ).
                                                                */
/* A triple is known if it is the head of a rule and the con-
/* ditions of the rule are satisfied. We put a rule that we
                                                                */
/* encounter at the head of the "why" trace, erasing any du-
                                                                */
/* plicates of the rule that are already in the "why" trace.
                                                                */
/* The "why" trace is maintained in the database, in a clause
                                                                */
/* of the form 'why_trace(<List of goals and rules>)'. This
                                                                */
                                                                */
/* differs from the "how" trace, which is handed as an argu-
                                                                */
/* ment from goal to goal.
is_known(Triple,Trace,[was_proved::[Triple,Rule]|Trace]) :-
 member(Rule:: Triple derived_from _Conds, Trace),
  ! .
```

```
is_known(Trpl,Trc,[Rule:: Trpl derived_from Conds|Trc1]) :-
  Rule:: if Conds then Trpl,
  ( verbose.
    writelist(['Trying ', Rule, ':: ', Trpl]),nl
    not verbose),
  why_trace(WhyTrace),
  remove(Rule:: Trpl derived_from Conds, WhyTrace, PartWhy),
  append([Rule:: Trpl derived_from Conds], PartWhy, NewWhy),
  retract(why_trace(_)),
  asserta(why_trace(NewWhy)),
  is_known(Conds,Trc,Trc1),
  ( verbose,
    writelist(['Proved ', Rule, ':: ', Trpl]),nl
   not verbose ),
  !.
/* A condition involving "and", "or", or "not" is known if its */
/* parts are known in suitable combinations.
is_known(Triples1 and Triples2, Trace, Trace2) :-
  is_known(Triples1, Trace, Trace1),
  is_known(Triples2, Trace1, Trace2).
is_known(Triples1 or _Triples2, Trace, Trace1) :-
  is_known(Triples1,Trace,Trace1).
is_known(_Triples1 or Triples2,Trace,Trace2) :-
  is_known(Triples2,Trace,Trace2).
is_known(not Triple,Trace,[confirmed_not::Triple|Trace]) :-
  not is_known(Triple,Trace,_Trace1).
/* A triple is known if (a) the rule-base classifies it as
                                                                */
/* "askable" and if (b) the user confirms it. The user may
                                                                */
/* request a "why" explanation before responding to the ques-
                                                                */
/* tion.
                                                                */
is_known([0,A,V],Trace,[was_told::[0,A,V]|Trace]):-
                       /* 'ask_about' causes the side-effect
                                                                 */
  askable:: [0,A,_],
  ask_about([0,A,V]), /* of confirming or denying [0,A,V] in
                                                                */
                       /* working memory. The clause succeeds */
  confirmed([0,A,V]). /* if the triple was confirmed.
                                                                */
```

```
/*----*/
/* If the user is asked about a triple [0,A,V] in which V is
                                                             */
/* a variable, we assume that only one value of V is allowed
                                                             */
/* for that triple. The askable-fact in the rule-base is to
                                                             */
/* have the form 'askable::[0,A,LegalVals]',where LegalVals is */
/* either a string describing legal values or a list enumer-
/* ating such values. When the user supplies a legal value for */
/* V, the triple is confirmed in working memory.
                                                             */
ask_about([Obj,Attr,Val]) :-
 var(Val),
  !,
 not confirmed([Obj,Attr,_]),
 nl, writelist([Obj,' ',Attr,'? ']),nl,
 askable:: [Obj,Attr,LegalValues],
 write('Legal values: '), write(LegalValues), nl,
 write('> '), read(Reply),
  ( /* If the user replies 'why.', give him an explanation and */
   /* ask again for a value.
     means (Reply, why),
     explain_why([Obj,Attr,Val]),
     ask_about([Obj,Attr,Val])
   )
   /* If LegalValues is a list, check that the reply is in
                                                             */
   /* the list.
                                                             */
     atom(LegalValues)
                                  /* LegalValues is a string.*/
     LegalValues = [_|_],
                                  /* LegalValues is a list. */
     member(Reply, LegalValues)
   ),
   assertz(confirmed([Obj,Attr,Reply]))
   write('Please re-enter your reply.'), nl,
   ask_about([Obj,Attr,Val])
 ).
/* If we get to this clause, the user is being asked to reply */
/* yes or no concerning a triple [0,A,V] in which V is not a
```

```
/* variable. For given 0 and V, working memory may store more */
/* than one triple, confirmed or denied, having different val- */
/* ues of V.
ask_about([Obj,Attr,Val]) :-
 not confirmed([Obj,Attr,Val]),
 not denied([Obj,Attr,Val]),
 writelist([Obj,'',Attr,'',Val,'? (yes./no./why.)']),
 nl, write('> '), read(Reply),
   means(Reply, yes),
   assertz(confirmed([Obj,Attr,Val])), !
   means(Reply,no),
   assertz(denied([Obj,Attr,Va.])), !
   means(Reply, why),
   explain_why([Obj,Attr,Val]),
   ask_about([Obj,Attr,Val])
   write('Please re-enter your reply.'), nl,
   ask_about([Obj,Attr,Val])
 ).
/*----*/
ask_about_verbose :-
 retractall(verbose),
 write(' Question: Do you want verbose operation(y./n.)?'),
 read(Reply), nl, nl,
 means (Reply, yes),
 assert(verbose).
/*----*/
ask_about_checking_type :-
 write(' Question:
 Do you wish to check ACTIVITY BOX, BOUNDARY ARROWS '), nl,
                  or to have HELP MESSAGES ?'), nl, nl,
 write('
               To check ACTIVITY BOX
                                      -> Enter a.'),nl,
 write('
                To check BOUNDARY ARROWS -> Enter b.'), nl,
 write('
 write('
               To have HELP MESSAGE -> Enter h.'), nl,
```

```
write('Choice : '),
 read(Reply), nl,
 ( ( Reply == a,
     load_sarule('ACTIVITYSARULE.PRO'),
     load_working_memory('CHECKBOX.PRO'),
     !)
   ( Reply == b,
     load_sarule('BOUNDARYSARULE.PRO'),
     load_working_memory('CHECKBOUNDARY.PRO'),
   ( Reply == h,
    help_messages,
     ask_about_checking_type )
   ( write(' Please re-enter your choice!!!!'),nl,nl,
     ask_about_checking_type )
 ).
/*----*/
ask_about_trace(Trace) :-
 nl,nl,
 write(' Question: Do you wish to see how this answer '), nl,
 write('
                was arrived at(y./n.)? '),
 read(Reply),
 ( means(Reply, yes), !,
   write_trace(Trace)
   true ).
/*----*/
ask_about_saving_working_memory :- nl,
 write('/************!!! WORNING !!!****************/'),nl,
 write('/* After this session, all working memory elements will */'),nl,
 write('/* be erased except for elements being protected by
 write('/* keep statements in the knowledge base.
 write(' Question: Do you wish to save the current working memory'),nl,
 write('in a file(y./n.)?'),
 read(Reply), nl,
```

```
( means(Reply, yes),
    save_working_memory,
    erase_working_memory, !
    erase_working_memory ).
/*-----*/
explain_why(Triple) :-
  why_trace(WhyTrace),
  write('Because::'),nl,
  justify(Triple, WhyTrace).
justify(Triple, WhyTrace) :-
  member(goal::Goal, WhyTrace),
  Triple = Goal,
  writelist(['This will satisfy the goal ',Goal]),nl,
  nl,
  1.
justify(Triple, WhyTrace) :-
  member(Rule::Head derived_from Cs, WhyTrace),
  ( among(Triple,Cs),
   writelist(['I can use ',Triple]),nl
   among(not Triple,Cs),
   writelist(['I can use NOT ',Triple]),nl
  remove(Rule::Head derived_from Cs, WhyTrace, NewTrace),
  list_known_triples(Cs),
                  to help satisfy ',Rule,':: ',Head]),nl,nl,
  writelist(['
  justify(Head, NewTrace).
list_known_triples(Cs) :-
  among(Triple,Cs),
   confirmed(Triple)
   fact:: Triple
  ).
 writelist(['
                  knowing ',Triple]),nl,
list_known_triples(_).
among(Triple,Conditions) :-
 Triple = Conditions.
```

```
among(Triple, FirstTriple and OtherConditions) :-
 Triple = FirstTriple
 among(Triple,OtherConditions).
among(Triple, FirstTriple or OtherConditions) :-
 Triple = FirstTriple
 among(Triple,OtherConditions).
                        /* Diagnostic utilities for
                                                          */
why :-
 why_trace(Trace),
                       /* the why-trace.
                                                           */
 write_trace(Trace).
list_why :-
 why_trace(Trace),
 member(M, Trace),
 write(M),nl,nl,
 fail.
why_candidates :-
 why_trace(Trace),
 member(_Rule:: Head derived_from Cs,Trace),
  among(Triple,Cs),
 write('Head = '), write(Head), nl,
 write('Triple = '), write(Triple), nl, nl,
 fail.
/*----*/
write_trace([]) :-
 nl.
write_trace([Tag::[0,A,V]|Rest]) :-
  ( Tag == goal, !, write('GOAL:: ')
   Tag == fact, !, write('FACT::
                                     ')
   Tag == solved, !, write('SOLVED:: ')
                                    ')
   Tag == was_told, !, write('TOLD::
   Tag == confirmed_not, !, write('CONTRADICTED:: ')
 write([0,A,V]),nl,
 write_trace(Rest).
```

```
write_trace([in_trace::[Tag,Triple]|Rest]) :-
  ١,
  write('KNOWN:: '), write(Tag), write(':: '), write(Triple), nl,
  write_trace(Rest).
write_trace([was_proved::[Triple,Rule]|Rest]) :- !,
  write('PROVED:: '), write(Triple), write(' using '), write(Rule), nl,
  write_trace(Rest).
write_trace([Rule:: Triple derived_from Conditions|Rest]) :- !,
  writelist([Rule,':: ',Triple,' Was Derived From']),nl,
  write_conditions(Conditions).
  write_trace(Rest).
write_trace([X[Rest]) :-
  write(X),nl,
  write_trace(Rest).
write_conditions([X,Y,Z]) :-
  tab(8), write([X,Y,Z]),nl.
write_conditions(not [X,Y,Z]) :-
  tab(4),write('NOT '),write([X,Y,Z]),nl.
write_conditions([X,Y,Z] and Conditions) :- !,
  tab(8), write([X,Y,Z]), write('AND'), nl,
  write_conditions(Conditions).
write_conditions(not [X,Y,Z] and Conditions) :- !,
  t \(\b(4)\), \(\text{write('NOT')}\), \(\text{write([X,Y,Z])}\), \(\text{write('AND')}\), \(\text{nl}\),
  write_conditions(Conditions).
write_conditions(Conditions1 and Conditions2) :-
  write_conditions(Conditions1),tab(8),write('AND'),nl,
  write_conditions(Conditions2).
write_conditions([X,Y,Z] or Conditions) :- !,
  tab(8), write([X,Y,Z]), write('OR'), nl,
  write_conditions(Conditions).
write_conditions(Conditions1 or Conditions2) :-
  write_conditions(Conditions1), tab(8), write('OR'), nl,
  write_conditions(Conditions2).
write_conditions(not [X,Y,Z] or Conditions) :-
  tab(4),write('NOT'),write([X,Y,Z]),write('OR'),nl,
  write_conditions(Conditions).
/*---- FILE I/O -----
get_filename(Filename) :-
  write('Please supply a filename: '),
  read(Filename).
```

```
load_sarule(SAruleType) :-
 retractall(_::_).
 see(SAruleType),
 load_file,
 seen.
load_file :-
 read(Term),
 load(Term).
load(end_of_file) :- !.
load(Term) :-
 assertz(Term).
 load_file.
load_working_memory(CheckFileType) :-
 retractall(confirmed(_)),
 retractall(denied(_)),
 see(CheckFileType),
 load_file,
 seen.
save_working_memory :-
 get_filename(Filename),
 tell(Filename),
 save_wme,
 told.
save_wme :-
 confirmed(Triple),
 writeq(confirmed(Triple)),
 write('.'),nl,
 fail.
save_wme :-
 denied(Triple),
 writeq(denied(Triple)),
 write('.'),nl,
 fail.
save_wme.
/*----*/
check_reserved_words(Attr) :-
```

```
member(Attr,[' ',is,input_is,has_input_number,output_is,
 has_output_number.control_is,has_control_number,
 mechanism_is, has_mechanism_number, number_is,
 should_be,has_number]).
show_head_message :-
 nl,nl,
  write(' /********************************/'),
 nl.
help_messages:-
            reconsult('HELP.PRO').
not(Predicate) :-
   call(Predicate), !, fail.
not(_).
writelist([]).
writelist([X|L]) :-
  write(X).
  writelist(L).
member(X,[X]_]).
member(X,[_|L]) :-
  member(X,L).
append([],L,L).
append([X|L],M,[X|N]) :-
  append(L,M,N).
remove(_,[],[]).
remove(X,[X|L],M) :-
  ١,
  remove(X,L,M).
remove(X,[Y|L],[Y|M]) :-
  remove(X,L,M).
prefix([],[]).
prefix([Goal|Goals],[goal::Goal|PrefixedGoals]) :-
  prefix(Goals,PrefixedGoals).
list_working_memory :-
  confirmed(Triple),
  write(confirmed(Triple)), write('.'), nl,
  fail.
```

```
list_working_memory :-
 denied(Triple),
 write(denied(Triple)),
 write('.'),nl,
 fail.
list_working_memory.
means(Reply,yes) :-
 member(Reply,[y,yes]).
means(Reply,no) :-
 member(Reply, [n, no]).
means(Reply, why) :-
 member(Reply,[why,w]).
start_message :-
  write('/*
                                                     */'),nl,
                                                     */'),nl,
  write('/*
            WELCOME TO IDEFO SYNTAX EXPERT SYSTEMS
                                                     */'),nl,
  write('/*
  write('/* I.Type start. to begin a new session.
                                                     */'),nl,
  write('/*
                                                     */'),nl,
  write('/* II. Answer all questions using lower case, ending with*/'),nl,
  write('/*
                                                     */'),nl,
           a period.
                                                     */'),nl,
  write('/*
  write('/* III. Type halt. to exit prolog session.
                                                     */'),nl,
                                                     */'),nl,
  nl.
```

?- start\_message.

## Activity IDEF<sub>0</sub> Syntax Rules

```
* DATE: 25 Apr. 1990
* VERSION: 1.0
* NAME: ACTIVITYSARULE.PRO
* TITLE: Rule base for an activity box
* COORDINATOR: Intack Kim
* PROJECT: Knowledge base
* OPERATING SYSTEM: UNIX 4.3
* LANGUAGE: Quintus Prolog
* FILE PROCESSING: This module should be used with an inference
               engine, BC3.
* CONTENTS: Rules for checking IDEFO syntax of an activity box.
* HISTORY:
************************
/* These lists of goal present the resulting message. */
 goals:: [['Name', '', _Name],
                             /* Goal for NAME of the box */
  ['Input', '', _Input], /* Goal for INPUT of box
        ['Output', '', _Output], /* Goal for OUTPUT of box
        ['Control', '', _Control],/* Goal for CONTROL of box */
        ['Mechanism', '', _Mechanism], /* Goal for MECHANISM */
        ['Number', '', _Number]]. /* Goal for the box NUMBER */
rule1 :: if [activityname, is, '']
 then ['Name', '', ' --> ERROR: No Activity Name.
   Each box must have an activity name'].
rule2 :: if [activityname, is, Activity]
   and [Activity, \==, '']
       then ['Name', '', ' --> CORRECT: Activity Name is OK'].
/************* About Input ***************/
rule3 :: if [activityname, is, Activity]
         and [Activity, input_is, null]
       then ['Input', '', '--> CORRECT: No Input Arrows, however,
             Input is OK'].
rule4 :: if [activityname, is, Activity]
```

```
and [Activity, input_is, '']
         then ['Input', '', ' --> ERROR: No Input Label
Each Input arrow must have an input label'].
rule5 :: if [activityname, is, Activity]
   and [Activity, has_input_number, InputNumber]
           and [InputNumber, >, 5]
        then ['Input', '', ' --> RECOMMEND:
               You would better reduce the number of Input arrows
               from 0 to 5'].
rule6 :: if [activityname, is, _Activity]
        then ['Input', '', ' --> CORRECT: Input is OK'].
/************* About Output **************/
/* If there is a box and the output of the box is empty, */
/* then there is no output/name.
                                                       */
rule7 :: if [activityname, is, Activity]
            and [Activity, output_is, null]
     then ['Output', '', ' --> ERROR: You should have at least
            one output arrow'].
rule8 :: if [activityname, is, Activity]
            and [Activity, output_is, '']
        then ['Output', '', ' --> ERROR: No Output Label.
               Each Output Arrow should have an output Label'].
rule9 :: if [activityname, is, Activity]
           and [Activity, has_output_number, OutputNumber]
   and [OutputNumber, >, 5]
        then ['Output', '', ' --> RECOMMEND:
       You would better reduce the number of Output arrows
               from 1 to 5'].
rule10 :: if [activityname, is, _Activity]
            then ['Output', '', ' --> CORRECT: Output is OK'].
/************* About Control ***************/
rule11 :: if [activityname, is, Activity]
            and [Activity, control_is, null]
   then ['Control', '', ' --> ERROR: You should have at least
          one control arrow'].
rule12 :: if [activityname, is, Activity]
```

```
and [Activity, control_is, '']
         then ['Control', '', ' --> ERROR: No Control Label.
                Each Control Arrow should have a control Label'].
rule13 :: if [activityname, is, Activity]
           and [Activity, has_control_number, ControlNumber]
    and [ControlNumber, >, 5]
        then ['Control', '', ' --> RECOMMEND:
     You would better reduce the number of Control arrows
             from 1 to 5'7.
rule14 :: if [activityname, is, _Activity]
         then ['Control', '', ' --> CORRECT: Control is OK'].
/************ About Mechanism ************/
rule15 :: if [activityname, is, Activity]
            and [Activity, mechanism_is, '']
         then ['Mechanism', '', '--> ERROR: No Mechanism Label.
             Each Mechanism Arrow should have a mechanism Label'].
rule16 :: if [activityname, is, Activity]
            and [Activity, mechanism_is, null]
  then ['Mechanism', '', '--> CORRECT: No Mechanism Arrows, however,
         Mechanism is OK'].
rule17 :: if [activityname, is, Activity]
    and [Activity, has_mechanism_number, MechanismNumber]
           and [MechanismNumber, >, 5]
        then ['Mechanism', '', '--> RECOMMEND:
               You would better reduce the number of Mechanism arrows
       within 0 to 5'].
rule18 :: if [activityname, is, _Activity]
         then ['Mechanism', '', '--> CORRECT: Mechanism is OK'].
/****************** Activity Number **************/
/* If there is a box and the number of the box is 0, that is, */
/* the box has no number, the box must be the top most box. */
rule19 :: if [title, is, Activity]
     and [Activity, number_is, Number]
     and [Number, ==, 0]
          then ['Number', '', '--> CORRECT: Activity number is OK.
 This activity must be the top most level box'].
```

```
rule20 :: if [title, is, Activity]
      and [Activity, number_is, Number]
      and [Number, \==, 0]
           then ['Number', '', '--> ERROR: This activity must be
 the top most level box.
 This activity don't need a box number'].
/* Each activity box must have a number for the box within */
/* 1 through 6 except for the top most level activity box. */
/* If there is a box and the number of the box is 0, then */
/* there is no number in the box.
rule21 :: if [Activity, number_is, Number]
               and [title, is, ParentActivity]
       and [ParentActivity, \==, Activity]
       and [Number, ==, 0]
           then ['Number', '', '--> ERROR: Activity box has no number.
        The activity box should have a box number from 1 to 6.
                 (If the activity box is the top most level one, then
                  ignore this message)'].
rule22 :: if [Activity, number_is, _Number]
       and [title, is, ParentActivity]
       and [ParentActivity, \==, Activity]
       and [activity_number, is, in_legal_range]
            then ['Number', '', '--> CORRECT: Activity box number
   is OK.'].
rule23 :: if [Activity, number_is, _Number]
       and [title, is, ParentActivity]
       and [ParentActivity, \==, Activity]
       and [activity_number, is, in_illegal_range]
        then ['Number', '', ' --> ERROR: Activity box number is
not proper.
        The recommended range of number is 1 to 6'].
/* These rules are the second level rules.
rule24 :: if [_Activity, number_is, Number]
       and [Number, >, 0]
               and [Number, <, 7]
            then [activity_number, is, in_legal_range].
rule25 :: if [_Activity, number_is, Number]
     and [Number, <, 0]
```

then [activity\_number, is, in\_illegal\_range].

## Boundary IDEF<sub>0</sub> Syntax Rules

```
/*****************
 * DATE: 25 Apr. 1990
 * VERSION: 1.0
 * FILE NAME: BOUNDARYSARULE.PRO
 * TITLE: Rule base for boundary arrows
 * COORDINATOR: Intack Kim
* PROJECT: Knowledge base
* OPERATING SYSTEM: UNIX 4.3
 * LANGUAGE: Quintus Prolog
 * FILE PROCESSING: This module should be used an inference engine *
                   BC3.
 * CONTENTS: Rules for checking IDEFO syntax of boundary arrows.
 * HISTORY:
/************** GOALS ****************/
/* These lists of goal present the resulting message. */
goals:: [['Boundary Input', ' ', _Input],
  /* Goal for boundary input in any IDEFO diagram */
         ['Boundary Output', '', _Output],
  /* Goal for boundary output in any IDEFO diagram */
         ['Boundary Control', '', _Control],
  /* Goal for boundary control in any IDEFO diagram */
         ['Boundary Mechanism', '', _Mechanism]
  /* Goal for boundary mechanism in any IDEFO diagram */
 ].
/****************** RULES **************/
/**** When there is no information of parent box ***/
rule1 :: if [child_title, is, ParentBox]
            and not [activityname, is, ParentBox]
         then [boundarysarule, is, stalled].
rule2 :: if [boundarysarule, is, stalled]
         then ['Boundary Input', '', '
 --> !!! THIS IS A FATAL ERROR !!!'].
rule3 :: if [boundarysarule, is, stalled]
         then ['Boundary Output', '', '
--> There is nothing about Parent activity'].
rule4 :: if [boundarysarule, is, stalled]
```

```
then ['Boundary Control', '', '
 --> Maybe you have tried to check syntax with
         a file without PARENT ACTIVITY BOX information'].
rule5 :: if [boundarysarule, is, stalled]
          then ['Boundary Mechanism', '', '
 --> PLEASE START AGAIN !!!'].
/************ About Boundary Input ************/
/* No label on boundary input arrow */
rule6 :: if [boundary_input, is, '']
          then ['Boundary Input', '','
 --> ERROR: No boundary input label'].
/* No label on input arrow of the parent box */
rule7 :: if [child_title. is. ParentBox]
            and [ParentBox, input_is, '']
        then ['Boundary Input',' ','
--> ERROR: Parent Input has no label'].
     The number of Input arrow(s) of the Parent Activity */
/* box is differ from that of the Boundary Input arrow(s).*/
    The number of the Parent Input arrow(s) > The number */
/* of the Boundary Input arrow(s).
                                                          */
rule8 :: if [child_title, is, ParentBox]
            and [ParentBox, has_input_number, InputNumber]
            and [boundary_input, has_number, BoundInNumber]
            and [InputNumber, >, BoundInNumber]
         then ['Boundary Input',' ','
--> ERROR: The number of Input arrow(s) of
               Parent Activity box is greater than that of
         Boundary Input arrow(s) -- Must have the same number'].
    The number of Input arrow(s) of the Parent Activity box is */
/* differ from that of the Boundary Input arrow(s).
                                                                */
/* The number of the Parent Input arrow(s) < The number of the */
                                                                */
/* Boundary Input arrow(s).
rule9 :: if [child_title, is, ParentBox]
            and [ParentBox, has_input_number, InputNumber]
            and [boundary_input, has_number, BoundInNumber]
            and [InputNumber, <, BoundInNumber]</pre>
        then ['Boundary Input',' ','
--> ERROR: The number of Input arrow(s) of
               Parent Activity box is less than that of
         Boundary Input arrow(s) -- Must have the same number'].
```

```
If the number of arrows is greater than 1 've, recomend */
/* about how many the number of arrows exists.
rule10 :: if [boundary_input, has_number, BoundInNumber]
           and [child_title, is, ParentBox]
           and [ParentBox, has_input_number, BoundInNumber]
           and [BoundInNumber, >, 5]
         then ['Boundary Input', '', '
                                          --> RECOMMEND:
        You would better reduce the number of arrows to six
        below'].
/* No boundary Input arrow and No Input at Parent Box
                                                                */
rule11 :: if [child_title, is, ParentBox]
           and [boundary_input, is, null]
           and [ParentBox, input_is, null]
         then ['Boundary Input', '','
--> CORRECT: Boundary Input is OK'].
/* Consider the correct case acording to the Number of */
/* Input arrow(s).
rule12 :: if [boundary_input, has_number, BoundInNumber]
             and [child_title, is, ParentBox]
             and [ParentBox, has_input_number, BoundInNumber]
        then [case_of_boundary_in, is, BoundInNumber].
/* Case 1: The number of Boundary Input arrow is 1. */
rule13 :: if [case_of_boundary_in, is, 1]
             and [child_title, is, ParentBox]
             and [ParentBox, input_is, ParentInput]
             and [boundary_input1, is, ParentInput]
        then ['Boundary Input', '','
--> CORRECT: Boundary Input is OK'].
/* Case 2: The number of Boundary Input arrows is 2. */
rule14 :: if [case_of_boundary_in, is, 2]
             and [boundary_input1, is, ParentInput1]
             and [boundary_input2, is, ParentInput2]
             and [child_title, is, ParentBox]
             and [ParentBox, input_is, ParentInput1]
             and [ParentBox, input_is, ParentInput2]
        then ['Boundary Input', '','
--> CORRECT: Boundary Input is OK'].
/* Case 3: The number of Boundary Input arrows is 3. */
rule15 :: if [case_of_boundary_in, is, 3]
```

```
and [boundary_input1, is, ParentInput1]
             and [boundary_input2, is, ParentInput2]
             and [boundary_input3, is, ParentInput3]
             and [child_title, is, ParentBox]
             and [ParentBox, input_is, ParentInput1]
             and [ParentBox, input_is, ParentInput2]
             and [ParentBox, input_is, ParentInput3]
         then ['Boundary Input', '','
--> CORRECT: Boundary Input is OK'].
 /* Case 4: The number of Boundary Input arrows is 4.
 rule16 :: if [case_of_boundary_in, is, 4]
             and [boundary_input1, is, ParentInput1]
             and [boundary_input2, is, ParentInput2]
             and [boundary_input3, is, ParentInput3]
             and [boundary_input4, is, ParentInput4]
             and [child_title, is, ParentBox]
             and [ParentBox, input_is, ParentInput1]
             and [ParentBox, input_is, ParentInput2]
             and [ParentBox, input_is, ParentInput3]
             and [ParentBox, input_is, ParentInput4]
         then ['Boundary Input', '','
--> CORRECT: Boundary Input is OK'].
 /* Case 5: The number of Boundary Input arrows is 5. */
 rule17 :: if [case_of_boundary_in, is, 5]
             and [boundary_input1, is, ParentInput1]
             and [boundary_input2, is, ParentInput2]
             and [boundary_input3, is, ParentInput3]
             and [boundary_input4, is, ParentInput4]
             and [boundary_input5, is, ParentInput5]
             and [child_title, is, ParentBox]
             and [ParentBox, input_is, ParentInput1]
             and [ParentBox, input_is, ParentInput2]
             and [ParentBox, input_is, ParentInput3]
             and [ParentBox, input_is, ParentInput4]
             and [ParentBox, input_is, ParentInput5]
         then ['Boundary Input', '',' --> CORRECT: Boundary Input
is OK'].
/* Boundary Input is not matched */
/* This rule includes all No_match case though both of */
/* parent and child have the same number of input
                                                        */
/* arrow(s).
                                                        */
rule18 :: if [child_title, is, ParentBox]
```

```
and [activityname, is, ParentBox]
          then ['Boundary Input',' ','--> ERROR: Boundary Input is not
          matched. -- You may have at least one unmatched label'].
/************* About Boundary Output ****************/
 rule19 :: if [boundary_output, is, '']
        then ['Boundary Output', '',' --> ERROR: No boundary output
       label'].
 rule20 :: if [child_title, is, ParentBox]
             and [ParentBox, output_is, '']
        then ['Boundary Output',' ',' --> ERROR: Parent Output has no
       label'].
 rule21 :: if [boundary_output, is, null]
          then ['Boundary Output',' ',' --> ERROR: No boundary output
 arrow.
                Should have at least one boundary output arrow'].
 rule22 :: if [child_title, is, ParentBox]
             and [ParentBox, output_is, null]
          then ['Boundary Output', '', '
 --> ERROR: No parent output arrow.
                Should have at least one parent output arrow'].
    The number of Output arrow(s) of the Parent Activity */
/* box is differ from that of the Boundary Output arrow(s).*/
    The number of the Parent Output arrow(s) > The number */
/* of the Boundary Output arrow(s).
rule23 :: if [child_title, is, ParentBox]
            and [ParentBox, has_output_number, OutputNumber]
            and [boundary_output, has_number, BoundOutNumber]
            and [OutputNumber, >, BoundOutNumber]
         then ['Boundary Output',' ',' --> ERROR: The number of Output
         arrow(s) of Parent Activity box is greater than that of
        Boundary Output arrow(s) -- Must have the same number'].
/* The number of Output arrow(s) of the Parent Activity box is */
/* differ from that of the Boundary Output arrow(s).
                                                                */
/* The number of the Parent Output arrow(s) < The number of the */
/* Boundary Output arrow(s).
                                                                */
rule24 :: if [child_title, is, ParentBox]
           and [ParentBox, has_output_number, OutputNumber]
           and [boundary_output, has_number, BoundOutNumber]
           and [OutputNumber, <, BoundOutNumber]</pre>
```

```
then ['Boundary Output',' ',' --> ERROR: The number of Output
         arrow(s) of Parent Activity box is less than that of
         Boundary Output arrow(s) -- Must have the same number'].
/* If the number of arrows is greater than five, recomend */
/* about how many the number of arrows exists.
                                                             */
rule25 :: if [boundary_output, has_number, BoundOutNumber]
            and [child_title, is, ParentBox]
            and [ParentBox, has_output_number, BoundOutNumber]
            and [BoundOutNumber, >, 5]
         then ['Boundary Output', ''.'
                                         --> RECOMMEND:
         You would better reduce the number of arrows to six
         below'].
/* Consider the correct case acording to the Number of */
/* Output arrow(s).
                                                        */
rule26 :: if [boundary_output, has_number, BoundOutNumber]
            and [child_title, is, ParentBox]
             and [ParentBox, has_output_number, BoundOutNumber]
         then [case_of_boundary_out, is, BoundOutNumber].
 /* Case 1: The number of Boundary Output arrow is 1. */
 rule27 :: if [case_of_boundary_out, is, 1]
            and [boundary_output1, is, ParentOutput]
            and [child_title, is, ParentBox]
            and [ParentBox, output_is, ParentOutput]
         then ['Boundary Output', '',' --> CORRECT:
                        Boundary Output is OK'].
 /* Case 2: The number of Boundary Output arrows is 2. */
rule28 :: if [case_of_boundary_out, is, 2]
            and [boundary_output1, is, ParentOutput1]
            and [boundary_output2, is, ParentOutput2]
            and [child_title, is, ParentBox]
            and [ParentBox, output_is, ParentOutput1]
             and [ParentBox, output_is, ParentOutput2]
         then ['Boundary Output', '',' --> CORRECT:
                        Boundary Output is OK'].
 /* Case 3: The number of Boundary Output arrows is 3. */
 rule29 :: if [case_of_boundary_out, is, 3]
            and [boundary_output1, is, ParentOutput1]
            and [boundary_output2, is, ParentOutput2]
            and [boundary_output3, is, ParentOutput3]
            and [child_title, is, ParentBox]
```

```
and [ParentBox, output_is, ParentOutput1]
            and [ParentBox, output_is, ParentOutput2]
            and [ParentBox, output_is, ParentOutput3]
        then ['Boundary Output', '',' --> CORRECT:
                        Boundary Output is OK'].
/* Case 4: The number of Boundary Output arrows is 4. */
rule30 :: if [case_of_boundary_out, is, 4]
            and [boundary_output1, is, ParentOutput1]
            and [boundary_output2, is, ParentOutput2]
            and [boundary_output3, is, ParentOutput3]
            and [boundary_output4, is, ParentOutput4]
            and [child_title, is, ParentBox]
            and [ParentBox, output_is, ParentOutput1]
            and [ParentBox, output_is, ParentOutput2]
            and [ParentBox, output_is, ParentOutput3]
            and [ParentBox, output_is, ParentOutput4]
        then ['Boundary Output', '',' --> CORRECT:
                        Boundary Output is OK'].
/* Case 5: The number of Boundary Output arrows is 5. */
rule31 :: if [case_of_boundary_out, is, 5]
            and [boundary_output1, is, ParentOutput1]
            and [boundary_output2, is, ParentOutput2]
            and [boundary_output3, is, ParentOutput3]
            and [boundary_output4, is, ParentOutput4]
            and [boundary_output5, is, ParentOutput5]
            and [child_title, is, ParentBox]
            and [ParentBox, output_is, ParentOutput1]
            and [ParentBox, output_is, ParentOutput2]
            and [ParentBox, output_is, ParentOutput3]
            and [ParentBox, output_is, ParentOutput4]
            and [ParentBox, output_is, ParentOutput5]
        then ['Boundary Output', '',' --> CORRECT:
                        Boundary Output is OK'].
/* Boundary Output is not matched */
/* This rule includes all No_match case though both of */
/* parent and child have the same number of output
                                                       */
                                                       */
/* arrow(s).
rule32 :: if [child_title, is, ParentBox]
             and [activityname, is, ParentBox]
         then ['Boundary Output',' ','--> ERROR: Boundary Output is not
         matched. -- You may have at least one unmatched label'].
```

```
/******** About Boundary Control ************/
 rule33 :: if [boundary_control, is, '']
      then ['Boundary Control', '','
      --> ERROR: No boundary control label'].
 rule34 :: if [child_title, is, ParentBox]
             and [ParentBox, control_is, '']
     then ['Boundary Control',' ','
    --> ERROR: Parent Control has no label'].
rule35 :: if [boundary_control, is, null]
          then ['Boundary Control',' ','
 --> ERROR: No boundary control arrow.
                 Should have at least one boundary control arrow'].
rule36 :: if [child_title, is, ParentBox]
             and [ParentBox, control_is, null]
          then ['Boundary Control', '', '
 --> ERROR: No parent control arrow.
                 Should have at least one parent control arrow'].
    If the number of arrows is greater than five, recomend */
/* about how many the number of arrows exists.
                                                             */
 rule37 :: if [boundary_control, has_number, BoundConNumber]
            and [child_title, is, ParentBox]
            and [ParentBox, has_control_number, BoundConNumber]
            and [BoundConNumber, >, 5]
         then ['Boundary Control', '', ' --> RECOMMEND:
         You would better reduce the number of arrows to six
         below'].
/* Consider the correct case acording to the Number of */
                                                        */
/* Control arrow(s).
rule38 :: if [boundary_control, has_number, BoundConNumber]
             and [child_title, is, ParentBox]
             and [ParentBox, has_control_number, BoundConNumber]
         then [case_of_boundary_con, is, BoundConNumber].
 /* Case 1: The number of Boundary Control arrow is 1. */
 rule39 :: if [case_of_boundary_con, is, 1]
             and [boundary_control1, is, ParentControl]
             and [child_title, is, ParentBox]
             and [ParentBox, control_is, ParentControl]
         then ['Boundary Control', '',' --> CORRECT: Boundary
                         Control is OK'].
```

```
/* Case 2: The number of Boundary Control arrows is 2. */
rule40 :: if [case_of_boundary_con, is, 2]
            and [boundary_control1, is, ParentControl1]
            and [boundary_control2, is, ParentControl2]
            and [child_title, is, ParentBox]
            and [ParentBox, control_is, ParentControl1]
            and [ParentBox, control_is, ParentControl2]
        then ['Boundary Control', '',' --> CORRECT: Boundary
                        Control is OK'].
/* Case 3: The number of Boundary Control arrows is 3. */
rule41 :: if [case_of_boundary_con, is, 3]
            and [boundary_control1, is, ParentControl1]
            and [boundary_control2, is, ParentControl2]
            and [boundary_control3, is, ParentControl3]
            and [child_title, is, ParentBox]
            and [ParentBox, control_is, ParentControl1]
            and [ParentBox, control_is, ParentControl2]
            and [ParentBox, control_is, ParentControl3]
        then ['Boundary Control', '',' --> CORRECT: Boundary
                        Control is OK'].
/* Case 4: The number of Boundary Control arrows is 4. */
rule42 :: if [case_of_boundary_con, is, 4]
            and [boundary_control1, is, ParentControl1]
            and [boundary_control2, is, ParentControl2]
            and [boundary_control3, is, ParentControl3]
            and [boundary_control4, is, ParentControl4]
            and [child_title, is, ParentBox]
            and [ParentBox, control_is, ParentControl1]
            and [ParentBox, control_is, ParentControl2]
            and [ParentBox, control_is, ParentControl3]
            and [ParentBox, control_is, ParentControl4]
        then ['Boundary Control', '',' --> CORRECT: Boundary
                        Control is OK'].
/* Case 5: The number of Boundary Control arrows is 5. */
rule43 :: if [case_of_boundary_con, is, 5]
            and [boundary_control1, is, ParentControl1]
            and [boundary_control2, is, ParentControl2]
            and [boundary_control3, is, ParentControl3]
            and [boundary_control4, is, ParentControl4]
            and [boundary_control5, is, ParentControl5]
            and [child_title, is, ParentBox]
```

```
and [ParentBox, control_is, ParentControl1]
             and [ParentBox, control_is, ParentControl2]
             and [ParentBox, control_is, ParentControl3]
             and [ParentBox, control_is, ParentControl4]
             and [ParentBox, control_is, ParentControl5]
         then ['Boundary Control', '',' --> CORRECT: Boundary
                         Control is OK'].
/* Boundary Control is not matched */
/* This rule includes all No_match case though both of */
/* parent and child have the same number of control
                                                        */
/* arrow(s).
                                                        */
rule44 :: if [child_title, is, ParentBox]
              and [activityname, is, ParentBox]
         then ['Boundary Control',' ','-> ERROR: Boundary Control is not
         matched. -- You may have at least one unmatched label'].
/******** About Boundary Mechanism **********/
rule45 :: if [boundary_mechanism, is, '']
    then ['Boundary Mechanism', '','
    --> ERROR: No boundary mechanism label'].
rule46 :: if [child_title, is, ParentBox]
             and [ParentBox, mechanism_is, '']
    then ['Boundary Mechanism',' ','
    --> ERROR: Parent Mechanism has no label'].
    The number of Mechanism arrow(s) of the Parent Activity */
/* box is differ from that of the Boundary Mechanism arrow(s).*/
/* The number of the Parent Mechanism arrow(s) > The number */
/* of the Boundary Mechanism arrow(s).
rule47 :: if [child_title, is, ParentBox]
            and [ParentBox, has_mechanism_number, MecNumber]
            and [boundary_mechanism, has_number, BoundMecNumber]
            and [MecNumber, >, BoundMecNumber]
         then ['Boundary Mechanism',' ','--> ERROR: The number of
         Mechanism arrow(s) of Parent Activity box is greater than that
         of Boundary Mechanism arrow(s) -- Must have the same
        number'].
/* The number of Mechanism arrow(s) of the Parent Activity box is */
/* differ from that of the Boundary Mechanism arrow(s).
/* The number of the Parent Mechanism arrow(s) < The number of the */
/* Boundary Mechanism arrow(s).
                                                                    */
rule48 :: if [child_title, is, ParentBox]
```

```
and [boundary_mechanism, has_number, BoundMecNumber]
            and [MecNumber, <, BoundMecNumber]
         then ['Boundary Mechanism',' ','--> ERROR: The number of
        Mechanism arrow(s) of Parent Activity box is less than
        that of Boundary Mechanism arrow(s) -- Must have the same
        number'].
    If the number of arrows is greater than five, recomend */
/* about how many the number of arrows exists.
                                                             */
rule49 :: if [boundary_mechanism, has_number, BoundMecNumber]
            and [child_title, is, ParentBox]
            and [ParentBox, has_mechanism_number, BoundMecNumber]
            and [BoundMecNumber, >, 5]
        then ['Boundary Mechanism', '', '--> RECOMMEND:
        You would better reduce the number of arrows to six
        below'].
/* No boundary Mechanism arrow and No Mechanism at Parent Box */
rule50 :: if [child_title, is, ParentBox]
            and [boundary_mechanism, is, null]
            and [ParentBox, mechanism_is, null]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
                        Mechanism is OK'].
/* Consider the correct case acording to the Number of */
/* Mechanism arrow(s).
rule51 :: if [boundary_mechanism, has_number, BoundMecNumber]
             and [child_title, is, ParentBox]
             and [ParentBox, has_mechanism_number, BoundMecNumber]
        then [case_of_boundary_mech, is, BoundMecNumber].
/* Case 1: The number of Boundary Mechanism arrow is 1. */
rule52 :: if [case_of_boundary_mech, is, 1]
             and [child_title, is, ParentBox]
             and [ParentBox, mechanism_is, ParentMech]
             and [boundary_mechanism1, is, ParentMech]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
                        Mechanism is OK'].
/* Case 2: The number of Boundary Mechanism arrows is 2. */
rule53 :: if [case_of_boundary_mech, is, 2]
            and [boundary_mechanism1, is, ParentMech1]
            and [boundary_mechanism2, is, ParentMech2]
             and [child_title, is, ParentBox]
```

and [ParentBox, has\_mechanism\_number, MecNumber]

```
and [ParentBox, mechanism_is, ParentMech1]
            and [ParentBox, mechanism_is, ParentMech2]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
                        Mechanism is OK'].
/* Case 3: The number of Boundary Mechanism arrows is 3. */
rule54 :: if [case_of_boundary_mech, is, 3]
            and [boundary_mechanism1, is, ParentMech1]
            and [boundary_mechanism2, is, ParentMech2]
            and [boundary_mechanism3, is, ParentMech3]
            and [child_title, is, ParentBox]
            and [ParentBox, mechanism_is, ParentMech1]
            and [ParentBox, mechanism_is, ParentMech2]
            and [ParentBox, mechanism_is, ParentMech3]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
                        Mechanism is OK'].
/* Case 4: The number of Boundary Mechanism arrows is 4. */
rule55 :: if [case_of_boundary_mech, is, 4]
            and [boundary_mechanism1, is, ParentMech1]
            and [boundary_mechanism2, is, ParentMech2]
            and [boundary_mechanism3, is, ParentMech3]
            and [boundary_mechanism4, is, ParentMech4]
            and [child_title, is, ParentBox]
            and [ParentBox, mechanism_is, ParentMech1]
            and [ParentBox, mechanism_is, ParentMech2]
            and [ParentBox, mechanism_is, ParentMech3]
            and [ParentBox, mechanism_is, ParentMech4]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
                        Mechanism is OK'].
/* Case 5: The number of Boundary Mechanism arrows is 5. */
rule56 :: if [case_of_boundary_mech, is, 5]
            and [boundary_mechanism1, is, ParentMech1]
            and [boundary_mechanism2, is, ParentMech2]
            and [boundary_mechanism3, is, ParentMech3]
            and [boundary_mechanism4, is, ParentMech4]
            and [boundary_mechanism5, is, ParentMech5]
            and [child_title, is, ParentBox]
            and [ParentBox, mechanism_is, ParentMech1]
            and [ParentBox, mechanism_is, ParentMech2]
            and [ParentBox, mechanism_is, ParentMech3]
            and [ParentBox, mechanism_is, ParentMech4]
            and [ParentBox, mechanism_is, ParentMech5]
        then ['Boundary Mechanism', '','--> CORRECT: Boundary
```

## Mechanism is OK'].

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## Vita

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Design Technique (SADT) method. The tool provides the requirements							
analyst and the designer with a means of checking the syntax of IDEFO							
diagram. The tool was implemented using expert system thechnique.							
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